Dinah Rosenberg · Eilon Solan · Nicolas Vieille

# Stopping games with randomized strategies

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**Abstract.** We study stopping games in the setup of Neveu. We prove the existence of a uniform value (in a sense defined below), by allowing the players to use randomized strategies. In constrast with previous work, we make no comparison assumption on the payoff processes. Moreover, we prove that the value is the limit of discounted values, and we construct  $\epsilon$ -optimal strategies.

# 1. Introduction

Dynkin (1969) introduced the following optimization problem. Two players observe stochastic sequences  $(r(n), x(n))_n$ . Player 1 (*resp.* player 2) is allowed to stop whenever  $x(n) \le 0$  (*resp.* x(n) > 0). The two players choose stopping times  $\mu_1$  and  $\mu_2$  which obey this rule, and the payoff is given by

$$\gamma(\mu_1, \mu_2) = \mathbf{E}\{1_{\mu_1 < \mu_2} r(\mu_1) + 1_{\mu_1 > \mu_2} r(\mu_2)\}.$$

The goal of player 1 is to maximize  $\gamma(\mu_1, \mu_2)$ , whereas player 2 tries to minimize  $\gamma(\mu_1, \mu_2)$ . Dynkin proved that this game has a value if  $\sup_n |r(n)| \in L^1$ , and constructed  $\epsilon$ -optimal strategies for the two players.

Kiefer (1971) and Neveu (1975) gave other sufficient conditions for existence of the value in this zero-sum game and in a variant of it. Neveu extended the game by allowing the players to stop simultaneously: a process  $(a_n, b_n, c_n)$  is given (with  $\sup_n \sup(|a_n|, |b_n|, |c_n|) \in L^1$ ), the two players choose stopping times  $\mu_1$  and  $\mu_2$ , and the payoff to player 1 is

$$\mathbf{E}\{a_{\mu_1}\mathbf{1}_{\mu_1<\mu_2}+b_{\mu_2}\mathbf{1}_{\mu_2<\mu_1}+c_{\mu_1}\mathbf{1}_{\mu_1=\mu_2<+\infty}\}$$

He proved that, under the assumption  $a_n = c_n \le b_n$ , the game has a value.

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D. Rosenberg: Laboratoire d'Analyse, Géométrie et Applications, Institut Galilée, Université Paris Nord, avenue Jean-Baptiste Clément, 93430 Villetaneuse, France. e-mail: dinah@math.univ-paris13.fr

E. Solan: MEDS Department, Kellogg Graduate School of Management, Northwestern University, 2001 Sheridan Rd., Evanston, IL 60208, USA, and School of Mathematical Sciences, Tel Aviv University, Israel. e-mail: eilons@post.tau.ac.il

N. Vieille: HEC, Département Economie et Finance, and Laboratoire d'Econométrie de l'Ecole Polytechnique, 1 rue Descartes, 75005 Paris, France. e-mail: vieille @poly.polytechnique.fr

There is a broad literature on continuous time Dynkin games giving sufficient conditions for the existence of the value and optimal strategies: Bismut (1979) proved that under the hypothesis  $a_n = c_n \leq b_n$ , some regularity assumption and Mokobodski's hypothesis (namely that there exist positive bounded supermartingales z and z' satisfying  $a \leq z - z' \leq b$ ) the value exists. The regularity assumption was weakened by Alario-Nazaret, Lepeltier and Marchal (1982), and then Lepeltier and Maingueneau (1984) established the existence of the value and optimal strategies without Mokobodski's hypothesis, assuming only  $a_n = c_n \leq b_n$ .

In the present paper, we focus on discrete time Dynkin games and we allow the players to use randomized stopping times. We prove the existence of the value, under the single integrability condition.

This result is related to a result due to Maitra and Sudderth (1993), for general stochastic games. In such games, the players receive a payoff in each stage. Maitra and Sudderth define the payoff associated to a play as the lim sup of the payoffs received along the play. They prove that such games have a value, provided the payoffs are bounded and deterministic functions of the state.

It is clear that, under some regularity assumptions on the processes  $(a_n)$ ,  $(b_n)$  and  $(c_n)$ , stopping games may be viewed as general stochastic games with a very specific transition structure (note however that boundedness of the payoff function will not be satisfied). Thus, the result of Maitra and Sudderth has some bite in stopping games. We emphasize that our method bears no relation to their approach (which is based on transfinite induction).

Our contribution is threefold. (i) We prove that the value exists under the single integrability requirement, and, moreover, it is uniform in a sense defined below. (ii) We prove that the value is the limit of the so-called discounted values, studied by Yasuda (1985). In particular, it follows that the discounted values converge. (iii) We construct  $\epsilon$ -optimal strategies for the players.

Our method is to construct a strategy for player 1 that guarantees him an expected payoff which is, up to an  $\epsilon$ , the limit of some sequence of discounted values. We provide two different constructions for an  $\epsilon$ -optimal strategy. In the first construction the player plays at each stage an optimal discounted strategy, where the discount factor may change from time to time. In the second construction, which has the flavor of Dynkin's construction, the player plays almost the limit of the optimal discounted strategies.

The paper is arranged as follows. In section 2 we present the model and the main results, in section 3 we introduce few tools, in section 4 we explain the main ideas of the two constructions, and finally, in sections 5.2 and 5.3 we provide the two constructions of  $\epsilon$ -optimal strategies. Section 6 concludes the paper by discussing related issues.

#### 2. The model and the main results

Let  $(\Omega, \mathscr{A}, \mathbf{P})$  be a probability space, and  $(\mathscr{F}_n)$  be a filtration over  $(\Omega, \mathscr{A}, \mathbf{P})$  (the information available at stage *n*). Let  $(a_n), (b_n), (c_n)$  be processes, defined over

 $(\Omega, \mathscr{A}, \mathbf{P})$ . We assume

$$\sup_{n} |a_n|, \sup_{n} |b_n|, \sup_{n} |c_n| \in L^1(\mathbf{P}).$$

$$\tag{1}$$

We also assume that  $(a_n)$ ,  $(b_n)$ , and  $(c_n)$  are adapted. This assumption can be dispensed with. One needs only replace everywhere  $(a_n)$ ,  $(b_n)$ , and  $(c_n)$  by their conditional expectations given  $\mathscr{F}_n$ . It is also convenient to assume  $\mathscr{A} = \sigma(\mathscr{F}_n, n \ge 0)$ .

By properly enlarging the probability space  $(\Omega, \mathscr{A}, \mathbf{P})$ , one can assume w.l.o.g. that it supports a double sequence  $(X_n, Y_n)_{n=0}^{\infty}$  of *iid* variables, uniformly distributed over [0, 1], such that, for each *n*: (i)  $(X_n, Y_n)$  is independent of the process  $(a_k, b_k, c_k)_k$ ; (ii)  $(X_n, Y_n)$  is  $\mathscr{F}_{n+1}$ -measurable, and independent of  $\mathscr{F}_n$ .

Define the stopping game as follows. A *strategy* for player 1 (*resp.* player 2) is a [0, 1]-valued, adapted process  $\mathbf{x} = (x_n)$  (*resp.*  $\mathbf{y} = (y_n)$ ):  $x_n$  is the probability that player 1 stops at stage *n*, conditional on stopping occurs after n - 1. The interpretation of a strategy as a randomized stopping time will be discussed in Section 6.

Given strategies (**x**, **y**), define the stopping stages of players 1 and 2 by  $t_1 = \inf\{n \ge 0, X_n \le x_n\}, t_2 = \inf\{n \ge 0, Y_n \le y_n\}$ , and set

$$t = \min(t_1, t_2). \tag{2}$$

Notice that t + 1 is a stopping time, but t needs not be.

We set  $r(\mathbf{x}, \mathbf{y}) = a_{t_1} \mathbf{1}_{t_1 < t_2} + b_{t_2} \mathbf{1}_{t_2 < t_1} + c_{t_1} \mathbf{1}_{t_1 = t_2 < +\infty}$ . The payoff of the game is  $\gamma(\mathbf{x}, \mathbf{y}) = \mathbf{E}(r(\mathbf{x}, \mathbf{y}))$ . The goal of player 1 is to maximize  $\gamma(\mathbf{x}, \mathbf{y})$ , and the goal of player 2 is to minimize it.

**Definition 2.1.**  $v \in \mathbf{R}$  is the value of the game if  $v = \sup_{\mathbf{x}} \inf_{\mathbf{y}} \gamma(\mathbf{x}, \mathbf{y}) = \inf_{\mathbf{y}} \sup_{\mathbf{x}} \gamma(\mathbf{x}, \mathbf{y})$ . Let  $\epsilon > 0$ . A strategy  $\mathbf{x}$  that satisfies  $\inf_{\mathbf{y}} \gamma(\mathbf{x}, \mathbf{y}) \ge v - \epsilon$  is an  $\epsilon$ -optimal strategy for player 1. A strategy  $\mathbf{y}$  that satisfies  $\sup_{\mathbf{x}} \gamma(\mathbf{x}, \mathbf{y}) \le v + \epsilon$  is an  $\epsilon$ -optimal strategy for player 2.

We will establish the following:

**Theorem 2.2.** Every zero-sum stopping game that satisfies (1) has a value v.

Let  $\lambda \in ]0, 1[$ . Define the  $\lambda$ -discounted payoff by  $r_{\lambda}(\mathbf{x}, \mathbf{y}) = (1 - \lambda)^{t+1} r(\mathbf{x}, \mathbf{y})$ and  $\gamma_{\lambda}(\mathbf{x}, \mathbf{y}) = \mathbf{E}(r_{\lambda}(\mathbf{x}, \mathbf{y})).$ 

**Definition 2.3.**  $v_{\lambda}$  is the  $\lambda$ -discounted value of the game if

 $v_{\lambda} = \sup_{\mathbf{x}} \inf_{\mathbf{y}} \gamma_{\lambda}(\mathbf{x}, \mathbf{y}) = \inf_{\mathbf{y}} \sup_{\mathbf{x}} \gamma_{\lambda}(\mathbf{x}, \mathbf{y}).$ 

Yasuda (1985) proves that the  $\lambda$ -discounted value always exists. In the sequel we prove that

**Theorem 2.4.**  $v = \lim_{\lambda \to 0} v_{\lambda}$ .

In particular,  $\lim_{\lambda \to 0} v_{\lambda}$  exists.

Set  $\gamma_n(\mathbf{x}, \mathbf{y}) = \mathbf{E}(\frac{n-t}{n}r(\mathbf{x}, \mathbf{y})\mathbf{1}_{t < n})$ . The natural interpretation of  $\gamma_n(\mathbf{x}, \mathbf{y})$  is in terms of average payoffs: for  $k \in \mathbf{N}$ , set  $g_k = r(\mathbf{x}, \mathbf{y})$  on  $\{t < k\}$  and  $g_k = 0$  otherwise. Then  $\gamma_n(\mathbf{x}, \mathbf{y}) = \mathbf{E}(\frac{1}{n}\sum_{k=0}^{n-1}g_k)$ .

By dominated convergence,  $\lim_{n} \gamma_n(\mathbf{x}, \mathbf{y}) = \gamma(\mathbf{x}, \mathbf{y})$ . Therefore, if  $\mathbf{x}^*$  is an  $\epsilon$ -optimal strategy of player 1, then for every  $\mathbf{y}$  there exists a stage N such that  $\gamma_n(\mathbf{x}^*, \mathbf{y}) \ge v - 2\epsilon$  holds for every  $n \ge N$ .

We prove that the value v is uniform in the sense below.

**Theorem 2.5.** For every  $\epsilon > 0$ , there exist  $\mathbf{x}^*$  and  $N \in \mathbf{N}$ , such that, for every  $\mathbf{y}$  and every  $n \ge N$ ,  $\gamma_n(\mathbf{x}^*, \mathbf{y}) \ge v - \epsilon$ . A symmetric result holds for player 2.

Thus, Theorem 2.5 is a strengthening of Theorem 2.2. It can be shown that it also implies Theorem 2.4. We then say that v is the *uniform* value of the game.

Theorem 2.5 was proved by Mertens and Neyman (1981) for general stochastic games with bounded payoffs, in which the function  $\lambda \mapsto v_{\lambda}$  satisfies some bounded variation property. In the case of recursive games with bounded payoffs, Rosenberg and Vieille (2000) proved that Theorem 2.5 holds, if  $(v_{\lambda})$  converge uniformly as  $\lambda$  goes to 0 (the uniformity is with respect to the initial state of the game). Our proof does not require any conditions on the discounted values.

## 3. Local games

#### 3.1. Reminder and definitions

Let  $g : A \times B \to \mathbf{R}$ , where *A* and *B* are finite sets (*g* is the payoff function of a zero-sum matrix game with action sets *A* and *B*). Denote by  $\Delta(A)$  and  $\Delta(B)$  the sets of probability distributions over *A* and *B*, and still by *g* the bilinear extension of *g* to  $\Delta(A) \times \Delta(B)$ .

The min max theorem states that  $\sup_{x \in \Delta(A)} \inf_{y \in \Delta(B)} g(x, y) = \inf_{y \in \Delta(B)} \sup_{x \in \Delta(A)} g(x, y)$ , which we denote by val g. Any x (*resp.* y) which achieves the sup on the left side (*resp.* inf on the right side) is called an optimal strategy of player 1 (*resp.* player 2). It is well known that the operator val is non-decreasing and non-expansive: val  $f \leq val g$  if  $f \leq g$ , and  $|val f - val g| \leq \sup_{A \times B} |f - g|$ .

For any real-valued  $\mathscr{F}_n$ -measurable function f, we let  $G_n(f)$  be the 0-sum game with ( $\mathscr{F}_n$ -measurable) payoff matrix

f	$b_n$
$a_n$	$c_n$

in which player 1 chooses a row and player 2 a column.

A strategy of player 1 in this game is a [0, 1]-valued,  $\mathcal{F}_n$ -measurable variable  $x_n$ , to be interpreted as the probability that player 1 chooses the bottom row. A strategy of player 2 is defined analogously.

Define  $G_n(x_n, y_n; f)$  to be the ( $\mathscr{F}_n$ -measurable) payoff to player 1 when the players use strategies  $x_n$  and  $y_n$ :

$$G_n(x_n, y_n; f) = x_n(1 - y_n)a_n + y_n(1 - x_n)b_n + x_ny_nc_n + (1 - x_n)(1 - y_n)f.$$

By the min max theorem, for every  $\omega \in \Omega$  the game with payoff matrix

$f(\omega)$	$b_n(\omega)$
$a_n(\omega)$	$c_n(\omega)$

has a value, denoted by val  $G_n(f)(\omega)$ , for every  $\omega \in \Omega$ .

We now argue that each player has an optimal strategy in  $G_n(f)$ .

**Proposition 3.1.** Let f be  $\mathcal{F}_n$ -measurable and real-valued. There exists a strategy  $x_n$  in  $G_n(f)$ , such that, for every y,

$$G_n(x_n, y; f) \ge \text{val } G_n(f) \text{ everywhere.}$$

A symmetric property holds for player 2.

*Proof.* For every  $\omega \in \Omega$ , the game with payoff matrix

$$\frac{f(\omega)}{a_n(\omega)} \frac{b_n(\omega)}{c_n(\omega)}$$

has optimal strategies for both players. Since  $f, a_n, b_n$  and  $c_n$  are all  $\mathscr{F}_n$ -measurable, the map which associates to each  $\omega$  the set of optimal strategies for player 1 is upper-semi-continuous and  $\mathscr{F}_n$ -measurable. By Kuratowski and Ryll-Nardzewski (1965) it has an  $\mathscr{F}_n$ -measurable selection.

Any  $x_n$  that satisfies the conclusion of Proposition 3.1 is said to be optimal in the game  $G_n(f)$ . If  $x_n$  and  $y_n$  are optimal strategies in  $G_n(f)$ , one has  $G_n(x_n, y_n; f) =$  val  $G_n(f)$  everywhere. In particular, val  $G_n(f)$  is  $\mathscr{F}_n$ -measurable.

#### 3.2. Local games and discounted values

It is useful to extend the notions of discounted values to the game *starting at stage n*.

For  $n \in \mathbf{N}$ , set  $\Sigma_n = {\mathbf{x}, x_p = 0, \forall p < n}$ , and  $T_n = {\mathbf{y}, y_p = 0, \forall p < n}$ . Those are strategies where the probability that the players stop before stage *n* is zero. Set

$$\underline{v}_n(\lambda) = \operatorname{esssup}_{\Sigma_n} \operatorname{essinf}_{T_n} \mathbf{E}[(1-\lambda)^{-n} r_\lambda(\mathbf{x}, \mathbf{y}) | \mathscr{F}_n],$$

and

$$\overline{v}_n(\lambda) = \operatorname{essinf}_{T_n} \operatorname{esssup}_{\Sigma_n} \mathbf{E}[(1-\lambda)^{-n} r_\lambda(\mathbf{x}, \mathbf{y}) | \mathscr{F}_n].$$

The proposition below contains obvious properties.

**Proposition 3.2.**  $(\overline{v}_n(\lambda))_n$  and  $(\underline{v}_n(\lambda))_n$  are adapted processes. Moreover,  $\sup_n |\overline{v}_n(\lambda)|, \sup_n |\underline{v}_n(\lambda)| \in L^1(\mathbf{P}).$ 

Yasuda (1985) proves that  $(\overline{v}_n(\lambda))_n$  and  $(\underline{v}_n(\lambda))_n$  are both solutions of the recursive equation

$$v_n(\lambda) = (1 - \lambda) \text{val } G_n(\mathbf{E}[v_{n+1}(\lambda)|\mathscr{F}_n]), \ \mathbf{P} - a.s.$$
(3)

He then proves that any solution of this sequence of equations is at most  $(\underline{v}_n(\lambda))$ and at least  $(\overline{v}_n(\lambda))$ . Since  $\overline{v}_n(\lambda) \ge \underline{v}_n(\lambda)$  it follows that the two are equal, **P**-a.s. We give a shorter argument, adapted from Shapley (1953). Since the value operator is non-expansive,

$$\begin{aligned} |\overline{v}_{n}(\lambda) - \underline{v}_{n}(\lambda)| &\leq (1-\lambda) |\mathbf{E}[\overline{v}_{n+1}(\lambda) - \underline{v}_{n+1}(\lambda)|\mathscr{F}_{n}]| \\ &\leq (1-\lambda) \mathbf{E}[|\overline{v}_{n+1}(\lambda) - v_{n+1}(\lambda)||\mathscr{F}_{n}] \end{aligned}$$

By taking expectations, one obtains

$$\begin{aligned} \|\overline{v}_n(\lambda) - \underline{v}_n(\lambda)\|_1 &\leq (1-\lambda) \|\overline{v}_{n+1}(\lambda) - \underline{v}_{n+1}(\lambda)\|_1 \\ &\leq (1-\lambda)^p \|\overline{v}_{n+p}(\lambda) - \underline{v}_{n+p}(\lambda)\|_1 \end{aligned}$$

for each  $p \in \mathbf{N}$ . Since  $\sup_n |\overline{v}_n(\lambda)|$ ,  $\sup_n |\underline{v}_n(\lambda)| \in L^1(\mathbf{P})$ , one obtains by letting  $p \to \infty$  that  $\overline{v}_n(\lambda) = \underline{v}_n(\lambda)$ , **P**-a.s. We define  $v_n(\lambda) = \underline{v}_n(\lambda)$  (=  $\overline{v}_n(\lambda)$ ) to be the  $\lambda$ -discounted value of the game starting at stage *n*. Notice that  $v(\lambda) = \mathbf{E}[v_0(\lambda)]$ .

We now let  $(\bar{\lambda}_p)_p$  be any decreasing sequence which converges to 0. Set  $v_n = \limsup_{p \to \infty} v_n(\bar{\lambda}_p)$ , and  $w = \mathbf{E}[v_0]$ . We shall prove the next proposition.

**Proposition 3.3.** For every  $\varepsilon > 0$ , there is a strategy  $\overline{\mathbf{x}}$  of player 1, and a positive integer N such that

$$\forall \mathbf{y}, \forall n \geq N, \ \gamma_n(\mathbf{\overline{x}}, \mathbf{y}) \geq w - \varepsilon.$$

We now explain why Proposition 3.3 implies Theorem 2.5 – w is the value of the game. Define  $z_0 = \liminf_{p\to\infty} v_0(\bar{\lambda}_p)$ , and  $z = \mathbf{E}[z_0]$ . By symmetry, for each  $\epsilon$ , there exists a strategy  $\bar{\mathbf{y}}$  such that  $\gamma_n(\mathbf{x}, \bar{\mathbf{y}}) \leq z + \epsilon$  for each  $\mathbf{x}$ , provided n is large enough. This readily implies  $w - \epsilon \leq z + \epsilon$ . Since  $z \leq w$ , and  $\epsilon$  is arbitrary, one obtains w = z. This shows that w is the uniform value of the game. The claim about the limit of discounted values is now immediate, since the sequence  $(\bar{\lambda}_p)$  used to define w is arbitrary.

The following result will be used later.

**Proposition 3.4.** One has  $v_n \leq \text{val } G_n(\mathbb{E}[v_{n+1}|\mathcal{F}_n])$ , for every *n*.

*Proof.* Recall that  $v_n(\lambda) = (1 - \lambda)$ val  $G_n(\mathbf{E}[v_{n+1}(\lambda)|\mathscr{F}_n])$ . By monotonicity of the value operator,

$$v_n(\bar{\lambda}_q) \le \alpha_q \text{val } G_n(\mathbf{E}[\sup_{p \ge q} v_{n+1}(\bar{\lambda}_p) | \mathscr{F}_n]), \text{ for each } q, \tag{4}$$

where  $\alpha_q = 1 - \bar{\lambda}_q$  if the val is negative, and 1 otherwise. By dominated convergence,  $\lim_{q \to +\infty} \mathbf{E}[\sup_{p \ge q} v_{n+1}(\bar{\lambda}_p) | \mathscr{F}_n] = \mathbf{E}[v_{n+1} | \mathscr{F}_n]$ . Since the val operator is non-expansive, the right-hand side of (4) converges to val  $G_n(\mathbf{E}[v_{n+1} | \mathscr{F}_n])$ , **P**-a.s. The result follows.

### 3.3. Locally optimal strategies and martingale properties

Denote by  $x_n(\lambda)$  and by  $x_n^*$  optimal strategies of player 1 in the local games  $G_n(\mathbf{E}[v_{n+1}(\lambda)|\mathscr{F}_n])$  and  $G_n(\mathbf{E}[v_{n+1}|\mathscr{F}_n])$ , which exist by Proposition 3.1.

Thus, for every strategy **y** and every  $n \ge 0$ , one has

$$G_n(x_n^*, y_n; \mathbf{E}[v_{n+1}|\mathscr{F}_n]) \ge v_n, \ \mathbf{P}\text{-a.s.}$$
(5)

and

$$(1 - \lambda)G_n(x_n(\lambda), y_n; \mathbf{E}[v_{n+1}(\lambda)|\mathscr{F}_n]) \ge v_n(\lambda) \mathbf{P}\text{-a.s.}$$
(6)

Recall that  $v_n(\lambda)$  is to be interpreted as the value of the (discounted) game starting in stage *n*, *conditional* on the fact that the game has not been stopped. Define the strategies  $\mathbf{x}(\lambda) = (x_n(\lambda))_n$  and  $\mathbf{x}^* = (x_n^*)_n$ .

Equation 3 and Proposition 3.4 provide recursive formulas for  $(v_n)_n$  and  $(v_n(\lambda))_n$ . In order to interpret these formulas in terms of submartingale properties, we use auxiliary processes.

For clarity of exposition, given any two events *E* and *A* in  $\mathscr{A}$ , we say that *E* holds **P**-a.s. on *A* if  $\mathbf{P}(A \cap E^c) = 0$ . We will frequently omit the qualification **P**-a.s.

Let  $(\alpha_n)_n$  be an adapted integrable process on  $(\Omega, \mathscr{A}, (\mathscr{F}_n), \mathbf{P})$ , and  $s_1 \leq s_2$ two stopping times (with values in  $\mathbf{N} \cup \{+\infty\}$ ). We say that  $(\alpha_n)_n$  is a submartingale between  $s_1$  and  $s_2$  if, for every  $n \geq 0$ , the inequality  $\mathbf{E}[\alpha_{n+1}|\mathscr{F}_n] \geq \alpha_n$  holds **P**-a.s. on the event  $\{s_1 \leq n < s_2\}$ . The process  $(\alpha_n)_n$  is a submartingale up to  $s_2$  if it is a submartingale between 0 and  $s_2$ . It is straightforward to adapt the sampling theorem as follows. Let  $(\alpha_n)$  be a submartingale between  $s_1$  and  $s_2$ . Let s be a stopping time, with **P**-a.s. finite values, such that  $s \leq s_2$ . Denote by  $\mathscr{F}_{s_1}$  the  $\sigma$ -algebra of events known at stage  $s_1$ . Then one has  $\mathbf{E}[\alpha_s|\mathscr{F}_{s_1}] \geq \alpha_{s_1}$ , **P**-a.s. on the event  $\{s_1 \leq s\}$ .

Let  $(\mathbf{x}, \mathbf{y})$  be a pair of strategies and t the induced stopping stage defined by (2). We define  $(\tilde{\alpha}_n)$  as  $\tilde{\alpha}_n = \alpha_n$  on  $\{t \ge n\}$  and  $\tilde{\alpha}_n = r(\mathbf{x}, \mathbf{y})$  if t < n. The process  $(\tilde{\alpha}_n)$  depends on  $(\mathbf{x}, \mathbf{y})$ . To avoid ambiguity, we will sometimes write: under  $(\mathbf{x}, \mathbf{y})$ , the process  $(\tilde{\alpha}_n)$  *etc*, when we wish to emphasize which strategies are being used in the definition of  $(\tilde{\alpha}_n)$ . With a (convenient) abuse of terminology, we refer to  $(\tilde{\alpha}_n)$  as the process  $(\alpha_n)$  stopped at t.

We use repeatedly the following relation, which holds **P**-a.s. on the event  $\{t \ge n\}$ :

$$\mathbf{E}[\tilde{\alpha}_{n+1}|\mathscr{F}_n] = G_n\left(x_n, y_n; \mathbf{E}\left[\alpha_{n+1}|\mathscr{F}_n\right]\right)$$
(7)

if  $(X_n, Y_n)$  is independent of  $\alpha_{n+1}$ . This latter independence property holds in all cases of interest, for instance if  $\alpha_{n+1} = v_{n+1}$  or  $\alpha_{n+1} = v_{n+1}(\lambda)$ , so that we shall apply (7) without further justification.

Set  $\mathscr{F}_n^2 = \sigma(\mathscr{F}_n, Y_n)$ , so that  $\mathscr{F}_n^2$  includes past and present values of the payoff processes, past "decisions" of the players and the decision of player 2 at stage *n*.

**Lemma 3.5.** Let  $\mathbf{y}$  be a strategy of player 2, and  $\lambda \in ]0, 1[$ . Under  $(\mathbf{x}(\lambda), \mathbf{y}), ((1 - \lambda)^n \tilde{v}_n(\lambda))_n$  is a submartingale up to t + 1. Under  $(\mathbf{x}^*, \mathbf{y}), (\tilde{v}_n)_n$  is a submartingale, both for  $(\mathcal{F}_n)$  and  $(\mathcal{F}_n^2)_n$ .

Notice that  $\sup_{n} |\tilde{v}_{n}(\lambda)|$  and  $\sup_{n} |\tilde{v}_{n}|$  belong to  $L^{1}(\mathbf{P})$ , for every choice of  $(\mathbf{x}, \mathbf{y})$ .

*Proof.* Let  $n \ge 0$ . On the event  $\{t \ge n\}$ ,

$$\mathbf{E}[(1-\lambda)\tilde{v}_{n+1}(\lambda)|\mathscr{F}_n] = (1-\lambda)G_n(x_n(\lambda), y_n; \mathbf{E}[v_{n+1}(\lambda)|\mathscr{F}_n]),$$

which is at least  $v_n(\lambda)$ , by (6). This proves the first claim since  $\tilde{v}_n(\lambda) = v_n(\lambda)$  if  $t \ge n$ .

For a similar reason, using (5),

$$\mathbf{E}[\tilde{v}_{n+1}|\mathscr{F}_n] \geq \tilde{v}_n,$$

on the event  $\{t \ge n\}$ . On  $\{t < n\}$ ,  $\tilde{v}_{n+1} = \tilde{v}_n$ . The same computation works also for the filtration  $(\mathscr{F}_n^2)_n$ . This completes the proof.

**Corollary 3.6.** For every  $\mathbf{y}$ ,  $\gamma_{\lambda}(\mathbf{x}(\lambda), \mathbf{y}) \geq \mathbf{E}(v_0(\lambda))$ .

*Proof.* Fix a strategy **y** of player 2. Let  $n \ge 0$ , and apply the submartingale property with the stopping time min(t + 1, n):

$$\mathbf{E}[(1-\lambda)^{\min(t+1,n)}\tilde{v}_{\min(t+1,n)}] \ge \mathbf{E}(v_0(\lambda)),$$

that is, using the definition of the stopped process  $(\tilde{v}_n)_n$ :

$$\mathbf{E}[(1-\lambda)^n v_n(\lambda)\mathbf{1}_{t>n} + (1-\lambda)^{t+1} r(\mathbf{x}(\lambda), \mathbf{y})\mathbf{1}_{t< n}] \ge \mathbf{E}(v_0(\lambda)).$$

By dominated convergence, the left-hand side converges to  $\gamma_{\lambda}(\mathbf{x}(\lambda), \mathbf{y})$ .

A similar proof proves the following.

**Corollary 3.7.** Let  $n \in \mathbb{N}$ . Let  $\tilde{\mathbf{x}}(\lambda)$  be the strategy that is identically 0 until stage n, and coincides with  $\mathbf{x}(\lambda)$  afterwards. Let  $\mathbf{y}$  be any strategy of player 2 that is identically 0 until stage n. Then

$$\mathbf{E}[(1-\lambda)^{t+1-n}r(\mathbf{x}(\lambda),\mathbf{y})|\mathscr{F}_n] \ge v_n(\lambda).$$

Corollary 3.6 implies that in the discounted game it is an optimal strategy for player 1 to play  $\mathbf{x}(\lambda)$ . No such result holds for the original problem: playing  $\mathbf{x}^*$  needs not be an optimal strategy.

Example

This matrix notation is a shortcut for the stopping game with payoffs  $a_n = b_n = 1$ ,  $c_n = 0$ , **P**-a.s. for every *n*. Clearly  $v_n$  and  $v_n(\lambda)$  are independent of *n* and constant, so we simply write *v* and  $v(\lambda)$ . The real number  $0 \le v(\lambda) \le 1$  is a solution to the equation  $v(\lambda) = (1 - \lambda) \operatorname{val} G(v(\lambda))$ , from which it is easily derived

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 $v(\lambda) = 1 - \sqrt{\lambda}$ , and  $\mathbf{x}(\lambda) = \sqrt{\lambda}/(1 + \sqrt{\lambda})$ . Therefore v = 1. Denote by **0** the strategy (of either player 1 or player 2) that never stops  $(0_n = 0 \text{ for all } n)$ . Then  $\mathbf{x}^* = \mathbf{0}$ . However,  $\gamma(\mathbf{x}^*, \mathbf{0}) = 0$ .

Nevertheless, if  $t_1$  is **P**-a.s. finite under  $\mathbf{x}^*$ , then  $\mathbf{x}^*$  is optimal for player 1.

**Lemma 3.8.** If  $\mathbf{P}(t_1 < +\infty) = 1$  under  $\mathbf{x}^*$ , then  $\mathbf{x}^*$  guarantees w for player 1.

*Proof.* Let **y** be an arbitrary strategy of player 2. By Lemma 3.5,  $(\tilde{v}_n)$  is a submartingale under  $(\mathbf{x}^*, \mathbf{y})$ . Since  $\mathbf{P}(t_1 < +\infty) = 1$ ,  $\mathbf{P}(t < +\infty) = 1$  as well, hence it follows that

$$\mathbf{E}[r(\mathbf{x}^*, \mathbf{y})]_{t < +\infty}] = \mathbf{E}[\tilde{v}_{\infty}] \ge \mathbf{E}[v_0] = w,$$

as desired.

#### 4. The main ideas of the proofs

We give a detailed sketch of the proofs in the deterministic case. Many technical issues disappear in that case. Therefore the main ideas appear, hopefully more clearly. Assume that  $(a_n)_n$ ,  $(b_n)_n$ ,  $(c_n)_n$ , and therefore also  $(v_n)_n$  and  $(v_n(\lambda))_n$ , are sequences of real numbers, bounded by 1.

For every **y**,  $(\tilde{v}_n)_n$  is a bounded submartingale under  $(\mathbf{x}^*, \mathbf{y})$ , thus

$$\mathbf{E}\left[\tilde{v}_{\infty}\right] \ge \mathbf{E}[v_0] = w \tag{8}$$

with  $\tilde{v}_{\infty} = \lim_{n \to \infty} \tilde{v}_{n}$ .

For  $\mathbf{y} = \mathbf{0}$ ,  $(\tilde{v}_n)$  coincides with  $v_n$  up to  $t_1$ . Thus,  $t_1 < +\infty$ , **P**-a.s., or  $(v_n)_n$  is a convergent sequence. In the first case,  $\mathbf{x}^*$  is optimal by (8).

We now assume that  $(v_n)$  is a convergent sequence, and given  $\varepsilon > 0$ , we choose  $N_0$  such that  $\sup_{n,m \ge N_0} |v_n - v_m| \le \varepsilon/2$ . We also assume for simplicity  $N_0 = 0$  (in the general case, the strategies below would be supplemented by: play  $\mathbf{x}^*$  up to  $N_0$ ). If  $w \le \varepsilon$ ,  $\tilde{v}_{\infty} \le 3\varepsilon/2$ , so that  $\mathbf{x}^*$  is  $3\varepsilon/2$ -optimal by (8). We are thus led to consider the case  $w > \varepsilon$ .

*First proof.* Choose  $\lambda_0$  such that  $v(\lambda_0) \ge w - \varepsilon/3$  and  $\varepsilon' \in (0, \varepsilon/6)$ . Player 1 starts playing according to  $\mathbf{x}(\lambda_0)$ . For each  $\mathbf{y}$ ,  $((1 - \lambda_0)^n \tilde{v}_n(\lambda_0))_n$  is a submartingale up to *t*. Set  $s_1 = \inf\{n, v_n(\lambda_0) \le \varepsilon'\}$ . Since  $(\tilde{v}_n(\lambda_0))_n$  is bounded,  $\min(t, s_1)$  is **P**-a.s. finite. Moreover, since  $v_{s_1}(\lambda_0) \le v(\lambda_0) - (\varepsilon/6 - \varepsilon')$  if  $s_1 \le t$ , the probability that  $t < s_1$  is bounded away from 0.

At stage  $s_1$ , the approximation of  $(v_n)$  by  $(v_n(\lambda_0))_n$  gets poor, so we switch to a new discount factor:  $\lambda_0$  is replaced by  $\lambda_1$ , with  $v_{s_1}(\lambda_1) \ge v_{s_1} - \varepsilon/3 \ge \varepsilon/6$ , and  $\mathbf{x}(\lambda_1)$  is played until  $s_2 = \inf\{n > s_1, v_n(\lambda_1) \le \varepsilon'\}$ , where we again switch from  $\lambda_1$  to  $\lambda_2$ , and so on.

Call  $\overline{\mathbf{x}}$  the resulting strategy. Under  $(\overline{\mathbf{x}}, \mathbf{y})$ , *t* is **P**-a.s. finite, since for every *n*, the probability of stopping between  $s_n$  and  $s_{n+1}$  is bounded away from 0. Introduce

the sequence  $(w_n)_n$ , where  $w_n = v_n(\lambda_p)$  if  $s_p \le n < s_{p+1}$ . By construction,  $w_0 \ge v - \varepsilon/3$  and  $(\tilde{w}_n)_n$  is a submartingale. Since  $t < +\infty$ , it converges to  $r(\bar{\mathbf{x}}, \mathbf{y})_{t < +\infty}$ , therefore  $\gamma(\bar{\mathbf{x}}, \mathbf{y}) \ge w - \varepsilon/3$ .

Second proof. The definition of  $\overline{\mathbf{x}}$  here is motivated by the observation

$$\limsup_{n} a_n \ge w - \varepsilon \tag{9}$$

which is derived as follows. For each  $\lambda$ , under  $(\mathbf{x}(\lambda), \mathbf{0})$ ,  $t = t_1$  and  $r(\mathbf{x}(\lambda), \mathbf{0}) = a_t$  if  $t < +\infty$ ; thus,

$$\mathbf{E}\left[(1-\lambda)^{t+1}a_t\mathbf{1}_{t<+\infty}\right] = \gamma_{\lambda}(\mathbf{x}(\lambda), \mathbf{0}) \ge v(\lambda).$$

The left-hand side lies in the closed convex hull of  $\{0, a_n, n \in \mathbf{N}\}$ . Given any  $\delta > 0$ ,  $v(\lambda) \ge w - \delta$ , for a suitable  $\lambda$ . Therefore,  $\sup_n a_n \ge w - \delta$ . Since  $v_n \ge w - \varepsilon$  for every *n*, this proof may be repeated, and (9) holds.

We define  $\overline{\mathbf{x}}$  by  $\overline{x}_n = x_n^* + \varepsilon$  if  $a_n \ge w - 2\varepsilon$ , and  $\overline{x}_n = x_n^*$  otherwise. Since (9) holds,  $t_1 < +\infty$  **P**-a.s. under  $\overline{\mathbf{x}}$ . To see that this strategy guarantees player 1 an expected payoff of w, we note that the following points hold:

- 1. If player 2 stops the game  $(t = t_2)$ , then the expected payoff of player 1 is at least w (up to an  $\epsilon$ ).
- 2. In the case that player 2 always continues, since player 1 changes his strategy *only* when a unilateral stopping is favorable for him,  $\mathbf{E}[v_n] \ge w \epsilon$ .

#### 5. Two $\epsilon$ -optimal strategies

## 5.1. Preliminaries

For the rest of the section we fix  $\epsilon > 0$ . Set  $m = \sup_n (\sup(|a_n|, |b_n|, |c_n|))$ . Since  $m \in L^1(\mathbf{P})$ , there exists  $\eta > 0$  such that, for every  $A \in \mathcal{A}$ ,

$$\mathbf{P}(A) < \eta \Rightarrow \mathbf{E}(m1_A) < \epsilon. \tag{10}$$

Notice that  $|v_n(\lambda)|, |v_n| \leq \mathbf{E}[m \mid \mathscr{F}_n], \mathbf{P}$ -a.s. for every *n*.

The sequence  $(v_n)$  needs not converge. On the other hand, the process  $(\tilde{v}_n)$ , being a submartingale under  $(\mathbf{x}^*, \mathbf{y})$  (with sup  $\tilde{v}_n \in L^1(\mathbf{P})$ ) converges **P**-a.s. and in  $L^1(\mathbf{P})$ , for every **y**.

The stopping time  $t_1$  is a function of player 1's strategy. Under  $(\mathbf{x}^*, \mathbf{0}), t = t_1$ , **P**-a.s. This implies that  $(v_n)$  converges **P**-a.s. on the set  $\{t_1 = +\infty\}$ .

Choose  $N_0 \in \mathbf{N}$  such that

$$\mathbf{P}\{\sup_{n,m\geq N_0} |v_n - v_m| > \epsilon/2, t_1 \geq N_0\} < \eta.$$
(11)

Thus, after stage  $N_0$ , with high probability  $v_n$  does not change by much.

#### 5.2. An $\epsilon$ -optimal strategy for player 1 - I

We first define the switching stages  $(s_p)$  and the approximating discount factors  $(\lambda_p)$ :  $v(\lambda_p)$  approximates v between  $s_p$  and  $s_{p+1}$ . Set  $s_0 = N_0$  if  $v_{N_0} > \epsilon$ , and  $s_0 = +\infty$  otherwise. Choose  $\varepsilon' \in (0, \varepsilon/6)$  and an  $\mathscr{F}_{s_0}$ -measurable function  $\lambda_0$  with  $v_{s_0}(\lambda_0) > v_{s_0} - \epsilon/3$  if  $s_0 < +\infty$ .

Set  $s_{p+1} = \inf\{n > s_p, v_n(\lambda_p) \le \varepsilon'\}$  and choose an  $\mathscr{F}_{s_{p+1}}$ -measurable function  $\lambda_{p+1}$ , such that  $v_{s_{p+1}}(\lambda_{p+1}) > v_{s_{p+1}} - \epsilon/3$  if  $s_{p+1} < +\infty$ .

Let  $\overline{\mathbf{x}}$  be the strategy that coincides with  $\mathbf{x}^*$  until  $s_0$ , and with  $\mathbf{x}(\lambda_p)$  between  $s_p$  and  $s_{p+1}$ :

$$\overline{x}_n = \begin{cases} x_n^* & n < s_0\\ x_n(\lambda_p) & s_p \le n < s_{p+1} \end{cases}$$

We shall prove that  $\overline{\mathbf{x}}$  is  $7\varepsilon$ -optimal.

By Lemma 3.5, for every **y**,  $(\tilde{v}_n)_n$  is a submartingale up to  $s_0$ , and  $((1 - \lambda_p)^n \tilde{v}_n(\lambda_p))_n$  is a submartingale between  $\min(s_p, t+1)$  and  $\min(s_{p+1}, t+1)$ , for each p.

We introduce an auxiliary variable  $z_n$  defined as

$$z_n = \begin{cases} v_n - \epsilon/3 & n < s_0\\ v_n(\lambda_p) & s_p \le n < s_{p+1} \end{cases}$$

Intuitively,  $z_n$  is (up to  $\epsilon/3$ ), the parameter we are interested in: the limit  $v_n$  before stage  $s_0$ , and the  $\lambda_p$ -discounted value for  $s_p \le n < s_{p+1}$ .

We ultimately wish to get a submartingale. A minor adjustment is needed. Define the stopping time *s* by  $s = +\infty$  if  $s_0 = +\infty$  and  $s = \inf\{n \ge N_0, v_n \le \varepsilon/2\}$ otherwise. By the definition of  $N_0$ ,  $\mathbf{P}(s < +\infty, t_1 \ge N_0) < \eta$ . We use *s* to define a process  $(w_n)$  by

$$w_n = \begin{cases} \mathbf{E} \left[ m | \mathscr{F}_n \right] & s \le n \\ z_n & \text{otherwise} \end{cases}$$

Observe that

$$\tilde{w}_{n+1} \ge \tilde{v}_{n+1}(\lambda_p) \text{ on the event } \{s_p \le n < s_{p+1}\}.$$
(12)

Indeed, this is clear if  $s \le n + 1$  or if t < n + 1. If not :

$$\tilde{w}_{n+1} = \begin{cases} v_{n+1}(\lambda_p) & n+1 < s_{p+1} \\ v_{n+1}(\lambda_{p+1}) & n+1 = s_{p+1} \end{cases}$$

If  $n + 1 < s_{p+1}$ , then  $\tilde{w}_{n+1} = \tilde{v}_{n+1}(\lambda_p)$ , while if  $n + 1 = s_{p+1}$ ,

$$\tilde{w}_{n+1} \ge v_{n+1} - \epsilon/3 \ge \varepsilon' \ge v_{n+1}(\lambda_p).$$

We set  $\overline{t} + 1 = \min(t + 1, s)$ . Observe that  $\mathbf{P}(t = \overline{t}) \ge 1 - \eta$ .

**Lemma 5.1.** For every y,  $(\tilde{w}_n)$  is a submartingale up to  $\bar{t} + 1$  under  $(\bar{x}, y)$ .

*Proof.* Fix a strategy **y** of player 2. Let  $n \in \mathbf{N}$ . We prove that  $\mathbf{E}[\tilde{w}_{n+1}|\mathcal{F}_n] \ge \tilde{w}_n$ , **P**-a.s. on the event  $\{\bar{t} + 1 > n\}$ .

If  $n < s_0$ ,  $w_n = v_n - \epsilon/3$ ,  $w_{n+1} \ge v_{n+1} - \epsilon/3$  (with equality if  $n+1 < s_0$ ), and  $\overline{x}_n = x_n^*$ . Thus  $\mathbf{E}[\tilde{w}_{n+1}|\mathscr{F}_n] \ge G_n(x_n^*, y_n; \mathbf{E}[v_{n+1} - \epsilon/3|\mathscr{F}_n]) \ge v_n - \epsilon/3$ , where the second inequality follows from the inequality  $G_n(x_n^*, y_n; \mathbf{E}[v_{n+1}|\mathscr{F}_n]) \ge v_n$  and since the val operator is non-expansive.

If  $s_p \le n < s_{p+1}$ ,  $w_n = v_n(\lambda_p)$ , and  $\overline{x}_n = x_n(\lambda_p)$ . In that case, by (12),

$$\mathbf{E}(\tilde{w}_{n+1}|\mathscr{F}_n) \ge G_n(\overline{x}_n, y_n; \mathbf{E}[v_{n+1}(\lambda_p)|\mathscr{F}_n]) \ge \frac{1}{1-\lambda_p}v_n(\lambda_p) \ge v_n(\lambda_p) = w_n,$$

where the last inequality holds since  $v_n(\lambda_p) > 0$ .

**Lemma 5.2.** For every **y**, under  $(\bar{\mathbf{x}}, \mathbf{y}), \bar{t} < +\infty, \mathbf{P}$ -a.s. on the event  $s_0 = N_0$ .

*Proof.* Fix a strategy **y** of player 2. We proceed in two steps. We prove first that  $\min(s_{p+1}, t) < +\infty$ , **P**-a.s. on  $\{s_p < s\}$ . From  $\min(s_p, t+1)$  up to  $\min(s_{p+1}, t+1)$ ,  $((1 - \lambda_p)^n \tilde{w}_n)$  is a submartingale. Thus, for every  $N \in \mathbf{N}$  and  $n \leq N$ , the sampling property applied to the finite stopping time  $\min(s_{p+1}, t+1, N)$  yields

$$w_{n} \leq \frac{1}{(1-\lambda_{p})^{n}} \mathbf{E} \left[ m(1-\lambda_{p})^{\min(s_{p+1},t+1)} \mathbf{1}_{\min(s_{p+1},t+1) \leq N} + \tilde{w}_{N}(1-\lambda_{p})^{N} \mathbf{1}_{\min(s_{p+1},t+1) > N} | \mathscr{F}_{n} \right]$$

on  $\{s_p \le n < \min(s_{p+1}, t+1)\}$ .

By taking  $N \to +\infty$  and by dominated convergence for conditional expectations, one obtains

$$\varepsilon' < v_n(\lambda_p) = w_n \le \mathbf{E} \left[ m(1 - \lambda_p)^{\min(s_{p+1}, t+1) - n} \mathbb{1}_{\min(s_{p+1}, t+1) < +\infty} | \mathscr{F}_n \right]$$
(13)

on the event  $\{s_p \le n < \min(s_{p+1}, t+1)\}$ .

By taking the limit  $n \to \infty$  in (13), one gets  $\limsup w_n \le 0$ , **P**-a.s. on the event  $\{s_p < +\infty, t = s_{p+1} = +\infty\} \cap \{s_p < s\}$ . But on this event  $w_n \ge \varepsilon'$ , **P**-a.s. for every *n*. This ends the first step.

One can rephrase the conclusion of the first step as  $\min(s_{p+1}, \bar{t}) < +\infty$  if  $\min(s_p, \bar{t}) < +\infty$ , **P**-a.s. By induction,  $\min(s_p, \bar{t}) < +\infty$  if  $s_0 < +\infty$ , **P**-a.s. for every *p*.

Since  $(\tilde{v}_n(\lambda_p))_n$  is a submartingale between  $\min(s_p, \bar{t}+1)$  and  $\min(s_{p+1}, \bar{t}+1)$ , and since  $v_{s_{p+1}}(\lambda_p) \le \varepsilon'$ ,

$$v_{s_p}(\lambda_p) \leq \mathbf{E}[m\mathbf{1}_{\bar{t}+1 \leq s_{p+1}} + \varepsilon' \cdot \mathbf{1}_{s_{p+1} < \bar{t}+1} | \mathscr{F}_{s_p}]$$

on  $\{s_p < \overline{t} + 1\}$ . Since  $v_{s_p}(\lambda_p) \ge \epsilon/6$ , it follows by taking expectations that

$$\frac{\epsilon}{6} \mathbf{P}(s_p < \bar{t} + 1) \le \mathbf{E}(m \mathbf{1}_{s_p < \bar{t} + 1 < +\infty}) + \varepsilon' \mathbf{P}(s_{p+1} < \bar{t} + 1),$$

hence

$$\left(\frac{\epsilon}{6} - \epsilon'\right) \mathbf{P}(s_p < \bar{t} + 1) \le \mathbf{E}(m \mathbf{1}_{s_p < \bar{t} + 1 < +\infty})$$

As *p* goes to infinity, the left-hand side converges to  $(\epsilon/6 - \epsilon')\mathbf{P}(s_0 = N_0, \bar{t} = +\infty)$ , while the right-hand side converges to 0. The result follows.

**Proposition 5.3.** There exists  $N \in \mathbf{N}$  such that, for every  $\mathbf{y}$  and  $n \ge N$ , one has  $\gamma_n(\overline{\mathbf{x}}, \mathbf{y}) \ge w - 7\epsilon$ .

*Proof.* By Lemma 5.2 with  $\mathbf{y} = \mathbf{0}$ , there exists some positive integer  $N_1 \ge N_0$  such that under  $(\bar{\mathbf{x}}, \mathbf{0})$ 

$$\mathbf{P}(s_0 = N_0, t \ge N_1) < \eta.$$
(14)

This readily implies that (14) holds under  $(\bar{\mathbf{x}}, \mathbf{y})$ , for every  $\mathbf{y}$ .

Let now  $N_2$  be sufficiently large such that

$$\frac{N_1}{N_2} \mathbf{E}[m] < \epsilon. \tag{15}$$

Using (14), (10) and (15) we have:

$$\left| \mathbf{E} \left[ \frac{1}{n+1} \sum_{k=0}^{n} g_k \right] - \mathbf{E} \left[ \frac{1}{n-N_1+1} \sum_{k=N_1}^{n} g_k \right] \right| < 2\varepsilon, \text{ provided } n \ge N_2.$$
(16)

\_ .

Fix  $n \ge N_2$  and any strategy **y**.

By definition,  $\gamma_{n+1}(\overline{\mathbf{x}}, \mathbf{y}) = \mathbf{E}[\frac{1}{n+1}\sum_{k=0}^{n} g_k]$ . We will evaluate  $\mathbf{E}\left[\frac{1}{n-N_1+1}\sum_{k=N_1}^{n} g_k\right]$ . Let  $N_1 \le k \le n$ .

On  $\{t < N_1\}, g_k = \tilde{w}_{N_1}.$ 

On  $\{t \ge N_1, s_0 = N_0\}, |g_k| \le m$ , but this event has a probability at most  $2\eta$ .

Consider now the event  $\{t \ge N_1, s_0 = +\infty\}$ . The event  $\{t \ge N_1, s_0 = +\infty, \sup_{q\ge N_1} v_q > 3\varepsilon/2\}$  has probability at most  $\eta$ . On the event  $\{t \ge N_1, s_0 = +\infty, \sup_{q\ge N_1} v_q \le 3\varepsilon/2\}$ ,  $g_k = \tilde{w}_k$  if k > t, while  $g_k = 0 \ge \tilde{w}_k - 3\varepsilon/2$  if  $k \le t$ . Therefore,

$$\mathbf{E}\left[g_{k}\mathbf{1}_{t\geq N_{1},s_{0}=+\infty}\right] \geq \mathbf{E}\left[\tilde{w}_{k}\mathbf{1}_{t\geq N_{1},s_{0}=+\infty}\right] - 3\varepsilon/2 - \varepsilon$$
$$\geq \mathbf{E}\left[\tilde{w}_{N_{1}}\mathbf{1}_{t\geq N_{1},s_{0}=+\infty}\right] - 5\varepsilon/2,$$

where the second inequality uses the fact that  $\{t \ge N_1, s_0 = +\infty\} \in \mathscr{F}_{N_1}$ , and the submartingale property of  $(\tilde{w}_n)_n$ .

Thus,

$$\mathbf{E}\left[g_k\right] \geq \mathbf{E}\left[\tilde{w}_{N_1}\right] - 5\varepsilon/2 - 2\varepsilon \geq w - 9\varepsilon/2,$$

where the second inequality uses  $w = \mathbf{E}[w_0]$  and the submartingale property of  $(\tilde{w}_n)_n$ . The result follows from (16).

5.3. An  $\epsilon$ -optimal strategy for player 1 – II

By Lemma 3.8, if  $\mathbf{P}(t_1 < +\infty) = 1$  under  $\mathbf{x}^*$ , then  $\mathbf{x}^*$  guarantees w for player 1. Therefore, we assume from now on that under  $\mathbf{x}^*$ 

$$\mathbf{P}(t_1 < +\infty) < 1. \tag{17}$$

Recall that  $\varepsilon > 0$  is given, and that  $\eta > 0$  is such that

$$\mathbf{P}(A) < \eta \Rightarrow \mathbf{E}[m1_A] < \varepsilon. \tag{18}$$

Assume moreover that  $\eta \mathbf{E}[m] \leq \varepsilon$ .

Recall also that  $N_0$  is such that, under  $(\mathbf{x}^*, \mathbf{0})$ ,

$$\mathbf{P}\left(\sup_{n,m\geq N_0}|v_n-v_m|>\varepsilon/2, t_1\geq N_0\right)<\eta.$$
(19)

By (17) we can assume w.l.o.g. that  $N_0$  is sufficiently large so that under  $\mathbf{x}^*$ ,

$$\mathbf{P}(t_1 < +\infty \mid t_1 \ge N_0) < \eta.$$

Define the strategy  $\hat{\mathbf{x}}$  by

$$\hat{x}_n = \begin{cases} \min\{x_n^* + \eta, 1\} & \text{if } n \ge N_0 \text{ and } \epsilon < v_{N_0} < a_n + \epsilon \\ x_n^* & \text{otherwise.} \end{cases}$$

We will prove that  $\gamma(\hat{\mathbf{x}}, \mathbf{y}) \ge w - 9\varepsilon$ , for every  $\mathbf{y}$ . The stronger statement:  $\gamma_n(\hat{\mathbf{x}}, \mathbf{y}) \ge v - 6\varepsilon$ , for every  $n \ge N_1$  and every  $\mathbf{y}$  also holds, provided  $N_1$  is large enough. We will not provide a proof.

## Lemma 5.4. One has

 $\limsup_{n} a_n \ge \limsup_{n} v_n \text{ on the event } \{\limsup_{n} v_n > 0\}.$ 

*Proof.* Let  $\lambda > 0$  and  $q \in \mathbf{N}$  be given. Denote by  $\tilde{\mathbf{x}}(\lambda)$  the strategy that coincides with **0** for n < q, and with  $\mathbf{x}(\lambda)$  for  $n \ge q$ . From Corollary 3.7, under  $(\tilde{\mathbf{x}}(\lambda), \mathbf{0})$ ,

$$(1-\lambda)^{-q}\mathbf{E}\left[r(\tilde{\mathbf{x}}(\lambda),\mathbf{0})(1-\lambda)^{t+1}\mathbf{1}_{t<+\infty}|\mathscr{F}_q\right] \ge v_q(\lambda).$$
(20)

Since player 2 never stops,  $t = t_1$  and  $r = a_{t_1}$  on  $t < +\infty$ . Since  $(1 - \lambda)^{t+1-q} \le 1$ , **P**-a.s., the left-hand side of (20) is at most

$$\mathbf{E}\left[a_t^+ \mathbf{1}_{t<+\infty} | \mathscr{F}_q\right] \leq \mathbf{E}\left[\sup_{n\geq q} a_n^+ | \mathscr{F}_q\right],$$

with  $a_n^+ = \max(a_n, 0)$ . Using (20),

$$\mathbf{E}\left[\sup_{n\geq q}a_n^+|\mathscr{F}_q\right]\geq v_q(\lambda).$$

By letting  $\lambda$  go to zero, one obtains  $\mathbf{E}\left[\sup_{n\geq q} a_n^+ | \mathscr{F}_q\right] \geq v_q$ . The sequence  $(\mathbf{E}\left[\sup_{n\geq q} a_n^+ | \mathscr{F}_q\right])_q$  converges **P**-a.s. to  $\limsup_q a_q^+$ . Therefore  $\limsup_q a_q^+ \geq \limsup_q v_q$ . Thus, on the event  $\{\limsup_n v_n > 0\}$ ,  $\limsup_n a_n \geq \limsup_n v_n$ , as desired.

Set 
$$\Omega_1 = \{t \ge N_0, v_{N_0} > \varepsilon\} \in \mathscr{F}_{N_0}$$
.

Proposition 5.5. Let y be given. One has

$$\mathbf{E}\left[r(\hat{\mathbf{x}},\mathbf{y})\mathbf{1}_{\Omega_{1}}\mathbf{1}_{t_{2}=t<+\infty}\right] \geq \mathbf{E}\left[v_{N_{0}}\mathbf{1}_{\Omega_{1}}\mathbf{1}_{t_{2}=t<+\infty}\right] - 4\varepsilon$$

under  $(\hat{\mathbf{x}}, \mathbf{y})$ .

*Proof.* We explicit the idea that, if player 2 stops at stage n, the corresponding expected payoff (where the expectation is taken with respect to player 1's decision) is at least  $v_n$ , up to  $\eta m$ , since player 1 plays  $x_n^*$  up to  $\eta$ .

Recall that  $\mathscr{F}_n^2 = \sigma(\mathscr{F}_n, Y_n)$ , so that  $\mathscr{F}_n^2$  includes past and present values of the payoff processes, past "decisions" of the players and the decision of player 2 at stage *n*. Observe that  $\{t_2 = t = n\} \in \mathscr{F}_n^2$ , and that by assumption  $X_n$  is independent of  $\mathscr{F}_n^2$ . Therefore, on the event  $\{t_2 = t = n\}$ ,

$$\mathbf{E}\left[r(\hat{\mathbf{x}},\mathbf{y})|\mathscr{F}_{n}^{2}\right] = G_{n}\left(\hat{x}_{n},1;\mathbf{E}\left[v_{n+1}|\mathscr{F}_{n}\right]\right)$$

(Note that the variable  $\mathbf{E}[v_{n+1}|\mathscr{F}_n]$  is here irrelevant). Since  $x_n$  is an optimal strategy in the local game  $G_n(\mathbf{E}[v_{n+1}|\mathscr{F}_n])$ , by Lemma 3.4,

$$G_n(x_n, 1; \mathbf{E}[v_{n+1}|\mathscr{F}_n]) \geq \operatorname{val} G_n(\mathbf{E}[v_{n+1}|\mathscr{F}_n]) \geq v_n.$$

Since  $|x_n - \hat{x}_n| \leq \eta$ ,

$$|G_n(x_n, 1; \mathbf{E}[v_{n+1}|\mathscr{F}_n]) - G_n(\hat{x}_n, 1; \mathbf{E}[v_{n+1}|\mathscr{F}_n])| \leq \eta m,$$

so that  $\mathbf{E}\left[r(\hat{\mathbf{x}}, \mathbf{y})|\mathscr{F}_n^2\right] \ge v_n - \eta m$  on the event  $\{t_2 = n = t\}$ . In other words,

$$\mathbf{E}\left[r(\hat{\mathbf{x}}, \mathbf{y})\mathbf{1}_{t_2=n=t} | \mathscr{F}_n^2\right] \ge (v_n - \eta m)\mathbf{1}_{t_2=n=t}, \mathbf{P}\text{-a.s.}$$

By first taking conditional expectations given  $\mathscr{F}_{N_0}$ , and then summing over  $n \ge N_0$ , one obtains

$$\mathbf{E}\left[r(\hat{\mathbf{x}}, \mathbf{y})\mathbf{1}_{N_0 \le t_2 = t < +\infty} | \mathscr{F}_{N_0}\right] \ge \mathbf{E}\left[\inf_{n \ge N_0} v_n \mathbf{1}_{t_2 = t = n} | \mathscr{F}_{N_0}\right] \\ -\eta \mathbf{E}\left[m\mathbf{1}_{N_0 \le t_2 = t < +\infty} | \mathscr{F}_{N_0}\right]$$

which yields

$$\mathbf{E}\left[r(\hat{\mathbf{x}},\mathbf{y})\mathbf{1}_{\Omega_{1}}\mathbf{1}_{t_{2}=t<+\infty}\right] \geq \mathbf{E}\left[\mathbf{1}_{\Omega_{1}}\inf_{n\geq N_{0}}v_{n}\mathbf{1}_{t_{2}=t=n}\right] - \eta \mathbf{E}\left[m\mathbf{1}_{\Omega_{1}}\mathbf{1}_{t_{2}=t<+\infty}\right].$$

Define  $\Omega_2 = \Omega_1 \cap \{\sup_{n,m>N_0} |v_n - v_m| \le \varepsilon/2\}$ . Thus,  $\mathbf{P}(\Omega_1 \setminus \Omega_2) < \eta$ , therefore

$$\mathbf{E}\left[\mathbf{1}_{\Omega_1}\inf_{n\geq N_0}v_n\mathbf{1}_{t_2=t=n}\right]-\mathbf{E}\left[\mathbf{1}_{\Omega_2}\inf_{n\geq N_0}v_n\mathbf{1}_{t_2=t=n}\right]\leq \mathbf{E}\left[\mathbf{1}_{\Omega_1\setminus\Omega_2}m\right]\leq \varepsilon.$$

On  $\Omega_2$ ,  $\inf_{n\geq N_0} v_n \geq v_{N_0} - \varepsilon/2$ . One finally gets

$$\mathbf{E}\left[r(\hat{\mathbf{x}}, \mathbf{y})\mathbf{1}_{\Omega_{1}}\mathbf{1}_{t_{2}=t<+\infty}\right] \geq \mathbf{E}\left[v_{N_{0}}\mathbf{1}_{\Omega_{1}}\mathbf{1}_{t_{2}=t<+\infty}\right] - \frac{\varepsilon}{2}\mathbf{P}(\Omega_{1} \cap \{t_{2}=t<+\infty\}) - 3\varepsilon.$$
(21)

Proposition 5.6. Let y be given. One has

$$\mathbf{E}\left[r(\hat{\mathbf{x}},\mathbf{y})\mathbf{1}_{\Omega_1}\mathbf{1}_{t_1 < t_2}\right] \geq \mathbf{E}\left[v_{N_0}\mathbf{1}_{\Omega_1}\mathbf{1}_{t_1 < t_2}\right] - 2\varepsilon.$$

under  $(\hat{\mathbf{x}}, \mathbf{y})$ .

*Proof.* Fix a strategy **y**. Note that  $\Omega_1 \cap \{t_1 < t_2\} = \{N_0 \le t_1 < t_2, v_{N_0} > \epsilon\}$ , and on this set,  $r(\hat{\mathbf{x}}, \mathbf{y}) = a_{t_1}$ .

By the definition of  $N_0$ ,  $\mathbf{P}(t_1 = t \ge N_0, a_t \le v_{N_0} - \epsilon) < \eta$ . In particular,  $\mathbf{P}(N_0 \le t_1 < t_2, a_{t_1} > v_{N_0} - \epsilon > 0) > \mathbf{P}(\Omega_1 \cap \{t_1 < t_2\}) - \eta$ . The result follows from (18).

**Lemma 5.7.** For every  $\mathbf{y}$ ,  $\gamma(\hat{\mathbf{x}}, \mathbf{y}) \ge w - 9\epsilon$ .

*Proof.* Define the stopping time  $\theta$  by  $\theta = N_0$  on  $\Omega_1 = \{t \ge N_0, v_{N_0} > \epsilon\}$ , and  $\theta = +\infty$  otherwise. The strategy  $\hat{\mathbf{x}}$  coincides with  $\mathbf{x}^*$  up to  $\theta$ . Therefore,  $(\tilde{v}_n)$  is a submartingale up to  $\theta$ .

Notice that  $\theta = +\infty$  if  $\theta > N_0$ ; therefore  $(\tilde{v}_n)$  converges, **P**-a.s. on the event  $\{\theta > N_0\}$ , say to  $\tilde{v}_{\infty}$ .

Given the integrability properties of  $(\tilde{v}_n)$ , one has

$$\mathbf{E}(\tilde{v}_{\theta}) \ge \mathbf{E}(\tilde{v}_{0}) = w. \tag{22}$$

By definition of  $(\tilde{v}_n)$ , one has  $\tilde{v}_{\infty} = r(\hat{\mathbf{x}}, \mathbf{y})$  if  $t < +\infty$ ,  $\tilde{v}_{\infty} \le 3\epsilon/2$  if  $t = +\infty$ and  $\sup_{n,m \ge N_0} |v_n - v_m| \le \epsilon/2$ , and  $\tilde{v}_{\infty} \le m$  otherwise. Thus, by (18) and (19),

$$\mathbf{E}[\tilde{v}_{\infty} \mathbf{1}_{\theta > N_0}] \leq \mathbf{E}[r(\hat{\mathbf{x}}, \mathbf{y}) \mathbf{1}_{t < +\infty} \mathbf{1}_{\theta > N_0}] + 3\epsilon/2 + \epsilon.$$

The inequality (22) may be rewritten as

$$\mathbf{E}[v_{N_0}\mathbf{1}_{\Omega_1} + \tilde{v}_{\infty}\mathbf{1}_{\theta > N_0}] \ge w.$$

and therefore, using Propositions 5.5 and 5.6,  $\mathbf{E}(r(\hat{\mathbf{x}}, \mathbf{y}) \mathbf{1}_{t < +\infty}) \ge w - \frac{17}{2}\epsilon$ , and the result follows.

## 6. Concluding remarks

This section contains a discussion of related issues. We first discuss an alternative way of introducing randomization into stopping games<sup>1</sup>. We then discuss a fairly easy extension of our main result.

We introduced randomization by allowing the players, at any stage, to stop with a probability between zero and one. These strategies are usually called *behavior strategies* in the game theory literature. We might as well consider the possibility for a player to select randomly a (deterministic) stopping time at the beginning of the game, thereby extending differently the set of available strategies. These strategies are called *mixed strategies*. For many classes of games, the two extensions are equivalent in a strong sense. The first equivalence result is due to Kuhn (1953).

For stopping games (as for many other games), the definition of mixed strategies as suggested here is problematic, since it requires to define a convenient measurable structure on the set of stopping times. There are two ways to avoid this problem.

Following Aumann (1964), one may enlarge the probability space from  $(\Omega, \mathcal{A}, \mathbf{P})$  to  $(\Omega \times [0, 1], \mathcal{A} \otimes \mathcal{B}, \mathbf{P} \otimes \lambda_1)$ , where  $\lambda_1$  is the Lebesgue measure. A mixed strategy (for Player 1) is then defined as an  $\mathcal{A} \otimes \mathcal{B}$ -measurable function  $\phi$  from  $\Omega \times [0, 1]$  to  $\mathbf{N} \cup \{+\infty\}$  such that

for  $\lambda_1$ -a.e.  $r \in [0, 1]$ ,  $\phi(r, \cdot)$  is a stopping time.

Intuitively, ([0, 1],  $\lambda_1$ ) is a randomizing device for player 1. We introduce an independent copy ([0, 1],  $\lambda_2$ ) for player 2.

We claim that these mixed strategies are equivalent to behavioral strategies. Denote  $\sigma_r = \phi(r, \cdot)$ . Then  $\sigma_r$  is  $\lambda_1$ -a.e. a stopping time. For each mixed strategy  $\phi$  and every  $n \in \mathbb{N}$ , define  $H(\phi)_n = \int \mathbb{1}_{\{\sigma_r \le n\}} \lambda_1(dr)$  the probability under  $\phi$  that player 1 stops prior to stage n + 1. Clearly,  $(H(\phi)_n)$  is  $(\mathcal{F}_n)$ -adapted. It can be viewed as the (random) distribution function corresponding to some behavior strategy  $\mathbf{x}$ , that we denote by  $h(\phi)$ . The map h from mixed to behavior strategies is onto. Indeed, given a behavior strategy  $\mathbf{x}$ , denote by  $F^{\mathbf{x}}$  the distribution function of  $t_1$ . Set  $\phi^{\mathbf{x}}(r, \omega) := \inf\{n \ge 0, F^{\mathbf{x}}(n, \omega) \ge r\}$ . Then  $\phi^{\mathbf{x}}$  is a mixed strategy, such that  $h(\phi^x) = \mathbf{x}$ . It is easy to verify that, for each pair  $(\phi, \psi)$  of mixed strategies, the expected payoff under  $(\phi, \psi)$  coincides with the expected payoff under the pair  $(h(\phi), h(\psi))$  of behavior strategies. For more details, see Touzi and Vieille (1999).

Another approach to define mixed strategies is due to Bismut (1977): it consists of interpreting such a strategy as an element of the dual space of a Banach space containing the stopping times, and of using functional analysis methods.

We argue now that the first proof of the main result can be extended to handle a larger class of stochastic games.<sup>2</sup> The class of games we consider now is the following. Each player has finitely many actions. The sets of actions are respectively *A* and *B* for the two players. The two players choose repeatedly elements from *A* and *B*. For each pair  $(a, b) \in A \times B$ , two processes  $(g_n^{a,b})_n$  and  $(p_n^{a,b})_n$  are given:  $p_n^{a,b}$  is the probability that the game stops in stage *n*, if (a, b) is played in that stage

<sup>&</sup>lt;sup>1</sup> We thank a referee for pointing out the issue.

<sup>&</sup>lt;sup>2</sup> We thank Sylvain Sorin for suggesting this generalization.

and the game has not stopped earlier;  $g_n^{a,b}$  is the payoff that is received by player 1 in that case. The payoff is zero if the game never stops.

In words, those are games where the actions of the players may influence the probability of termination and the terminal payoff, but, if the game continues, they do not influence the information of the players at the next stage.

Clearly, stopping games belong to this class, with  $A = B = \{\text{stop, continue}\}$ , and  $p_n^{a,b} = 0$  if a = b = continue, and  $p_n^{a,b} = 1$  otherwise. To specify properly the game, we need to tell what is known at stage *n* about past choices of the players. This turns here to be irrelevant (in contrast with other classes of stochastic games).

We briefly sketch how the proof in Section 5.2 has to be adapted. All notations are the same. The only difficulty lies in defining  $N_0$ , since, loosely speaking, there exists no least terminating strategy. Partition  $\Omega$  into  $\Omega_c$  and  $\Omega_d$ , where  $\Omega_c \in \mathscr{A} = F_{\infty}$  is the convergence set of the sequence  $(v_n)_n$ , and  $\Omega_d = \Omega \setminus \Omega_c$ . We choose an integer  $N_0$  large enough and an event  $F \in \mathscr{F}_{N_0}$  such that  $\mathbf{P}(F \Delta \Omega_c) < \eta$ .

We define a strategy  $\hat{\mathbf{x}}$  that has the following features: it coincides with  $\mathbf{x}^*$  unless F occurs and  $v_{N_0} > \epsilon$ ; in that case, it switches at stage  $N_0$  to the strategy we defined in section 5.2, *i.e.*, it plays a sequence of locally optimal strategies in properly chosen discounted games if  $v_{N_0} > \epsilon$ .

It can be shown that  $\hat{\mathbf{x}}$  guarantees w up to  $7\varepsilon$ .

It is not clear whether the second proof can be generalized to this class of games.

We conclude with a brief discussion on our assumptions related to the filtration  $(\mathscr{F}_n)_n$ . We assumed that the payoff processes  $(a_n)$ ,  $(b_n)$  and  $(c_n)$  are adapted, and that  $X_n$  and  $Y_n$  are independent of  $\mathscr{F}_n$  but  $\mathscr{F}_{n+1}$ -measurable. As we argued previously, the first assumption can be totally dispensed with. Informally, the second assumption means that (i) in any stage, each player has no information about the action the other player is about to choose, and (ii) past choices are observed. The first part of the assumption is crucial, but the second is irrelevant. Observe indeed that our  $\varepsilon$ -optimal strategies make no use of the past actions of the opponent. Finally, it is crucial that both players have the same filtration. The existence of the value does not extend to the situation where the payoff processes are constant, and one of the players has more information than the other about their value.

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