# Probabilistic methods in combinatorics 

Homework assignment \#3
Due date: Monday, May 27, 2019

Problem 1. Let $\mathcal{P} \mathcal{D}\left(K_{4}\right)$ denote the following graph property: $G \in \mathcal{P} \mathcal{D}\left(K_{4}\right)$ if and only if $G$ contains a collection of $\lfloor|V(G)| / 8\rfloor$ pairwise vertex-disjoint copies of $K_{4}$. Find a function $\theta: \mathbb{N} \rightarrow[0,1]$ such that

$$
\lim _{n \rightarrow \infty} \operatorname{Pr}\left(G_{n, p(n)} \in \mathcal{P} \mathcal{D}\left(K_{4}\right)\right)= \begin{cases}1 & \text { if } p(n) \gg \theta(n)  \tag{1}\\ 0 & \text { if } p(n) \ll \theta(n)\end{cases}
$$

(The problem asks not only to determine $\theta$ but also to prove that this $\theta$ satisfies (1).)
Problem 2. For a graph $G$, let $\alpha_{3}(G)$ denote the largest size of a set $U \subseteq V(G)$ such that the subgraph $G[U]$ induced by $U$ contains no triangles. Show that there are constants $c, C>0$ such that

$$
\operatorname{Pr}\left(c \log n \leqslant \alpha_{3}\left(G_{n, 1 / 2}\right) \leqslant C \log n\right)=1-o(1) .
$$

Problem 3. Prove that, for every positive $\varepsilon$, there is an $n_{0}$ such that, for every $n>n_{0}$, there is an $n$-vertex graph that contains every graph on $\left\lfloor(2-\varepsilon) \log _{2} n\right\rfloor$ vertices as an induced subgraph.

Problem 4. Show that moment bounds for tail probabilities are always better than CramérChernoff bounds. More precisely, let $X$ be a nonnegative random variable such that $\mathbb{E}\left[e^{\lambda X}\right]<\infty$ for every $\lambda \geqslant 0$ and let $t>0$. Prove that

$$
\inf _{q \in \mathbb{N}} \frac{\mathbb{E}\left[X^{q}\right]}{t^{q}} \leqslant \inf _{\lambda>0} \frac{\mathbb{E}\left[e^{\lambda X}\right]}{e^{\lambda t}},
$$

where $\mathbb{N}$ is the set of natural numbers, that is, $\mathbb{N}=\{0,1,2, \ldots\}$.
Problem 5. Show that there exists an $n_{0}$ such that the following holds. Let $G$ be a graph with $n \geqslant n_{0}$ vertices and minimum degree $\delta(G) \geqslant(\log n)^{2}$. The vertex set of $G$ may be partitioned into three sets $V_{1}, V_{2}$, and $V_{3}$ such that $\delta\left(G\left[V_{i}\right]\right) \geqslant 0.33 \cdot \delta(G)$ for every $i$.

Problem 6. Let $G$ be a graph with $\chi(G)=2000$. Let $U$ be a subset of $V(G)$ selected uniformly at random and let $H=G[U]$ be the subgraph of $G$ induced by $U$. Prove that

$$
\operatorname{Pr}(\chi(H) \leqslant 900) \leqslant \frac{1}{10} .
$$

## Please do NOT submit written solutions to the following exercises:

Exercise 1. Let $X$ be a square-integrable real-valued random variable and let $m$ be a median of $X$, that is, $m$ is a number such that both $X \leqslant m$ and $X \geqslant m$ hold with probability at least $1 / 2$. Show that

$$
|\mathbb{E}[X]-m| \leqslant \sqrt{\operatorname{Var}(X)} .
$$

