

(B6) $2B_k(m)+2 \leq A_k(m+3)$, for $k \geq 1$, $m \geq 0$.

(B7) $A_{k-1}(m) \leq B_k(m) \leq 2B_k(m) \leq A_k(m+3)$, for $k \geq 4$, $m \geq 1$, so that the sequences of functions B_k and A_k have the same asymptotic order of growth.

Next we define another sequence of functions $\{C_k\}$ by putting $C_k(m) = 2B_k(m)$, for all k and m . An explicit recursive definition of these functions, for positive k and m , is

$$\begin{aligned} C_1(m) &= 1, & m \geq 1, \\ C_k(1) &= 2C_{k-1}(2), & k \geq 2, \\ C_k(m) &= C_k(m-1) \cdot C_{k-1}(C_k(m-1)), & k \geq 2, m \geq 2. \end{aligned} \quad (2.4)$$

It follows easily from the preceding analysis that

(C1) $C_2(m) = 2$, for $m \geq 0$.

(C2) $C_3(m) = 2^{m+1}$, for $m \geq 0$.

(C3) $C_4(m) \geq 2^{2^{m-2}}$, with $m+1$ 2's in the exponential tower.

(C4) $A_{k-1}(m) \leq C_k(m) \leq A_k(m+3)$, for $k \geq 4$, $m \geq 1$, so that the growth of the sequences of functions $\{C_k\}$ and $\{A_k\}$ are also of the same asymptotic order of magnitude.

In what follows we will often use the shorthand notations

$$\bar{\alpha} = C_k(m-1), \quad \bar{\beta} = C_{k-1}(C_k(m-1)) = C_{k-1}(\bar{\alpha})$$

and $\bar{\gamma} = C_k(m) = \bar{\alpha} \cdot \bar{\beta}$ (by definition).

2.3.2 Generation of superlinear-size sequences

For each $k, m \geq 1$, the sequence $S(k, m)$ that we are going to construct will satisfy the following two properties:

(1) $S(k, m)$ is composed of $N_k(m) = m \cdot C_k(m)$ distinct symbols. (These symbols are named (d, γ) , for $d = 1, \dots, m$, $\gamma = 1, \dots, \bar{\gamma}$, and are ordered in lexicographical order, so that $(d', \gamma') < (d'', \gamma'')$ if $\gamma' < \gamma''$ or $\gamma' = \gamma''$ and $d' < d''$.)

(2) $S(k, m)$ contains $C_k(m)$ fans of size m , where each fan is a contiguous subsequence of the form

$$((1, \gamma)(2, \gamma) \dots (m, \gamma)).$$

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Since fans are pairwise disjoint by definition, the naming scheme of the symbols of $S(k, m)$ can be interpreted as assigning to each symbol the index γ of the fan in which it appears, and its index d within that fan.

The construction proceeds by double induction on k and m , as follows.

1. $k = 1$: The sequence is a single fan of size m :

$$S(1, m) = ((1, 1)(2, 1) \dots (m, 1)).$$

Properties (1) and (2) clearly hold here ($C_1(m) = 1$).

2. $k = 2$: The sequence contains a pair of disjoint fans of size m , with a block following each of these fans. Specifically,

$$\begin{aligned} S(2, m) &= ((1, 1)(2, 1) \dots (m-1, 1)(m, 1)(m-1, 1) \dots (1, 1) \\ &\quad (1, 2)(2, 2) \dots (m-1, 2)(m, 2)(m-1, 2) \dots (1, 2)). \end{aligned}$$

Indeed, $S(2, m)$ contains $C_2(m) = 2$ fans and is composed of $2m$ symbols.

3. $k \geq 3, m = 1$: The sequence is identical to the sequence for $k' = k-1$ and $m' = 2$, except for renaming of its symbols and fans: $S(k-1, 2)$ contains $C_{k-1}(2) = \frac{1}{2}C_k(1)$ fans, each of which consists of two symbols; the symbol renaming in $S(k, 1)$ causes each of these two elements to become a 1-element fan. Properties (1) and (2) clearly hold.

4. The general case $k \geq 3, m > 1$:

- (i) Generate inductively the sequence $S' = S(k, m-1)$; by induction, it contains $\bar{\alpha}$ fans of size $m-1$ each and is composed of $(m-1) \cdot \bar{\alpha}$ symbols.
- (ii) Create $\bar{\beta}$ copies of S' whose sets of symbols are pairwise disjoint. For each $\beta \leq \bar{\beta}$, rename the symbols in the β th copy S'_β of S' as (d, α, β) where $1 \leq d \leq m-1$ is the index of the symbol in the fan of S'_β containing it, and $1 \leq \alpha \leq \bar{\alpha}$ is the index of this fan in S'_β .
- (iii) Generate inductively the sequence $S^* = S(k-1, \bar{\alpha})$ whose set of symbols is disjoint from that of any S'_β ; by induction, it contains $\bar{\beta}$ fans of size $\bar{\alpha}$ each. Rename the symbols of S^* as (m, α, β) (where α is the index of that symbol within its fan, and β is the index of that fan in S^*). Duplicate the last element $(m, \bar{\alpha}, \beta)$ in each of the $\bar{\beta}$ fans of S^* .
- (iv) For each $1 \leq \alpha \leq \bar{\alpha}$, $1 \leq \beta \leq \bar{\beta}$, extend the α th fan of S'_β by duplicating its last element $(m-1, \alpha, \beta)$, and by inserting the corresponding symbol (m, α, β) of S^* between these duplicated appearances of $(m-1, \alpha, \beta)$. This process extends the $(m-1)$ -fans of S'_β into m -fans and adds a new element after each extended fan.

- (v) Finally construct the desired sequence $S(k, m)$ by merging the β copies S'_β of S' with the sequence S^* . This is done by replacing, for each $1 \leq \beta < \bar{\beta}$, the β th fan of S^* by the corresponding copy S'_β of S' , as modified in (iv) above. Note that the duplicated copy of the last element in each fan of S^* (formed in step (iii) above) appears now after the copy S'_β that replaces this fan; see Figure 2.1 for an illustration of this process.

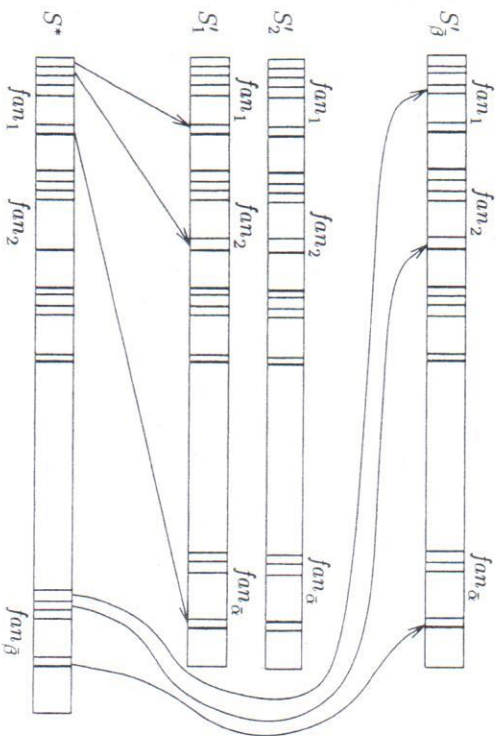


FIGURE 2.1. Merging the subsequences.

To establish property (1), note that $S(k, m)$ consists of

$$\begin{aligned} N_k(m) &= \bar{\beta} \cdot (m-1)\bar{\alpha} + \bar{\alpha}C_{k-1}(\bar{\alpha}) \\ &= (m-1)\bar{\alpha}\bar{\beta} + \bar{\alpha}\bar{\beta} \\ &= m\bar{\alpha}\bar{\beta} \\ &= mC_k(m) \end{aligned}$$

symbols. Property (2) is trivial, because the fans of $S(k, m)$ are precisely the extended fans of the copies S'_β of S' , and their number is $C_k(m-1) \cdot \bar{\beta} = C_k(m)$.

We now establish several important properties of the sequences $S(k, m)$. For our present purpose, property (a) is all we need. However, later on in Chapter 4 we will be concerned with geometric realization of the sequences $S(k, m)$, and there we will need to use the other properties.

Theorem 2.20 For each $k, m \geq 1$ the sequence $S = S(k, m)$ satisfies the following properties:

- (a) S is a $DS(N_k(m), 3)$ -sequence.
 (b) Each symbol of S appears in precisely one fan and makes its first (leftmost) appearance in S .
 (c) For $k \geq 2$ and for each $\gamma \leq \bar{\gamma}$, the last element (m, γ) of the γ th fan of S forms the beginning of a contiguous subsequence that is the reverse of that fan:

$$((m, \gamma)(m-1, \gamma) \cdots (2, \gamma)(1, \gamma)).$$

(Note that this sequence is the initial portion of a block of S .)

- (d) For each block c of S , let f be the rightmost fan preceding or including c and let c_1, c_2, \dots, c_t be the blocks appearing in S between f and c , for some $t \geq 0$. Let a be the first (leftmost) element of c ; then either this appearance of a is within f , or else a must also appear in one of the preceding blocks c_i .

Remark 2.21 (i) For each $\gamma \leq \bar{\gamma}$ and each $d < m$, the element (d, γ) in the γ th fan of S forms a 1-element block. Note that property (d) is trivially correct for these singleton blocks.

- (ii) Property (b) implies in particular that S starts with a fan.

(iii) Unless c is one of the singleton blocks mentioned in (i) above, the first block c_1 in property (d) is the block mentioned in (c) (whose initial portion is the reverse of the fan f). Note that property (d) clearly holds for the case $c = c_1$.

Proof. The proof proceeds by double induction on k and m . The base case $k = 1$ is trivial: $S(1, m)$ is plainly a $DS(m, 3)$ -sequence, (b) and (d) are trivial, and (c) is vacuous in this case.

The case $k = 2$ is also easy. Here $\bar{\gamma} = 2$ and $S(2, m)$ is obviously a $DS(2m, 3)$ -sequence, so (a) follows. Properties (b), (c), and (d) are also immediate.

Next consider the case $k > 2$, $m = 1$. Here $S(k, 1) = S(k-1, 2)$ (with its symbols being renamed), so property (a) holds by induction. Property (b) is also trivial because the only change in the fan structure between $S(k-1, 2)$ and $S(k, 1)$ is that each fan is split into two subfans. Since each fan is now of size 1, property (c) is trivial too. Finally, since the block structure in $S(k, 1)$ is identical to that in $S(k-1, 2)$, (d) also follows immediately by induction.

Finally consider the general case $k > 2$, $m > 1$. We first prove property (a). First note that no two adjacent elements of $S = S(k, m)$ are equal. Indeed, by the induction hypothesis, no two adjacent elements either in S^* or in any S'_β are equal; all these sequences have pairwise disjoint sets of symbols, and the