

1 Random walk on a graph

Assume that a connected finite oriented graph has m vertices, and each vertex has k outgoing edges and k incoming edges (the same k for all vertices). Denote the set of vertices by V and the set of edges by E ; $E \subset V \times V$ (it may intersect the diagonal). (Multiple edges are thus excluded, but all said can be easily generalized to graphs with multiple edges.) It is assumed that $E \cap ((A \times (V \setminus A)) \cup ((V \setminus A) \times A)) \neq \emptyset$ for every $A \subset V$ such that $A \neq \emptyset$ and $V \setminus A \neq \emptyset$ (weak connectedness). Also, $\#\{y \in V : (x, y) \in E\} = \#\{y \in V : (y, x) \in E\} = k$ for all $x \in V$. In addition, we assume *aperiodicity*: there exists no $p \in \{2, 3, \dots\}$ such that every loop length is divisible by p . (A loop is a sequence of vertices (y_0, y_1, \dots, y_t) such that $(y_0, y_1) \in E, \dots, (y_{t-1}, y_t) \in E$ and $y_t = y_0$; its length is t .)

A random walk started at a given vertex x_0 . A path (of length n) of the random walk is a sequence (s_0, \dots, s_n) of vertices such that the pairs (s_{k-1}, s_k) belong to E (for $k = 1, \dots, n$) and $s_0 = x_0$. There are k^n such paths; each has the probability k^{-n} (by definition). We have the probability space Ω of paths, and random variables $S_0, \dots, S_n : \Omega \rightarrow V$.

We start with some graph-theoretic (non-probabilistic) statements.

1.1 Lemma. For every $A \subset V$, the number of incoming edges is equal to the number of outgoing edges. That is,

$$\#(E \cap (A \times (V \setminus A))) = \#(E \cap ((V \setminus A) \times A)).$$

Proof. We have

$$\begin{aligned} E \cap (A \times V) &= E \cap (A \times (V \setminus A)) \uplus E \cap (A \times A), \\ E \cap (V \times A) &= E \cap ((V \setminus A) \times A) \uplus E \cap (A \times A) \end{aligned}$$

and

$$\#(E \cap (A \times V)) = k \cdot \#A = \#(E \cap (V \times A)).$$

□

1.2 Corollary. (Strong connectedness.)

$$E \cap (A \times (V \setminus A)) \neq \emptyset \text{ for every } A \subset V \text{ such that } A \neq \emptyset \text{ and } V \setminus A \neq \emptyset.$$

1.3 Corollary. For all $x, y \in V$ there exists a path (of *some* length) from x to y .

1.4 Lemma. There exists T such that for all $x, y \in V$, every $t \geq T$ is the length of some (at least one) path from x to y .

Proof. (Sketch.)

The set L_x of lengths of all loops from x to x is a semigroup, therefore $L_x - L_x$ is a group, $L_x - L_x = p_x\mathbb{Z}$ for some p_x . The period p_x does not depend on x . Thus, $p_x = 1$ for all x . It means existence of N such that $N \in L_x$ and $N + 1 \in L_x$. We take $T = N^2$ and note that $N^2 + kN + r = N(N + k) + r = N(N + k - r) + (N + 1)r \in L_x$. Generalization from $x = y$ to all x, y is easy. \square

Now we return to probability. We want to show that the initial point x_0 is ultimately forgotten by the Markov chain.

Given another starting point $x'_0 \in V$, we introduce the probability space Ω' of paths (of length n) starting at x'_0 , and random variables $S'_0, \dots, S'_n : \Omega' \rightarrow V$. We take the product

$$\tilde{\Omega} = \Omega \times \Omega'$$

and treat S_t, S'_t as maps $\tilde{\Omega} \rightarrow V$. We get two *independent* random walks, one starting at x_0 , the other at x'_0 . In addition, we let $\tilde{S}_t = (S_t, S'_t) : \tilde{\Omega} \rightarrow \tilde{V} = V \times V$.

Recall the reflection principle, instrumental in ‘Extremal values, etc.’ It will help again! The transformation $(x, y) \mapsto (y, x)$ of \tilde{V} will be treated as reflection, and the diagonal of \tilde{V} as the barrier. We define $M_n : \tilde{\Omega} \rightarrow \{0, 1\}$ by

$$M_n = \begin{cases} 0 & \text{if } S_0 \neq S'_0, S_1 \neq S'_1, \dots, S_n \neq S'_n, \\ 1 & \text{otherwise.} \end{cases}$$

1.5 Lemma. The conditional distribution of \tilde{S}_n given $M_n = 1$ is symmetric.

That is, $\mathbb{E}(f(\tilde{S}_n) | M_n = 1) = 0$ for every antisymmetric function $f : \tilde{V} \rightarrow \mathbb{R}$ (antisymmetric means $f(y, x) = -f(x, y)$).

The proof is similar to the proof of Lemma 1 in ‘Extremal values, etc.’

1.6 Exercise. $|\mathbb{P}(S_n = x) - \mathbb{P}(S'_n = x)| \leq \mathbb{P}(M_n = 0)$.

Prove it.

Hint: $f(a, b) = \mathbf{1}_{\{x\}}(a) - \mathbf{1}_{\{x\}}(b)$.

The probability of the event $M_n = 0$ depends on n , x_0 and x'_0 . We maximize it in x_0, x'_0 :

$$\varepsilon_n = \max_{x_0, x'_0 \in V} \mathbb{P}(M_n = 0).$$

1.7 Lemma. $\varepsilon_n \rightarrow 0$ as $n \rightarrow \infty$.

The proof will be given later.

Let $p_n(x, y)$ denote the n -step transition probability from x to y . (Thus, $\mathbb{P}(S_t = y) = p_t(x_0, y)$ and $\mathbb{P}(S'_t = y) = p_t(x'_0, y)$.)

1.8 Exercise. $\sum_{y \in V} p_1(x, y) = 1$ for all $x \in V$, and $\sum_{x \in V} p_1(x, y) = 1$ for all $y \in V$.

Prove it.

Hint: k times $1/k \dots$

1.9 Exercise. $\sum_{y \in V} p_n(x, y) = 1$ for all $x \in V$, and $\sum_{x \in V} p_n(x, y) = 1$ for all $y \in V$.

Prove it.

Hint: induction in n .

1.10 Theorem. For each vertex x of the graph,

$$\mathbb{P}(S_n = x) \rightarrow \frac{1}{m} \quad \text{as } n \rightarrow \infty.$$

Proof. By 1.6, $|p_n(x_0, y) - p_n(x'_0, y)| \leq \varepsilon_n$. By 1.9, $\frac{1}{m} \sum_{x'_0 \in V} p_n(x'_0, y) = \frac{1}{m}$. Thus, $|p_n(x_0, y) - \frac{1}{m}| \leq \varepsilon_n$; finally, $\varepsilon_n \rightarrow 0$ by 1.7. \square

Proof of Lemma 1.7. Lemma 1.4 gives us T such that $p_T(x, y) \neq 0$ for all x, y . Clearly, $p_T(x, y) \geq k^{-T}$. Thus,

$$\mathbb{P}(M_T = 1) \geq \mathbb{P}(S_T = y, S'_T = y) \geq k^{-2T},$$

no matter which y is used. We put $\theta = 1 - k^{-2T}$ and see that $\mathbb{P}(M_T = 0) \leq \theta$. But moreover, $\mathbb{P}(M_{t+T} = 0 | S_t = a, S'_t = b) \leq \theta$ for all a, b (provided that the condition is of non-zero probability). It follows that

$$\begin{aligned} \mathbb{P}(M_{t+T} = 0 | M_t = 0) &\leq \theta \quad \text{for all } t; \\ \mathbb{P}(M_{t+T} = 0) &\leq \theta \cdot \mathbb{P}(M_t = 0) \quad \text{for all } t; \\ \mathbb{P}(M_{jT} = 0) &\leq \theta^j \quad \text{for all } j; \end{aligned}$$

however, $\theta^j \rightarrow 0$ as $j \rightarrow \infty$. \square

Interestingly, $\varepsilon_n \rightarrow 0$ exponentially fast. However, the constant Tk^{2T} can be quite large.

2 Finite Markov chains

A *Markov chain* (discrete in space and time, and homogeneous in time) is described by a *transition probability matrix*

$$(p(x, y))_{x, y \in V}$$

satisfying

$$p(x, y) \geq 0; \quad \forall x \sum_y p(x, y) = 1.$$

The set V is assumed to be finite. We turn V into a graph putting

$$E = \{(x, y) \in V^2 : p(x, y) \neq 0\}$$

and define the probability of a path (s_0, \dots, s_n) as the product of n probabilities

$$p(s_0, \dots, s_n) = p(s_0, s_1) \dots p(s_{n-1}, s_n);$$

as before, s_0 must be equal to a given initial point $x_0 \in V$. Here are some definitions that depend on the graph only.

A set $A \subset V$ is *closed* if $E \cap (A \times (V \setminus A)) = \emptyset$.

A Markov chain is *irreducible* if \emptyset and V are the only closed sets. In other words: for all $x, y \in V$ there exists a path from x to y (recall 1.3).

An irreducible Markov chain is *aperiodic*, if there exists no $p \in \{2, 3, \dots\}$ such that every loop length is divisible by p . (This property does not depend on the initial point; recall the proof of 1.4.)

2.1 Lemma. If the Markov chain is irreducible then

$$\lim_{n \rightarrow \infty} \mathbb{P}(S_1 \neq y, \dots, S_n \neq y) = 0$$

for each $y \in V$.

Proof. We take T and ε such that for every $x \in V$ there exists a path from x to y of length $\leq T$ and of probability $\geq \varepsilon$. Then (assuming $\mathbb{P}(S_1 \neq y, \dots, S_n \neq y) \neq 0$ for all n),

$$\mathbb{P}(S_{t+1} \neq y, \dots, S_{t+T} \neq y \mid S_t = x) \leq 1 - \varepsilon.$$

Thus

$$\begin{aligned} \mathbb{P}(S_{t+1} \neq y, \dots, S_{t+T} \neq y \mid S_1 \neq y, \dots, S_t \neq y) &= \\ &= \sum_{x \in V} \mathbb{P}(S_{t+1} \neq y, \dots, S_{t+T} \neq y \mid S_t = x) \mathbb{P}(S_t = x \mid S_1 \neq y, \dots, S_t \neq y) \leq \\ &\leq (1 - \varepsilon) \end{aligned}$$

and

$$\mathbb{P}(S_1 \neq y, S_2 \neq y, \dots, S_{jT} \neq y) \leq (1 - \varepsilon)^j \quad \text{for } j = 1, 2, \dots$$

□

2.2 Exercise. Let y belong to each nonempty closed set. (In other words: for every x there exists a path from x to y .) Prove that

$$\lim_{n \rightarrow \infty} \mathbb{P}(S_1 \neq y, \dots, S_n \neq y) = 0.$$

2.3 Exercise. Let $A \subset V$ intersect every nonempty closed set. (In other words: for every x there exists a path from x to A .) Prove that

$$\lim_{n \rightarrow \infty} \mathbb{P}(S_1 \notin A, \dots, S_n \notin A) = 0.$$

2.4 Lemma. If the Markov chain is irreducible and aperiodic, then there exists T such that for all $x, y \in V$, every $t \geq T$ is the length of some (at least one) path from x to y . That is, $p_t(x, y) > 0$.

The proof is similar to that of 1.4. (Only the graph matters.)

We may consider two independent copies of the Markov chain:

$$V^2 = V \times V, \\ p^{(2)}((x_1, x_2), (y_1, y_2)) = p(x_1, y_1)p(x_2, y_2).$$

2.5 Exercise. (a) If the Markov chain (V, p) is irreducible and aperiodic, then the Markov chain $(V^2, p^{(2)})$ is irreducible and aperiodic;

(b) it may happen that (V, p) is irreducible but $(V^2, p^{(2)})$ is not.

Prove (a) and find a counterexample for (b).

Assume that the Markov chain is irreducible and aperiodic (from now on, till Theorem 2.10).

2.6 Lemma. There exist $\varepsilon_n \rightarrow 0$ such that

$$|p_n(x_1, y) - p_n(x_2, y)| \leq \varepsilon_n$$

for all $x_1, x_2, y \in V$ and $n = 0, 1, 2, \dots$

Proof. (Sketch.) We use the reflection-type argument similarly to 1.5, 1.6, 1.7. □

2.7 Exercise. For all probability measures μ on V ,

$$\left| \sum_{x_1} \mu(x_1) p_n(x_1, y) - p_n(x_2, y) \right| \leq \varepsilon_n$$

for all $x_2, y \in V$ and $n = 0, 1, \dots$

Prove it.

Hint: $|\sum_{x_1} \mu(x_1)(p_n(x_1, y) - p_n(x_2, y))|$.

Of course, $\mu(x)$ means $\mu(\{x\})$. Substituting for μ the distribution of S_t we get

$$|\mathbb{P}(S_{t+n} = y) - p_n(x_2, y)| \leq \varepsilon_n$$

for all $x_2, y \in V$ and $t, n \in \{0, 1, \dots\}$.

2.8 Exercise. For all probability measures μ, ν on V ,

$$\left| \sum_{x_1} \mu(x_1) p_n(x_1, y) - \sum_{x_2} \nu(x_2) p_n(x_2, y) \right| \leq \varepsilon_n$$

for all $y \in V$ and $n = 0, 1, \dots$

Prove it.

2.9 Corollary.

$$|\mathbb{P}(S_t = y) - \mathbb{P}(S_u = y)| \leq \varepsilon_n$$

for all $y \in V$ and $t, u \in \{n, n+1, \dots\}$.

2.10 Theorem. If a Markov chain is irreducible and aperiodic then the limit

$$\lim_n \mathbb{P}(S_n = x)$$

exists for each $x \in V$.

Proof. By 2.9, $(\mathbb{P}(S_n = x))_n$ is a Cauchy sequence. □

We still assume that the Markov chain is irreducible and aperiodic (from now on, till Theorem 2.15).

2.11 Definition. A probability measure μ on V is *stationary*, if

$$\mu(y) = \sum_{x \in V} \mu(x) p(x, y) \quad \text{for all } y \in V.$$

2.12 Exercise. The numbers

$$\mu(x) = \lim_{n \rightarrow \infty} \mathbb{P}(S_n = x)$$

are a stationary probability measure.

Prove it.

Hint: $\mathbb{P}(S_{n+1} = y) = \sum_x \mathbb{P}(S_n = x)p(x, y)$.

2.13 Exercise. $\mu(x) > 0$ for every x .

Prove it.

Hint: otherwise there exist x, y such that $\mu(x) > 0$, $\mu(y) = 0$ and $p(x, y) > 0$.

2.14 Exercise. The measure μ defined in 2.12 is the only stationary probability measure.

Prove it.

Hint: apply 2.8 to stationary μ, ν .

2.15 Theorem. If a Markov chain is irreducible and aperiodic then it has one and only one stationary probability measure μ , and

$$\sum_{x \in V} \nu(x)p_n(x, y) \rightarrow \mu(y) \quad \text{as } n \rightarrow \infty$$

for every probability measure ν on V .

2.16 Exercise. Prove Theorem 2.15.

If a Markov chain (V, p) is irreducible but periodic, with the (least) period d , then the limit

$$\mu(x) = \lim_n \mathbb{P}(S_{nd} = x)$$

exists for each $x \in V$. The numbers $\mu(x)$ are a probability measure satisfying

$$\mu(y) = \sum_{x \in V} \mu(x)p_d(x, y) \quad \text{for all } y \in V.$$

That is, μ is stationary for the Markov chain (V, p_d) . The measure

$$\nu(x) = \lim_n \frac{1}{d} (\mathbb{P}(S_{nd} = x) + \mathbb{P}(S_{nd+1} = x) + \cdots + \mathbb{P}(S_{nd+d-1} = x))$$

is stationary for (V, p) .

Here is another property related to the graph only.

2.17 Definition. A state $x \in V$ is *transient*, if there exists $y \in V$ such that a path from x to y exists, but a path from y to x does not exist. Otherwise, x is called *recurrent*.

2.18 Exercise. If x is transient then

$$\mathbb{P}(S_n = x) \rightarrow 0 \quad \text{as } n \rightarrow \infty.$$

Prove it.

Hint: apply 2.3 to the set A of all y such that there is no path from y to x .

Recurrent states x, y are called equivalent, if there exists a path from x to y , and a path from y to x . (Well, the latter follows from the former.) Equivalence classes are irreducible closed sets. . .