**Homework 5** 

- 1. Prove that  $\sqrt{F}(\rho, \sum_i \lambda_i \sigma_i) \ge \sum_i \lambda_i \sqrt{F}(\rho, \sigma_i)$ .
  - Prove that if the classical fidelity function F (over probability distributions) is concave, then so does the quantum fidelity (over density matrices).
  - Prove that  $F(\rho, \sum_i \lambda_i \sigma_i) \ge \sum_i \lambda_i F(\rho, \sigma_i)$ .
- 2. Prove that for every constant  $\varepsilon > 0$ ,  $Q_{\varepsilon}(IP) = \Omega(n)$ .

In What follows you may us the following inequalities:

**Sub-additivity** :  $S(AB) \leq S(A) + S(B)$ , and,

**Araki-Lieb inequality** :  $S(AB) \ge S(A) - S(B)$ .

## A fact about the Holevo quantity :

Suppose X is a random variable with  $p_x = \Pr(X = x)$ . Let  $\rho$  be any quantum encoding of X and denote  $\rho_x = \rho(x)$ . The Holevo quantity is  $\chi(\rho) = S(\sum_x p_x \rho_x) - \sum_x p_x S(\rho_x)$ . You may use the fact that if  $\rho$  is over quantum registers  $A \otimes B$ , then

$$\chi(\mathrm{Tr}_B(\rho)) \le \chi(\rho).$$

- 3. Show an example where I(A:B) > S(B). Prove that  $I(A:B) \le 2S(B)$ .
- 4. Alice wants to communicate an arbitrary x ∈ {0,1}<sup>n</sup> to Bob. Alice and Bob communicate in rounds, in each round Alice (or Bob) applies a unitary transformation on his/her part and transmits a quibt to the other side, until at the end Bob measures his state and tries to infer x. The protocol is successful if for every x ∈ {0,1}<sup>n</sup> Bob succeeds with probability 1. Let n<sub>A</sub> / n<sub>B</sub>) be the number of messages sent by Alice / Bob respectively.
  - (a) Show that for every  $n_A$  and  $n_B$  such that  $n_A \ge \lceil n/2 \rceil$  and  $n_A + n_B \ge n$  there exists a successful protocol with parameters  $n_A, n_B$ .
  - (b) Show that in any successful protocol n<sub>A</sub> ≥ [n/2] and n<sub>A</sub> + n<sub>B</sub> ≥ n.
    Hint: Follow how S(ρ<sub>i</sub>) and χ(ρ<sub>i</sub>) change with i, where ρ<sub>i</sub> is Bob's density matrix at round i.
  - (c) Define a protocol as *p*-successful if for every  $x \in \{0, 1\}^n$  Bob succeeds with probability *p*. Prove a statement similar to item (b) for *p*-successful protocols.
- 5. Assume the same situation as Q4 except that now Alice and Bob may share unlimited preprepared entanglement. Prove that in any successful protocol  $n_A \ge \lceil n/2 \rceil$ .