## Topics in Extremal Combinatorics (0366.4996)- Fall '21

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Home Assignment 4

Due date: 04/12/22

## Please submit organized and well written solutions!

**Problem 1.** Suppose  $\mathcal{H} = (V, E)$  is a 3-uniform 6-regular hypergraph. Show that it is possible to split E into two sets  $E_1, E_2$  so that both edge sets cover V, that it  $\bigcup_{e \in E_1} e = \bigcup_{e \in E_2} e = V$ .

**Problem 2.** For an integer k, let  $T_k$  be the complete k-ary tree of depth k. Let  $\mathcal{H}_1$  be the set system whose ground set is the edges of  $T_k$  and has a set for each path of  $T_k$  starting at the root and ending at a leaf. Let  $\mathcal{H}_2$  be the set system whose ground set is the edges of  $T_k$  and has a set for each internal vertex  $v \in T_k$  that consists of the edges connecting v to its children. Show that  $\operatorname{disc}(\mathcal{H}_2) \leq 1$ ,  $\operatorname{disc}(\mathcal{H}_2) \leq 1$  and  $\operatorname{disc}(\mathcal{H}_1 \cup \mathcal{H}_2) \geq k$ .

**Problem 3.** Let A be an  $n \times n$  matrix of  $\pm 1$ . Show that there is  $x \in \{1, -1\}^n$  so that for every  $1 \le i \le n$ , the *i*<sup>th</sup> entry of Ax is smaller, in absolute value, than 2i.

**Problem 4.** Show that if  $S_1, \ldots, S_n$  are subsets of [r], then there is a  $\{-1, 0, 1\}$  valued function that induces discrepancy at most  $C\sqrt{r\log(2n/r)}$  and attains the value 0 at most 9r/10 times.

**Problem 5.** Show that the eigenvalue bound for the discrepancy of a set system gives only a  $\Theta(1)$  lower bound for the set system defined by a Hadamard matrix.

**Problem 6.** Given a hypergraph  $\mathcal{H}$  let  $\mathcal{H}_t$  be the sub-hypergraph containing only the edges of size at least t, and let  $\Delta(\mathcal{H})$  denote  $\mathcal{H}$ 's maximum degree. We proved in class that if  $\mathcal{H}$  is an n-vertex hypergraph with m edges and there is some t, so that  $\Delta(\mathcal{H}_t) \leq t$ , then  $\mathcal{H}$ 's discrepancy is  $O(\sqrt{t} \cdot \log n \cdot \sqrt{\log m})$ . Use this to prove that the discrepancy of arithmetic progressions within  $\{1, \ldots, n\}$  is  $O(n^{1/4} \log^C n)$ .

Hint: Consider only arithmetic progressions whose length is a power of 2 with a special structure.