# A construction of almost Steiner systems

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#### Abstract

Let n, k, and t be integers satisfying  $n > k > t \ge 2$ . A Steiner system with parameters t, k, and n is a k-uniform hypergraph on n vertices in which every set of t distinct vertices is contained in exactly one edge. An outstanding problem in Design Theory is to determine whether a nontrivial Steiner system exists for  $t \ge 6$ .

In this note we prove that for every  $k > t \ge 2$  and sufficiently large n, there exists an almost Steiner system with parameters t, k, and n; that is, there exists a k-uniform hypergraph on n vertices such that every set of t distinct vertices is covered by either one or two edges.

## 1 Introduction

Let n, k, t, and  $\lambda$  be positive integers satisfying  $n > k > t \ge 2$ . A t- $(n, k, \lambda)$ -design is a k-uniform hypergraph  $\mathcal{H} = (X, \mathcal{F})$  on n vertices with the following property: every t-set of vertices  $A \subset X$  is contained in exactly  $\lambda$  edges  $F \in \mathcal{F}$ . The special case  $\lambda = 1$  is known as a Steiner system with parameters t, k, and n, named after Jakob Steiner who pondered the existence of such systems in 1853. Steiner systems, t-designs<sup>1</sup> and other combinatorial designs turn out to be useful in a multitude of applications, e.g., in coding theory, storage systems design, and wireless communication. For a survey of the subject, the reader is referred to [2].

A counting argument shows that a Steiner triple system — that is, a 2-(n, 3, 1)-design — can only exist when  $n \equiv 1$  or  $n \equiv 3 \pmod{6}$ . For every such n, this is achieved via constructions based on symmetric idempotent quasigroups. Geometric constructions over finite fields give rise to some further infinite families of Steiner systems with t = 2 and t = 3. For instance, for a prime power q and an integer  $m \geq 2$ , affine geometries yield 2- $(q^m, q, 1)$ -designs, projective geometries yield 2- $(q^m + \dots + q^2 + q + 1, q + 1, 1)$ -designs and spherical geometries yield 3- $(q^m + 1, q, 1)$ -designs.

For t = 4 and t = 5, only finitely many nontrivial constructions of Steiner systems are known; for  $t \ge 6$ , no constructions are known at all.

Before stating our result, let us extend the definition of t-designs as follows. Let n, k, and t be positive integers satisfying  $n > k > t \ge 2$  and let  $\Lambda$  be a set of positive integers. A t- $(n, k, \Lambda)$ -design is a k-uniform hypergraph  $\mathcal{H} = (X, \mathcal{F})$  on n vertices with the following property: for every t-set of vertices  $A \subset X$ , the number of edges  $F \in \mathcal{F}$  that contain A belongs to  $\Lambda$ . Clearly, when  $\Lambda = \{\lambda\}$  is a singleton, a t- $(n, k, \{\lambda\})$ -design coincides with a t- $(n, k, \lambda)$ -design as defined above.

Not able to construct Steiner systems for large t, Erdős and Hanani [3] aimed for large partial Steiner systems; that is, t- $(n, k, \{0, 1\})$ -designs with as many edges as possible. Since a Steiner

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<sup>&</sup>lt;sup>1</sup>That is, t- $(n, k, \lambda)$ -designs for some parameters n, k, and  $\lambda$ .

system has exactly  $\binom{n}{t}/\binom{k}{t}$  edges, they conjectured the existence of partial Steiner systems with  $(1-o(1))\binom{n}{t}/\binom{k}{t}$  edges. This was first proved by Rödl [8] in 1985, with further refinements [4, 5, 6] of the o(1) term, as stated in the following theorem:

**Theorem 1** (Rödl). Let k and t be integers such that  $k > t \ge 2$ . Then there exists a partial Steiner system with parameters t, k, and n covering all but  $o(n^t)$  of the t-sets.

Theorem 1 can also be rephrased in terms of a covering rather than a packing; that is, it asserts the existence of a system with  $(1+o(1))\binom{n}{t}/\binom{k}{t}$  edges such that every t-set is covered at least once (see, e.g., [1, page 56]). Nevertheless, some t-sets might be covered multiple times (perhaps even  $\omega(1)$  times). It is therefore natural to ask for t- $(n, k, \{1, ..., r\})$ -designs, where r is as small as possible. The main aim of this short note is to show how to extend Theorem 1 to cover all t-sets at least once but at most twice.

**Theorem 2.** Let k and t be integers such that  $k > t \ge 2$ . Then, for sufficiently large n, there exists a t- $(n, k, \{1, 2\})$ -design.

Our proof actually gives a stronger result: there exists a t- $(n, k, \{1, 2\})$ -design with  $(1 + o(1)) \binom{n}{t} / \binom{k}{t}$  edges.

# 2 Preliminaries

In this section we present results needed for the proof of Theorem 2.

Given a t- $(n, k, \{0, 1\})$ -design  $\mathcal{H} = (X, \mathcal{F})$ , we define the leave hypergraph  $(X, \mathcal{L}_{\mathcal{H}})$  to be the t-uniform hypergraph whose edges are the t-sets  $A \subset X$  not covered by any edge  $F \in \mathcal{F}$ .

Following closely the proof of Theorem 1 appearing in [4], we recover an extended form of the theorem, which is a key ingredient in the proof of our main result.

**Theorem 3.** Let k and t be integers such that  $k > t \ge 2$ . There exists a constant  $\varepsilon = \varepsilon(k, t) > 0$  such that for sufficiently large n, there exists a partial Steiner system  $\mathcal{H} = (X, \mathcal{F})$  with parameters t, k, and n satisfying the following property:

(\*) For every  $0 \le \ell < t$ , every set  $X' \subset X$  of size  $|X'| = \ell$  is contained in  $O\left(n^{t-\ell-\varepsilon}\right)$  edges of the leave hypergraph.

We also make use of the following probabilistic tool.

**Talagrand's inequality.** In its general form, Talagrand's inequality is an isoperimetric-type inequality for product probability spaces. We use the following formulation from [7, pages 232–233], suitable for showing that a random variable in a product space is unlikely to overshoot its expectation under two conditions:

**Theorem 4** (Talagrand). Let  $Z \ge 0$  be a non-trivial random variable, which is determined by n independent trials  $T_1, \ldots, T_n$ . Let c > 0 and suppose that the following properties hold:

- i. (c-Lipschitz) changing the outcome of one trial can affect Z by at most c, and
- ii. (Certifiable) for any s, if  $Z \ge s$  then there is a set of at most s trials whose outcomes certify that  $Z \ge s$ .

Then  $\Pr[Z > t] < 2 \exp(-t/16c^2)$  for any  $t \ge 2\mathbb{E}[Z] + 80c\sqrt{\mathbb{E}[Z]}$ .

### 3 Proof of the main result

In this section we prove Theorem 2.

#### 3.1 Outline

The construction is done in two phases:

- I. Apply Theorem 3 to get a t- $(n, k, \{0, 1\})$ -design  $\mathcal{H} = (X, \mathcal{F})$  with property  $(\clubsuit)$  with respect to some  $0 < \varepsilon < 1$ .
- II. Build another t- $(n, k, \{0, 1\})$ -design  $\mathcal{H}' = (X, \mathcal{F}')$  that covers the uncovered t-sets  $\mathcal{L}_{\mathcal{H}}$ .

Combining both designs, we get that every t-set is covered at least once but no more than twice; namely  $(X, \mathcal{F} \cup \mathcal{F}')$  is a t- $(n, k, \{1, 2\})$ -design, as required.

We now describe how to build  $\mathcal{H}'$ . For a set  $A \subset X$ , denote by  $\mathcal{T}_A = \{C \subseteq X : |C| = k \text{ and } A \subseteq C\}$  the family of possible continuations of A to a subset of X of cardinality k. Note that  $\mathcal{T}_A = \emptyset$  when |A| > k.

Consider the leave hypergraph  $(X, \mathcal{L}_{\mathcal{H}})$ . Our goal is to choose, for every uncovered t-set  $A \in \mathcal{L}_{\mathcal{H}}$ , a k-set  $A' \in \mathcal{T}_A$  such that  $|A' \cap B'| < t$  for every two distinct  $A, B \in \mathcal{L}_{\mathcal{H}}$ . This ensures that the obtained hypergraph  $\mathcal{H}' = (X, \{A' : A \in \mathcal{L}_{\mathcal{H}}\})$  is indeed a t- $(n, k, \{0, 1\})$ -design.

To this aim, for every  $A \in \mathcal{L}_{\mathcal{H}}$  we introduce intermediate lists  $\mathcal{R}_A \subseteq \mathcal{S}_A \subseteq \mathcal{T}_A$  that will help us control the cardinalities of pairwise intersections when choosing  $A' \in \mathcal{R}_A$ . First note that we surely cannot afford to consider continuations that fully contain some other  $B \in \mathcal{L}_{\mathcal{H}}$ , so we restrict ourselves to the list

$$S_A = \mathcal{T}_A \setminus \bigcup \{\mathcal{T}_B : B \in \mathcal{L}_H \setminus \{A\}\} = \{C \in \mathcal{T}_A : B \not\subseteq C \text{ for all } B \in \mathcal{L}_H, B \neq A\}.$$

Note that, by definition, the lists  $\mathcal{S}_A$  for different A are disjoint. Next, choose a much smaller sub-list  $\mathcal{R}_A \subseteq \mathcal{S}_A$  by picking each  $C \in \mathcal{S}_A$  to  $\mathcal{R}_A$  independently at random with probability  $p = n^{t-k+\varepsilon/2}$  (we can of course assume here and later that  $\varepsilon < 1 \le k - t$ , and thus  $0 ). Finally, select <math>A' \in \mathcal{R}_A$  that has no intersection of size at least t with any  $C \in \mathcal{R}_B$  for any other  $B \in \mathcal{L}_H$ . If there is such a choice for every  $A \in \mathcal{L}_H$ , we get  $|A' \cap B'| < t$  for distinct  $A, B \in \mathcal{L}_H$ , as requested.

#### 3.2 Details

We start by showing that the lists  $S_A$  are large enough.

Claim 5. For every  $A \in \mathcal{L}_{\mathcal{H}}$  we have  $|\mathcal{S}_A| = \Theta(n^{k-t})$ .

*Proof.* Fix  $A \in \mathcal{L}_{\mathcal{H}}$ . Obviously  $|\mathcal{T}_A| = \binom{n-t}{k-t} = \Theta\left(n^{k-t}\right)$ . Since  $\mathcal{S}_A \subseteq \mathcal{T}_A$ , it suffices to show that  $|\mathcal{T}_A \setminus \mathcal{S}_A| = o\left(n^{k-t}\right)$ .

Writing  $\mathcal{L}_H \setminus \{A\}$  as the disjoint union  $\bigcup_{\ell=0}^{t-1} \mathcal{B}_{\ell}$ , where  $\mathcal{B}_{\ell} = \{B \in \mathcal{L}_H : |A \cap B| = \ell\}$ , we have

$$\mathcal{T}_A \setminus \mathcal{S}_A = \{ C \in \mathcal{T}_A : \exists B \in \mathcal{L}_H \setminus \{A\} \text{ such that } C \in \mathcal{T}_B \}$$

$$= \bigcup_{\ell=0}^{t-1} \{ C \in \mathcal{T}_A : \exists B \in \mathcal{B}_\ell \text{ such that } C \in \mathcal{T}_B \}$$

$$= \bigcup_{\ell=0}^{t-1} \bigcup_{B \in \mathcal{B}_\ell} \mathcal{T}_{A \cup B}.$$

Note that for all  $0 \le \ell < t$  and for all  $B \in \mathcal{B}_{\ell}$ ,  $|A \cup B| = 2t - \ell$  and thus  $|\mathcal{T}_{A \cup B}| = \binom{n-2t+\ell}{k-2t+\ell} \le n^{k-2t+\ell}$ . Moreover,  $|\mathcal{B}_{\ell}| = \binom{t}{\ell} \cdot O\left(n^{t-\ell-\varepsilon}\right) = O\left(n^{t-\ell-\varepsilon}\right)$  by Property (4). Thus,

$$|\mathcal{T}_A \setminus \mathcal{S}_A| \le \sum_{\ell=0}^{t-1} |\mathcal{B}_\ell| \, n^{k-2t+\ell} = \ell \cdot O\left(n^{k-t-\varepsilon}\right) = o\left(n^{k-t}\right),$$

establishing the claim.

Recall that the sub-list  $\mathcal{R}_A \subseteq \mathcal{S}_A$  was obtained by picking each  $C \in \mathcal{S}_A$  to  $\mathcal{R}_A$  independently at random with probability  $p = n^{t-k+\varepsilon/2}$ . The next claim shows that  $\mathcal{R}_A$  typically contains many k-sets whose pairwise intersections are exactly A. This will be used in the proof of Claim 7.

Claim 6. Almost surely (i.e., with probability tending to 1 as n tends to infinity), for every  $A \in \mathcal{L}_H$ , the family  $\mathcal{R}_A$  contains a subset  $\mathcal{Q}_A \subseteq \mathcal{R}_A$  of size  $\Theta\left(n^{\varepsilon/3}\right)$  such that  $C_1 \cap C_2 = A$  for every two distinct  $C_1, C_2 \in \mathcal{Q}_A$ .

*Proof.* Fix  $A \in \mathcal{L}_H$ . Construct  $\mathcal{Q}_A$  greedily as follows: start with  $\mathcal{Q}_A = \emptyset$ ; as long as  $|\mathcal{Q}_A| < n^{\varepsilon/3}$  and there exists  $C \in \mathcal{R}_A \setminus \mathcal{Q}_A$  such that  $C \cap C' = A$  for all  $C' \in \mathcal{Q}_A$ , add C to  $\mathcal{Q}_A$ . It suffices to show that this process continues  $|n^{\varepsilon/3}|$  steps.

If the process halts after  $s < \lfloor n^{\varepsilon/3} \rfloor$  steps, then every k-tuple of  $\mathcal{R}_A$  intersects one of the s previously chosen sets in some vertex outside A. This means that there exists a subset  $X_A \subset X$  of cardinality  $|X_A| = kn^{\varepsilon/3}$  ( $X_A$  contains the union of these s previously picked sets) such that none of the edges C of  $\mathcal{S}_A$  satisfying  $C \cap X_A = A$  is chosen into  $\mathcal{R}_A$ . For bounding the number of such edges in  $\mathcal{S}_A$  from bellow, we need to subtract from  $|\mathcal{S}_A|$  the number of edges  $C \in \mathcal{T}_A$  with  $C \cap X_A \neq A$ . The latter can be bounded (from above) by  $\sum_{i=1}^{k-t} {|X_A| \choose k-t-i}$  (choose  $1 \leq i \leq k-t$  vertices from  $X_A$ , other than A, to be in C, and then choose the remaining k-t-i vertices from  $X \setminus X_A$ ). Since  $|\mathcal{S}_A| = \Theta(n^{k-t})$ , and since

$$\sum_{i=1}^{k-t} {|X_A| \choose i} {n-|X_A| \choose k-t-i} \le \sum_{i=1}^{k-t} \Theta(n^{i\varepsilon/3}) \cdot n^{k-t-i} = o(n^{k-t}),$$

we obtain that the number of such edges in  $S_A$  is at least  $|S_A| - \sum_{i=1}^{k-t} {|X_A| \choose i} {n-|X_A| \choose k-t-i} = \Theta(n^{k-t})$ . It thus follows that the probability of the latter event to happen for a given A is at most

$$\binom{n}{kn^{\varepsilon/3}} \left(1-p\right)^{\Theta\left(n^{k-t}\right)}$$

(choose  $X_A$  first, and then require all edges of  $S_A$  intersecting  $X_A$  only at A to be absent from  $\mathcal{R}_A$ ). The above estimate is clearly at most

$$n^{n^{\varepsilon/3}} \cdot e^{-\Theta\left(pn^{k-t}\right)} = \exp\left\{n^{\varepsilon/3} \ln n - \Theta\left(n^{\varepsilon/2}\right)\right\} < \exp\left\{-n^{\varepsilon/3}\right\} \,.$$

Taking the union bound over all  $(\leq \binom{n}{t})$  choices of A establishes the claim.

The last step is to select a well-behaved set  $A' \in \mathcal{R}_A$ . The next claim shows this is indeed possible. Claim 7. Almost surely for every  $A \in \mathcal{L}_H$  we can select  $A' \in \mathcal{R}_A$  such that  $|A' \cap C| < t$  for all  $C \in \bigcup \{\mathcal{R}_B : B \in \mathcal{L}_H \setminus \{A\}\}$ .

Proof. Fix  $A \in \mathcal{L}_H$  and fix all the random choices which determine the list  $\mathcal{R}_A$  such that it satisfies Claim 6. Let  $\mathcal{Q}_A \subseteq \mathcal{R}_A$  be as provided by Claim 6 and let  $\mathcal{R} = \bigcup \{\mathcal{R}_B : B \in \mathcal{L}_H \setminus \{A\}\}$  be the random family of all obstacle sets. Define the random variable Z to be the number of sets  $A' \in \mathcal{Q}_A$  for which  $|A' \cap C| \geq t$  for some  $C \in \mathcal{R}$ . Since  $\mathcal{S}_A$  is disjoint from  $\mathcal{S} = \bigcup \{\mathcal{S}_B : B \in \mathcal{L}_H \setminus \{A\}\}$ , we can view  $\mathcal{R}$  as a random subset of  $\mathcal{S}$ , with each element selected to  $\mathcal{R}$  independently with probability  $p = n^{t-k+\varepsilon/2}$ . Thus Z is determined by  $|\mathcal{S}|$  independent trials. We wish to show that Z is not too large via Theorem 4; for this, Z has to satisfy the two conditions therein.

- 1. If  $C \in \mathcal{S}$  satisfies  $|A' \cap C| \geq t$  for some  $A' \in \mathcal{Q}_A$  then  $C \setminus A$  must intersect  $A' \setminus A$  (since  $A \not\subseteq C$ ). However, the (k-t)-sets  $\{A' \setminus A : A' \in \mathcal{Q}_A\}$  are pairwise disjoint (by the definition of  $\mathcal{Q}_A$ ) so each C cannot rule out more than k different sets  $A' \in \mathcal{Q}_A$ . Thus Z is k-Lipschitz.
- 2. Assume that  $Z \geq s$ . Then, by definition, there exist distinct sets  $A'_1, \ldots, A'_s \in \mathcal{Q}_A$  and (not necessarily distinct) sets  $C_1, \ldots, C_s \in \mathcal{R}$  such that  $|A'_i \cap C_i| \geq t$  for  $i = 1, \ldots, s$ . These are at most s trials whose outcomes ensure that  $Z \geq s$ ; i.e., Z is certifiable.

Let us now calculate  $\mathbb{E}[Z]$ . Fix  $A' \in \mathcal{Q}_A$  and let  $Z_{A'}$  be the indicator random variable of the event  $E_{A'} = \{\exists C \in \mathcal{R} : |A' \cap C| \geq t\}$ . The only set in  $\mathcal{L}_H$  fully contained in A' is A, so we can write  $\mathcal{L}_H \setminus \{A\}$  as the disjoint union  $\bigcup_{\ell=0}^{t-1} \mathcal{B}'_{\ell}$ , where  $\mathcal{B}'_{\ell} = \{B \in \mathcal{L}_H : |A' \cap B| = \ell\}$ . For any  $0 \leq \ell < t$  and  $B \in \mathcal{B}'_{\ell}$ , the number of bad sets (i.e., sets that will trigger  $E_{A'}$ ) in  $\mathcal{S}_B$  is

$$\begin{aligned} \left| \left\{ C \in \mathcal{S}_B : \left| A' \cap C \right| \ge t \right\} \right| &\le \left| \left\{ C \in \mathcal{T}_B : \left| A' \cap C \right| \ge t \right\} \right| \\ &= \left| \left\{ C \in \mathcal{T}_B : \left| \left( A' \cap C \right) \setminus B \right| \ge t - \ell \right\} \right| \\ &= \sum_{i=t-\ell}^{k-t} \binom{k-\ell}{i} \binom{n-k-t+\ell}{k-t-i} = O\left(n^{k-2t+\ell}\right), \end{aligned}$$

since C contains B, |B| = t, together with  $i \ge t - \ell$  elements from  $A' \setminus B$  and the rest from  $X \setminus (A' \cup B)$ . Each such bad set ends up in  $\mathcal{R}_B$  with probability  $p = n^{\varepsilon/2 + t - k}$ , so the expected number of bad sets in  $\mathcal{R}_B$  is  $O\left(n^{\varepsilon/2 - t + \ell}\right)$ . By Property ( $\clubsuit$ ) we have  $|\mathcal{B}'_{\ell}| = O\left(n^{t - \ell - \varepsilon}\right)$  and thus the total expected number of bad sets in  $\mathcal{R}$  is  $\ell \cdot O\left(n^{\varepsilon/2 - \varepsilon}\right) = O\left(n^{-\varepsilon/2}\right)$ . By Markov's inequality we have

$$\Pr\left[E_{A'}\right] = O\left(n^{-\varepsilon/2}\right).$$

Now, Z is a sum of  $|\mathcal{Q}_A| = \Theta(n^{\varepsilon/3})$  random variables  $Z_{A'}$  and thus

$$\mathbb{E}\left[Z\right] = |\mathcal{Q}_A| \cdot \mathbb{E}\left[Z_{A'}\right] = |\mathcal{Q}_A| \cdot \Pr\left[E_{A'}\right] = O\left(n^{-\varepsilon/6}\right) = o\left(1\right).$$

Applying Theorem 4 with  $t = |\mathcal{Q}_A| = \Theta(n^{\varepsilon/3})$  and c = k = O(1), we get that

$$\Pr\left[\text{all } A' \in \mathcal{R}_A \text{ are ruled out}\right] \leq \Pr\left[Z \geq |\mathcal{Q}_A|\right] < 2\exp\left(-\Omega\left(n^{\varepsilon/3}\right)\right).$$

Taking the union bound over all  $|\mathcal{L}_H| \leq \binom{n}{t}$  choices of  $A \in \mathcal{L}_H$  establishes the claim.

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