# A UNIFORM TAUBERIAN THEOREM IN DYNAMIC PROGRAMMING\*

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We prove that, in dynamic programming framework, uniform convergence of  $v_{\lambda}$  implies uniform convergence of  $v_n$  and vice versa. Moreover, both have the same limit.

1. Introduction. A deterministic dynamic programming problem is defined by a set of states S, a (nonvoid) correspondence  $\Gamma$  from S to itself and a bounded real function on S, say with values in [0,1].

Given the state s, one chooses t in  $\Gamma(s)$  and gets a payoff of f(s). A strategy is such a sequence of (history dependent) choices at each stage  $n=0,1,\ldots$ . Any strategy induces a play at s, i.e., a sequence  $h=(s=s_0,s_1,\ldots,s_n,\ldots)$  with  $s_{n+1}\in\Gamma(s_n)$ . Each play h induces an n-average payoff  $f_n(h)=(1/n)\sum_{m=0}^{n-1}f(s_m)$  and a  $\lambda$ -discounted payoff  $f_{\lambda}(h)=(1-\lambda)\sum_{m=0}^{\infty}\lambda^m f(s^m)$ . Taking the supremum on all strategies of the above functions defines the n-stage value  $v_n(s)$  and the  $\lambda$ -discounted value  $v_{\lambda}(s)$ .

We will consider here the asymptotic behavior of these two families of functions (as  $n \to \infty$  or  $\lambda \to 1$ ) and prove that the uniform convergence of one implies the uniform convergence of the other, and both to the same limit. Note that without the uniformity condition  $\lim v_n$  and  $\lim v_\lambda$  may exist and differ (see Example (§2)). Moreover, this condition does not imply the equality with  $v_\infty$  (defined through  $f_\infty(h) = \liminf f_n(h)$ ) (Lehrer and Monderer 1989).

The proofs and the result extend easily to the general (stochastic) case (see §6). Formally, let

(A)  $v_{\lambda} \rightarrow_{\lambda \rightarrow 1} v$ , uniformly on S and

(B)  $v_n \to \int_{n \to \infty}^{\infty} w$ , uniformly on S.

The purpose of this paper is to prove the following.

THEOREM. (a) If (A) then (B) and v = w; (b) If (B) then (A) and v = w.

**2. Example.** Take  $S = \mathbb{N}^* \times \mathbb{N}$ , where  $\mathbb{N}^* = \mathbb{N} \setminus \{0\}$ ,  $\Gamma(n, 0) = \{(n + 1, 0), (n, 1)\}$  and  $\Gamma(n, m) = \{(n, m + 1)\}$  for m > 0. f(n, m) = 1 iff  $1 \le m \le n$  and 0 otherwise. In words, at each state (n, 0) either you choose to get 1 for the next n stages and then always 0 or you proceed to state (n + 1, 0) and get 0 at that stage.

Let s = (1, 0). The feasible sequences of payoffs are of the form: n times 0, n times 1 and then only 0 (say, on a play  $h_n$  at s); or always 0. Obviously,  $\lim v_n(s) = \frac{1}{2}$ .  $f_{\lambda}(h_n) = \lambda^n - \lambda^{2n}$ ; hence  $\lim v_{\lambda}(s) = \frac{1}{4}$  and, finally,  $v_{\infty}(s) = 0$ .

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3. Preliminary results. Let us begin by proving properties that hold in a general framework. The first one shows that lim sup is decreasing on plays.

PROPOSITION 1. For any play  $h = (s_m)$  at s one has

$$\limsup v_n(s_m) \le \limsup v_n(s)$$
 and

$$\limsup v_{\lambda}(s_m) \leq \limsup v_{\lambda}(s) \quad \text{for all } m.$$

PROOF. Given m, choose  $N > 2m/\epsilon$  and a play at  $s_m$ , h', satisfying  $f_n(h') \ge \limsup v_n(s_m) - \epsilon/2$  with  $n \ge N$ . Then

$$f_{n+N}(s_0, s_1, \dots, s_{m-1}, h') \ge \limsup v_n(s_m) - \epsilon.$$

Similarly, in the discounted case, let  $\lambda_0$  with  $\lambda_0^m > (1 - \epsilon/2)$  and h'' at  $s_m$  satisfying  $f_{\lambda}(h'') \ge \limsup v_{\lambda}(s_m) - \epsilon/2$  with  $\lambda \ge \lambda_0$ . Then

$$f_{\lambda}(s_0,\ldots,s_{m-1},h'') \ge \limsup v_{\lambda}(s_m) - \epsilon.$$
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The second property is that, given an n average on a play, there exists a state (on this play) from which all averages of small length (compared to n) give at least this amount. Note that this result may be useful for related problems (cf. Lehrer and Monderer 1989).

PROPOSITION 2. Let  $\epsilon > 0$ . For all s, n there exist a play  $h = (s_l)$  at s and a stage L such that:

$$(1/T)\sum_{m=0}^{T-1} f(s_{L+m}) \ge v_n(s) - \epsilon \quad \text{for all } 1 \le T \le \lfloor n\epsilon/2 \rfloor.$$

PROOF. Otherwise there exist s and n such that on each play h at s and each stage L, there is some T = T(h, L) such that:

$$1 \le T \le \lfloor n\epsilon/2 \rfloor$$
 and  $(1/T) \sum_{m=0}^{T-1} f(s_{L+m}) < v_n(s) - \epsilon$ ,

in particular, if h is a play at s satisfying  $f_n(h) \ge v_n(s) - \epsilon/2$ . But then we can divide this play into segments of length at most  $[n\epsilon/2]$  on each of which, except at most the last one, the average payoff is less than  $v_n(s) - \epsilon$ . (Taking  $L_0 = 0$ , then  $L_1 = T(h, L_0)$ ,  $L_2 = T(h, L_1)$ , and so on.) It follows that  $f_n(h) \le v_n - \epsilon + \epsilon/2$ , a contradiction.

We now compare  $v_n$  and  $v_\lambda$ . Recall that any normalized power series with parameter in [0,1) can be written also as a convex combination of the finite averages. Since we will need it later, we provide here the explicit formula.

If  $\{a_m\}$  is a bounded sequence and  $0 \le \lambda < 1$ , then for all  $n \in \mathbb{N} \cup \{+\infty\}$ 

(1) 
$$(1 - \lambda) \sum_{m=0}^{n} a_m \lambda^m = (1 - \lambda)^2 \sum_{m=0}^{n-1} \lambda^m (m+1) \left( \sum_{l=0}^{m} a_l / (m+1) \right)$$
$$+ (1 - \lambda) \lambda^n (n+1) \left( \sum_{l=0}^{n} a_l / (n+1) \right).$$

This relation will allow us to examine the asymptotic behavior of families of geometric distributions. We start with the following simple observation.

LEMMA 3. Let  $M(\alpha, \beta; \lambda) = (1 - \lambda)^2 \sum_{\alpha}^{\beta} \lambda^m (m + 1)$ :

(i) There exist  $N_0$  and  $\epsilon_0$  such that  $\forall n \geq N_0, \forall \epsilon \leq \epsilon_0$ ,

(2) 
$$M([(1-\epsilon)n], n; 1-1/n) \ge \epsilon/2e.$$

(ii)  $\forall \delta > 0$ , there exist  $N_0$  and  $\epsilon_0$  such that  $\forall n \geq N_0, \forall \epsilon \leq \epsilon_0$ ,

(3) 
$$M([\epsilon n], [(1-\epsilon)n]; 1-1/n\sqrt{\epsilon}) \geq 1-\delta.$$

PROOF. Use (1) with  $n = +\infty$  to get

$$M(\alpha,\beta;\lambda) = ((\alpha+1)\lambda^{\alpha} - \alpha\lambda^{\alpha+1}) - ((\beta+2)\lambda^{\beta+1} - (\beta+1)\lambda^{\beta+2}).$$

For  $\alpha = [(1 - \epsilon)n]$ ,  $\beta = n$  and  $\lambda = 1 - 1/n$  the first term is of the order of  $(2 - \epsilon) \exp(-1 + \epsilon)$  and the second of 2/e. For  $\alpha = [\epsilon n]$ ,  $\beta = [(1 - \epsilon)n]$  and  $\lambda = 1 - 1/n\sqrt{\epsilon}$  we obtain, respectively,  $(1 + \sqrt{\epsilon}) \exp(-\sqrt{\epsilon})$  and  $1/\sqrt{\epsilon} \exp(-1/\sqrt{\epsilon})$ , providing that n is much larger than  $1/\epsilon$ .

PROPOSITION 4.  $\forall \epsilon > 0, \forall N$ , there is  $\lambda_0$  such that for all  $\lambda \geq \lambda_0$  and all s in S there exists  $n \geq N$  with  $v_n(s) \geq v_{\lambda}(s) - \epsilon$ .

PROOF. Given  $\epsilon > 0$  and N let  $\lambda_0$  be such that:

$$(1-\lambda)^2 \sum_{m=0}^{N-1} \lambda^m (m+1) < \epsilon/2 \quad \text{for } \lambda \ge \lambda_0.$$

By (1) this implies that on an  $\epsilon/2$  optimal play at s for  $v_{\lambda}$ , say h, there exists some  $n \ge N$  with  $f_n(h) \ge v_{\lambda}(s) - \epsilon$ . //

Corollary 5.  $\limsup v_n \ge \limsup v_\lambda$ .

## 4. Proof of part (a). We assume (A).

Lemma 6.  $\forall \epsilon > 0$ , there is an N such that  $n \geq N$  implies  $v_n \leq v + \epsilon$ .

PROOF. Otherwise, let  $\epsilon > 0$  such that for all N there exist  $n \ge N$  and s in S with  $v_n(s) > v(s) + \epsilon$ . Let  $\lambda$  such that  $||v_{\lambda} - v|| \le \epsilon/8$  (by (A)) and N such that

$$(1-\lambda)^2 \sum_{m=0}^{\lfloor n\epsilon/4\rfloor-1} \lambda^m (m+1) \ge 1-\epsilon/8 \quad \text{for } n \ge N.$$

We now use Proposition 2 with  $\epsilon/2$  to get a play h at s and a stage L with:

$$(1/T)\sum_{m=0}^{T-1} f(s_{L+m}) \ge v_n(s) - \epsilon/2 > v(s) + \epsilon/2 \quad \text{for all } 1 \le T \le \lfloor n\epsilon/4 \rfloor.$$

(1) then implies that

$$v_{\lambda}(s_L) \ge v(s) + \epsilon/2 - \epsilon/8.$$

Hence,

$$v(s_t) \ge v(s) + \epsilon/2 - \epsilon/8 - \epsilon/8$$

a contradiction to Proposition 1. //

LEMMA 7.  $\forall \epsilon > 0$ , there is an N such that  $n \geq N$  implies  $v_n \geq v - \epsilon$ .

PROOF. Otherwise, let  $\epsilon > 0$  such that for all N there exist  $n \ge N$  and s with  $v_n(s) < v(s) - \epsilon$ . In particular, for N large enough, on any play h at s one has:

$$(1/T)\sum_{m=0}^{T-1} f(s_m) \le v(s) - \epsilon/2 \quad \text{for } \left[ (1 - \epsilon/2)n \right] \le T \le n.$$

Choose N large enough and  $\epsilon$  small enough to get that the weight of these stages is at least  $\delta = \epsilon/4e$  for  $\lambda = 1 - 1/n$ ; i.e.,  $M([(1 - \epsilon/2)n], n; 1 - 1/n) \ge \epsilon/4e$  (by (2) in Lemma 3). Finally, let K be such that  $v_K(s) \le v(s) + \delta \epsilon/8$  for  $n \ge K$ , by Lemma 6. Choose N large enough to guarantee for  $n \ge N$  and  $\lambda = 1 - 1/n$ :  $||v_{\lambda} - v|| < \delta \epsilon/5$  (by (A)), and furthermore,  $(1 - \lambda)^2 \sum_{m=0}^{K-1} \lambda^m (m+1) < \epsilon \delta/8$ . This implies by (1) that:

$$f_{\lambda}(h) \leq \epsilon \delta/8 + \delta(v(s) - \epsilon/2) + (1 - \delta - \epsilon \delta/8)(v(s) + \delta \epsilon/8)$$
  
$$\leq v(s) - \delta \epsilon/4,$$

a contradiction. //

Lemmas 6 and 7 give (a).

REMARK. Notice that the previous proof shows also that uniform convergence of a sequence  $V_{\lambda_i}$ , where  $\lambda_i \to 1$ , implies uniform convergence of the sequence  $V_{n_i}$ , where  $n_i = [1/(1 - \lambda_i)]$ . Moreover, both converge to the same limit.

## 5. Proof of part (b). We assume (B).

LEMMA 8. For any  $\epsilon > 0$  small enough, there exists N such that for all  $n \geq N$  and all s, there is a play  $h = (s_t)$  at s satisfying:

$$(1/T)\sum_{m=0}^{T-1} f(s_m) \geqslant w(s) - \epsilon \quad \text{for all } [\epsilon n] \leq T \leq [(1-\epsilon)n].$$

PROOF. Use (B) to get N such that  $||v_n - w|| \le \delta$  for  $n \ge [\epsilon N]$  with  $\delta = \epsilon^2/3$ . Given  $n \ge N$  let  $h = (s_l)$  at s with  $f_n(h) \ge v_n(s) - \delta$ . For  $T \le [(1 - \epsilon)n]$  we obtain on h:

$$v_{n-T}(s_T) \le w(s_T) + \delta \le w(s) + \delta$$

by Proposition 1. Thus,

$$n(v_n(s)-\delta) \leq \sum_{m=0}^{T-1} f(s_m) + (n-T)(w(s)+\delta).$$

Hence,

$$(1/T)\sum_{m=0}^{T-1} f(s_m) \ge w(s) - n/T \cdot 3\delta$$

$$\ge w(s) - \epsilon \quad \text{for } T \ge [\epsilon n]. \quad //$$

Lemma 9.  $\forall \delta > 0$ , there exists  $\lambda_0$  such that  $\lambda \geq \lambda_0$  implies  $v_{\lambda} \geq w - \delta$ .

PROOF. Choose  $\epsilon_0$  and  $N_0$  as in (3) (Lemma 3) with  $\delta/3$ . Then use Lemma 8 to get with  $\epsilon \leq \delta/3$  for any *n* large enough and any *s*, the existence of a play *h* at *s* with:

$$f_{\lambda}(h) \ge (1 - \lambda)^2 \sum_{T = [\epsilon n]}^{[(1 - \epsilon)n]} \lambda^T (T + 1) \left( (1/T) \sum_{m=0}^{T-1} f(s_m) \right)$$
$$\ge (1 - \delta/3) (w - \delta/3) \quad \text{for } \lambda = \lambda_n = 1 - 1/n\sqrt{\epsilon}.$$

Note that  $\lambda_n \leq \lambda \leq \lambda_{n+1}$  implies

$$f_{\lambda_n}(h)/(1-\lambda_n) \le f_{\lambda}(h)/(1-\lambda) \le f_{\lambda_{n+1}}(h)/(1-\lambda_{n+1}).$$

Hence, the result for  $\lambda$  large enough. //

LEMMA 10.  $\forall \delta > 0$ , there exists  $\lambda_0$  such that  $\lambda \geq \lambda_0$  implies  $v_{\lambda} \leq w + \delta$ .

PROOF. Follows from Proposition 4, using (B). //

**6.** Comments. The result extends to the stochastic case as follows. Consider first a countable-Borel framework where S is a countable set of states, A is a Borel set of actions, q is a transition probability from  $S \times A$  to S and f is a measurable bounded payoff function from  $S \times A$  to  $\mathbb{R}$ . (Recall that if S and A are finite, v, w and  $v_{\infty}$  exist, are equal and realized with a pure stationary strategy (Blackwell 1962).)

Define, for any Markov strategy  $\sigma$  (Blackwell 1965), a play starting from s by a sequence  $\{w_n\}$  of probabilities on S with  $w_0 = \delta_s$ ,  $w_{n+1}(S') = \int_S q(S' \mid t, \sigma_n(t)) w_n(dt)$ , for all Borel sets  $S' \subset S$ . The corresponding sequence of payoffs is  $\{x_n\}$  with

$$x_n = \int_{S} f(t, \sigma_n(t)) w_n(dt).$$

The proof goes then word-for-word.

If one leaves the countable state set up, a selection theorem is needed in Proposition 1. Hence, one can use an analytic framework (Blackwell, Freedman and Orkin 1974), where plays are defined by strategies.

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