

The Longest Cycle of a Graph with a Large Minimal Degree

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ABSTRACT

We show that every graph G on n vertices with minimal degree at least n/k contains a cycle of length at least $\lfloor n/(k-1) \rfloor$. This verifies a conjecture of Katchalski. When $k=2$ our result reduces to the classical theorem of Dirac that asserts that if all degrees are at least $\frac{1}{2}n$ then G is Hamiltonian.

1. INTRODUCTION

For a graph $G = \langle V(G), E(G) \rangle$ let $\delta(G)$ denote the minimal degree of vertex of G and let $c(G)$ denote the circumference, i.e., the size of a longest cycle of G . In this note we prove

Theorem 1.1. Suppose $n > k \geq 2$ are integers and that G is a graph on n vertices with $\delta(G) \geq n/k$. Then $c(G) \geq \lfloor n/(k-1) \rfloor$.

This result was conjectured by M. Katchalski [3], who also proved it for $k \leq 4$. For $k=2$ it reduces to the classical theorem of Dirac ([2]; see also [1, p. 135]) and asserts that G is Hamiltonian provided $\delta(G) \geq \frac{1}{2}n$.

The following example shows that Theorem 1.1 is almost sharp: Suppose $n > k \geq 2$ and put $r = \lfloor (n-1)/k \rfloor$. Then $n = kr + s$, where $1 \leq s \leq k$ and r is the largest integer $< n/k$. Let G be the graph consisting of k complete graphs on $r+1$ vertices, $k-s+1$ of which share a common vertex. Then G has n vertices, $\delta(G) = r$, and G contains no cycle of size $> r+1$. It is worth noting, though, that there is a possible (slight) strengthening of Theorem 1.1, namely, $c(G) \geq \lfloor n/(k-1) \rfloor$ (without the integer part), provided $\delta(G) \geq n/k$. At the moment we can prove this strengthened version only for $n \gg k$.

To prove Theorem 1.1 we need several known results and several new lemmas. These are presented in Section 2 and are applied in Section 3 to deduce the theorem.

2. LEMMAS

Lemma 2.1 (Dirac [2], see also [4, problems 10.21, 10.27, pp. 67–68]). Let G be a graph on n vertices.

- (a) If $\delta(G) > 1$ then $c(G) \geq \delta(G) + 1$.
- (b) If $\delta(G) \geq \frac{1}{2}n > 1$ then G is Hamiltonian.
- (c) If G is 2-connected then

$$c(G) \geq \min(n, 2\delta(G)).$$

Lemma 2.2 (Lovász [4, problem 10.19, p. 67]). Let x and y be two distinct vertices of a 2-connected graph G and suppose that each vertex of G other than x and y has degree at least d . Then there is a path of length $\geq d$ between x and y . ■

Lemma 2.3. Let z be a vertex of a 2-connected graph G and suppose that each vertex of G other than z has degree $\geq d$. Then either G contains a cycle of size $\geq 2d$ including z or any two distinct vertices of G are connected by a path of length $\geq d - 1$.

Proof. If $G - z$ is 2-connected the second possibility holds, by Lemma 2.2. Otherwise, let B and C be two endblocks of $G - z$ and let $b \in V(B)$ and $c \in V(C)$ be the corresponding cutvertices (possibly $b = c$). Clearly z has a neighbor $b' \in V(B) - b$ and a neighbor $c' \in V(C) - c$. By Lemma 2.2, B contains a (b', b) -path of length $\geq d - 1$ and C contains a (c', c) -path of length $\geq d - 1$. These paths, the edges zb' , zc' , and a (b, c) -path in $G - [z + (B - b) + (C - c)]$ form a cycle of length $\geq 2d$ containing z . ■

Remark. As noted by one of the referees, one can prove that, in fact, under the hypotheses of Lemma 2.3, there is a path of length at least d between any two vertices of G . For our purposes, however, Lemma 2.3 suffices.

To prove Theorem 1.1, we have to study the block structure of the graph G . This is done in the rest of this section.

Suppose $n > k(k - 1)$, where $k \geq 2$, and let G be a connected graph on n vertices with $\delta(G) \geq n/k$. Call a block of G *large* if it has at least $n/k + 1$ vertices, otherwise call it *small*. Let \mathcal{B} denote the set of all large blocks of G , and \mathcal{C} the set of all cutvertices of G that belong to at least two large blocks. Finally, let H denote the bipartite graph with classes of vertices \mathcal{B} and \mathcal{C} , in which $B \in \mathcal{B}$ is joined to $c \in \mathcal{C}$ iff $c \in V(B)$. Note that H is a subgraph of the block-cutvertex tree of G induced by $\mathcal{B} \cup \mathcal{C}$, and is therefore a forest.

Lemma 2.4

- (a) $|\mathcal{B}| \leq k - 1$.
- (b) $|\mathcal{C}| \leq k - 2$.
- (c) Every vertex of G belongs to at least one large block of G .
- (d) If y is a cutvertex of G that belongs to a unique large block, then y is joined to at least $n/k - k + 2 + |\mathcal{C}|$ vertices of this block.

Proof. Note that every endblock of G is a large block. Indeed, if B is an endblock it contains a vertex v which is not a cutvertex of G and all the $\geq n/k$ neighbors of v belong to $V(B) - v$. Hence B is large. For $c \in \mathcal{C}$ let $d(c)$ denote the degree of c in H , i.e., the number of large blocks of G that contain c . Let $\gamma = \gamma(H)$ denote the number of connected components of H . Since H is a forest with classes of vertices \mathcal{B} and \mathcal{C} ;

$$|E(H)| = \sum\{d(c) : c \in \mathcal{C}\} = |\mathcal{B}| + |\mathcal{C}| - \gamma. \tag{2.1}$$

Thus

$$\begin{aligned} n &\geq |\cup\{V(B) : B \in \mathcal{B}\}| = \sum\{|V(B)| : B \in \mathcal{B}\} - \sum\{[d(c) - 1] : c \in \mathcal{C}\} \\ &\geq |\mathcal{B}|(n/k + 1) + |\mathcal{C}| - \sum\{d(c) : c \in \mathcal{C}\} = |\mathcal{B}|(n/k) + \gamma \\ &> |\mathcal{B}|n/k. \end{aligned}$$

This verifies (a).

Equality (2.1) and (a) imply

$$|\mathcal{C}| \leq \sum\{[d(c) - 1] : c \in \mathcal{C}\} = |\mathcal{B}| - \gamma \leq k - 2,$$

which proves (b).

To prove (c) assume it is false and let $v \in V(G)$ be a counterexample. Note that every vertex that belongs to a small block of G is a cutvertex, since its degree is at least n/k . Thus v is a cutvertex. Let B_1, B_2, \dots, B_r be the blocks of G that contain v . Clearly $\sum_{i=1}^r |V(B_i) - v| \geq n/k$ and every vertex of $\cup_{i=1}^r V(B_i)$ is a cutvertex. Thus there are at least $n/k > k - 1$ vertices of H within a distance 2 (in H) from v . Therefore H has more than $k - 1$ endvertices, each of which is a large block of G , contradicting (a). This contradiction proves (c).

To prove (d), let y be a cutvertex of G that belongs to a unique large block of G . Consider the subgraph F obtained from G by deleting all edges that do not belong to large blocks. By (a) and (c) F has at most $k - 1 - r$ connected components, where $r = |\mathcal{C}|$ and y belongs to one of them. Clearly y has at most one neighbor in G in any other component, whereas in its own component it is joined in G only to vertices of its own large block. Thus the degree of y in this

block is at least $n/k - (k - 2 - r) = n/k - k + 2 + r$. This completes the proof of the lemma. ■

3. PROOF OF THEOREM 1.1

Suppose $n > k \geq 2$ and let G be a graph on n vertices with $\delta(G) \geq n/k$. We must show that

$$c(G) \geq \lceil n/(k - 1) \rceil. \tag{3.1}$$

If $\lceil n/(k - 1) \rceil \leq \lceil n/k \rceil + 1$ this follows from Lemma 2.1(a). Thus we may assume

$$n/(k - 1) \geq \lceil n/(k - 1) \rceil \geq \lceil n/k \rceil + 2 \geq n/k + 2,$$

i.e.,

$$n \geq 2k(k - 1). \tag{3.2}$$

Clearly we may also assume that G is connected; otherwise add bridges to make it connected. By Lemma 2.4 G has at least $n - k + 2$ vertices that belong to exactly one large block and it has at most $k - 1$ large blocks. Thus there is a block B containing $m \geq (n - k + 2)/(k - 1) > n/(k - 1) - 1$ such vertices. Let D be the induced subgraph of B on these $m (\geq \lceil n/(k - 1) \rceil)$ vertices. Since D is obtained from B by deleting vertices that belong to at least two large blocks of G , Lemma 2.4(b), (d) implies that the degree of every vertex of D is at least $n/k - k + 2$. If $2(n/k - k + 2) \geq m$, D is Hamiltonian, by Lemma 2.1(b), and since $m \geq \lceil n/(k - 1) \rceil$ the desired result follows. Thus we may assume that $m > 2(n/k - k + 2)$. Similarly, if D is 2-connected the desired result follows from Lemma 2.1(c) and inequality (3.2), which guarantees that $2(n/k - k + 2) \geq n/(k - 1)$. Thus we may assume that D is not 2-connected. We consider three possible cases, according to the block structure of D .

Case 1. Each block of D is a connected component of it. Let A and C be two distinct blocks of D . Since B is 2-connected there are two vertex-disjoint paths in B connecting $a_i \in V(A)$ to $c_i \in V(C) (i = 1, 2)$ and using no edges of A or C . By Lemma 2.2, A contains an (a_1, a_2) -path of length at least $n/k - k + 2$ and C contains a (c_1, c_2) -path of length at least $n/k - k + 2$. Altogether we have a cycle of size at least $2(n/k - k + 2) + 2 > n/(k - 1)$, as needed.

Case 2. D has two nonadjacent endblocks A and C in the same connected component. As in the previous case, the 2-connectivity of B implies the existence of two vertex-disjoint paths in B connecting $a_i \in V(A)$ to $c_i \in V(C) (i = 1, 2)$

and using no edges of A or C . If A or C contains a cycle of size at least $2(n/k - k + 2) \geq n/(k - 1)$, (3.1) follows. Otherwise, by Lemma 2.3, A contains an (a_1, a_2) -path of length at least $n/k - k + 1$ and C contains a (c_1, c_2) -path of length at least $n/k - k + 1$. This gives a cycle of size at least $2 + 2(n/k - k + 1) \geq n/(k - 1)$, implying (3.1).

Case 3. D contains two adjacent endblocks A, C . Let x be the unique common vertex of A and C . By the 2-connectivity of B , B contains a path between some $a \in A, a \neq x$, and some $c \in C, c \neq x$, using no edges of A or C . By Lemma 2.2 A contains an (a, x) -path of length at least $n/k - k + 2$ and C contains a (c, x) -path of length at least $n/k - k + 2$. These paths, together with the previous (a, c) -path, form a cycle of size at least $2(n/k - k + 2) + 1 > n/(k - 1)$. This completes the proof of the theorem. ■

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