

# Economical coverings of sets of lattice points

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## Abstract

Let  $t(n, d)$  be the minimum number  $t$  such that there are  $t$  of the  $n^d$  lattice points

$$\{(x_1, \dots, x_d) : 1 \leq x_i \leq n\}$$

so that the  $\binom{t}{2}$  lines that they determine cover all the above  $n^d$  lattice points. We prove that for every integer  $d \geq 2$  there are two positive constants  $c_1 = c_1(d)$  and  $c_2 = c_2(d)$  such that for every  $n$

$$c_1 n^{d(d-1)/(2d-1)} \leq t(n, d) \leq c_2 n^{d(d-1)/(2d-1)} \log n.$$

The special case  $d = 2$  settles a problem of Erdős and Purdy.

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# 1 Introduction

We say that a set of points  $S$  in an Euclidean space (of any dimension) *determines* a line  $l$  if  $l$  contains at least 2 points of  $S$ . Let  $L = L(n, 2) = \{(x_1, x_2) : 1 \leq x_1, x_2 \leq n\}$  denote the set of all points in the  $n$  by  $n$  square lattice. A subset  $S$  of  $L$  is called an  $(n, 2)$ -*covering set* if the union of all the lines determined by  $S$  contains all the points in  $L$ . Let  $t(n, 2)$  denote the minimum cardinality of an  $(n, 2)$ -covering set.

Erdős and Purdy (see [1]) raised the problem of estimating  $t(n, 2)$ . They mentioned that it is not hard to see that  $t(n, 2) \geq \Omega(n^{2/3})$  and asked if  $t(n, 2) = o(n)$ . In this short paper we answer this question affirmatively and show that  $t(n, 2) \leq O(n^{2/3} \log n)$ . Therefore,  $t(n, 2)$  is indeed  $o(n)$  and the gap between the upper and the lower bounds for this quantity is a  $\log n$  factor.

Our proof combines simple geometric and probabilistic arguments with some of the standard techniques used in the study of Diophantine Approximations. The method works in dimensions bigger than 2 as well. Let  $L(n, d)$  denote the set of all  $n^d$  vectors  $(x_1, \dots, x_d)$ , where  $1 \leq x_i \leq n$  are integers for all  $i$ . A subset  $S$  of  $L(n, d)$  is an  $(n, d)$ -*covering set* if the lines determined by  $S$  cover all points of  $L(n, d)$ . Let  $t(n, d)$  denote the minimum possible cardinality of an  $(n, d)$ -covering set. Our main result is the following.

**Theorem 1.1** *For every integer  $d \geq 2$  there are two positive constants  $c_1 = c_1(d)$  and  $c_2 = c_2(d)$  such that for every  $n$ :*

$$c_1 n^{d(d-1)/(2d-1)} \leq t(n, d) \leq c_2 n^{d(d-1)/(2d-1)} \log n.$$

## 2 The lower bound

In this section we prove the easy part of Theorem 1.1, by showing that for every  $d$  there exists a positive constant  $c_1 = c_1(d)$  such that  $t(n, d)$  is at least  $c_1 n^{d(d-1)/(2d-1)}$  for all  $n$ . Put  $L = L(n, d)$ . For a line  $l$  in  $R^d$ , let  $|l|$  denote the number of points of  $L$  contained in  $l$ . Observe that if  $|l| \geq 2$  then  $l$  can be presented by its parametric equations  $X_i = a_i + p_i z$ , ( $1 \leq i \leq d$ ), where  $a_i, p_i$  are integers and the greatest common divisor of  $p_1, \dots, p_d$  is 1. We say that a line  $l$  with the above presentation is of *type  $q$*  if  $q = \max |p_i|$ . The *direction* of  $l$  is the vector  $(p_1, \dots, p_d)$ . Note that this

vector and its inverse  $(-p_1, \dots, -p_d)$  represent the same direction. We need the following simple fact.

**Fact:** For each  $q \geq 1$ , the number of distinct directions of lines of type  $q$  is at most  $d(2q+1)^{d-1} \leq d(3q)^{d-1}$ .

**Proof** There are  $d$  possibilities for choosing an index  $i$  such that  $|p_i| = q$ , and since a vector and its inverse define the same direction we may assume that  $p_i = q$ . Any other  $p_j$  is an integer between  $-q$  and  $q$  implying the desired estimate.  $\square$

Let  $S$  be a subset of cardinality  $t$  of  $L$ . Let  $F$  denote the set of all lines determined by  $S$ . Clearly,  $|F| \leq \binom{t}{2}$ . Let  $P(S)$  denote the set of all the points in the lattice  $L$  which are covered by the union of all the lines in  $F$ . Obviously,

$$|P(S)| \leq \sum_{l \in F} |l|. \quad (1)$$

Clearly, for each line  $l$  of type  $q$ ,  $|l| \leq n/q$ . Also, the number of lines in  $F$  in each direction is at most  $t/2$ , since a set of  $t$  points cannot determine more than  $t/2$  parallel lines. Let  $f_q$  denote the number of lines of type  $q$  in  $F$ . By the Fact and the remark above  $f_q \leq \frac{t}{2}d(3q)^{d-1}$ . Therefore, the right hand side of (1) is at most the maximum possible value of  $\sum_{q \geq 1} f_q \frac{n}{q}$ , subject to the constraints  $\sum_{q \geq 1} f_q \leq \binom{t}{2}$  and  $f_q \leq \frac{t}{2}d(3q)^{d-1}$ . This last maximum is obviously attained when  $\sum_{q \geq 1} f_q = \binom{t}{2}$  and  $f_q$  is as large as it can be for all  $q < s$  and is 0 for all  $q > s$ , where  $s$  is an appropriately chosen integer. Therefore, at the maximum,  $f_q = \frac{t}{2}d(3q)^{d-1}$  for all  $q < s$  and  $f_q = 0$  for all  $q > s$ . Since  $\sum_{q \geq 1} f_q = \binom{t}{2}$ , a simple calculation shows that  $s \leq 1 + t^{1/d}/3^{(d-1)/d}$ . This implies that the above maximum is at most  $\sum_{q=1}^s \frac{t}{2}d(3q)^{d-1} \frac{n}{q} < c_3(d)nt^{(2d-1)/d}$ . We have thus proved that if  $|S| = t$  then  $|P(S)| \leq c_3(d)nt^{(2d-1)/d}$ . Since if  $S$  is an  $(n, d)$ -covering set then  $P(S) = L$  (and hence  $|P(S)| = n^d$ ), this gives the following:

**Lemma 2.1** *If  $S$  is a subset of cardinality  $t$  of  $L(n, d)$  then the number of points of  $L(n, d)$  covered by the lines determined by  $S$  is at most  $c_3nt^{(2d-1)/d}$ , where  $c_3$  is a positive constant depending only on  $d$ . Thus, for every  $d$  there is a positive constant  $c_1 = c_1(d)$  such that  $t(n, d) \geq c_1n^{d(d-1)/(2d-1)}$  for all  $n$ .  $\square$*

### 3 The upper bound

In this section we prove the upper bound in Theorem 1.1. For convenience we omit the floor and ceiling signs and assume that the fractional powers of  $n$  appearing here are all integers. Since we deal with fixed values of  $d$  and large values of  $n$ , it is easy to see that this can indeed be assumed without loss of generality. We start with the following somewhat technical lemma.

**Lemma 3.1** *For every two lattice points  $(x_1, \dots, x_d)$  and  $(a_1, \dots, a_d)$  in  $L(n, d)$  there are  $d$  integers  $p_1, \dots, p_d$  and a real number  $z$ , such that for  $q = \max\{|p_i| : 1 \leq i \leq d\}$ :*

$$1 \leq q \leq n^{(d-1)/(2d-1)}, \quad (2)$$

$$|(x_i - a_i) - p_i z| \leq \frac{n^{(2d-2)/(2d-1)}}{q} \quad \text{for all } 1 \leq i \leq d, \quad (3)$$

and there exists an index  $j$ ,  $1 \leq j \leq d$ , such that

$$p_j = q \quad \text{and} \quad (x_j - a_j) - p_j z = 0. \quad (4)$$

**Proof** The proof is based on the standard argument of Dirichlet used in the study of approximation of reals by rationals, (see, e.g., [2]). Let us change the indices, if needed, so that

$$|x_d - a_d| = \max\{|x_i - a_i| : 1 \leq i \leq d\}.$$

If this maximum is zero the result is trivial (since in this case we can take  $p_i = 1$  for all  $i$  and  $z = 0$ ). Otherwise put  $Q = n^{1/(2d-1)}$  and define  $\alpha_i = \frac{x_i - a_i}{x_d - a_d}$  for  $1 \leq i \leq d-1$ . Consider the  $(d-1)$ -dimensional unit cube  $0 \leq y_i < 1$ , ( $1 \leq i \leq d-1$ ) and split it into  $Q^{d-1}$  identical subcubes by drawing hyperplanes parallel to its facets, where the distance between each pair of consecutive parallel hyperplanes is  $1/Q = n^{-1/(2d-1)}$ . For each integer  $j$ ,  $0 \leq j \leq Q^{d-1}$ , let  $P_j$  be the point  $(j\alpha_1 \pmod{1}, \dots, j\alpha_{d-1} \pmod{1})$  in the above unit cube. Since there are  $Q^{d-1} + 1$  such points, there are two of them, say  $P_l$  and  $P_m$ , where  $l < m$ , that lie in the same subcube. Define  $q = m - l$  and define  $p_i$  by the equality  $q\alpha_i = p_i + \epsilon_i$ , where  $p_i$  is an integer and  $|\epsilon_i| \leq 1/Q$ .

Define  $p_d = q$  and  $z = (x_d - a_d)/q$ . Since the absolute value of each  $\alpha_i$  is at most 1 it follows that  $|p_i| \leq q \leq Q^{d-1} = n^{(d-1)/(2d-1)}$  for all  $1 \leq i \leq d$ . In addition, for each  $1 \leq i \leq d$

$$\left| q \frac{x_i - a_i}{x_d - a_d} - p_i \right| \leq \frac{1}{Q},$$

implying

$$|(x_i - a_i) - p_i z| = |(x_i - a_i) - p_i \frac{x_d - a_d}{q}| \leq \frac{|x_d - a_d|}{qQ} \leq \frac{n^{(2d-2)/(2d-1)}}{q}.$$

Moreover, for  $i = d$  we have  $(x_d - a_d) - p_d z = 0$ . This completes the proof of the lemma.  $\square$

Using the last lemma we prove the following:

**Lemma 3.2** *For every fixed  $d \geq 2$  there are two positive constants  $c_4 = c_4(d)$  and  $c_5 = c_5(d)$  such that the following holds. For every  $n \geq c_5$  and for every lattice point  $(x_1, \dots, x_d) \in L(n, d)$ , there are at least  $\frac{1}{2}n^d$  lattice points  $(a_1, \dots, a_d)$  in  $L(n, d)$  satisfying*

$$(c_4 + 2)n^{(d-1)/(2d-1)} \leq a_i \leq n - (c_4 + 2)n^{(d-1)/(2d-1)} \quad \text{for all } 1 \leq i \leq d \quad (5)$$

such that there exist  $d$  integers  $p_1, \dots, p_d$  and a real number  $z$  so that

$$1 \leq |p_i| \leq n^{(d-1)/(2d-1)}, \quad (6)$$

and

$$|(x_i - a_i) - p_i z| \leq c_4 n^{(d-1)/(2d-1)} \quad \text{for all } 1 \leq i \leq d. \quad (7)$$

**Proof** We prove the lemma with  $c_4 = 4d9^{d-1}$ , where  $c_5$  is chosen so that for  $n \geq c_5$  the number of lattice points  $(a_1, \dots, a_d) \in L(n, d)$  violating (5) is at most  $\frac{1}{4}n^d$ .

By Lemma 3.1, for every  $(a_1, \dots, a_d)$  in  $L(n, d)$  there are integers  $p_1, \dots, p_d$  and a real number  $z$  such that if  $q$  is the maximum absolute value of the integers  $p_i$ , then inequalities (2) and (3) hold. Therefore, inequality (6) holds, by (2). If  $q \geq \frac{1}{c_4}n^{(d-1)/(2d-1)}$  then inequality (7) follows from (3). Therefore, by the choice of  $c_5$ , in order to complete the proof of the lemma it suffices to show that the number of points  $(a_1, \dots, a_d)$  for which the above  $q$  is smaller than  $\frac{1}{c_4}n^{(d-1)/(2d-1)}$  is at most  $\frac{1}{4}n^d$ .

The number of these points can be estimated as follows: For each fixed value of  $q$ ,  $1 \leq q \leq \frac{1}{c_4}n^{(d-1)/(2d-1)}$ , there are  $d$  possibilities for choosing the index  $j$  for which (4) holds, and there are at most  $n$  possibilities for choosing  $a_j$ . Once these choices are made,  $z$  is determined. There are at most  $(2q + 1)^{d-1} \leq (3q)^{d-1}$  choices for the other integers  $p_j$ . Given these choices, each of the numbers  $x_i - a_i$  must fall into an interval of length  $2n^{(2d-2)/(2d-1)}/q$ . Obviously, this implies that there are at most  $2n^{(2d-2)/(2d-1)}/q + 1 \leq 3n^{(2d-2)/(2d-1)}/q$  ways to choose each  $a_i$ , for all  $i \neq j$ .

The total number of choices is thus at most

$$\begin{aligned} & \sum_{q=1}^{n^{(d-1)/(2d-1)}/c_4} dn(3q)^{d-1} \left(3 \frac{n^{(2d-2)/(2d-1)}}{q}\right)^{d-1} \\ & \sum_{q=1}^{n^{(d-1)/(2d-1)}/c_4} 9^{d-1} dn^{1+\frac{(2d-2)(d-1)}{2d-1}} \\ & = \frac{9^{d-1}d}{c_4} n^d = \frac{1}{4}n^d. \end{aligned}$$

This completes the proof.  $\square$

For a lattice point  $\mathbf{a} = (a_1, \dots, a_d)$  that satisfies (5), let  $B(\mathbf{a})$  denote the set of all points  $(y_1, \dots, y_d) \in L(n, d)$  that satisfy

$$|a_i - y_i| \leq (c_4 + 2)n^{(d-1)/(2d-1)} \quad \text{for all } 1 \leq i \leq d,$$

where  $c_4 = c_4(d)$  is the constant appearing in Lemma 3.2. The last ingredient we need for the proof of the upper bound in Theorem 1.1 is the following simple lemma.

**Lemma 3.3** *Let  $(x_1, \dots, x_d)$  be a lattice point in  $L(n, d)$ , and suppose that the lattice point  $\mathbf{a} = (a_1, \dots, a_d)$  satisfies (5). Suppose, further, that there are  $d$  integers  $p_1, \dots, p_d$  and a real number  $z$  such that (6) and (7) hold. Then, there is a line containing  $(x_1, \dots, x_d)$  that contains at least two points of  $B(\mathbf{a})$ .*

**Proof** Let  $B$  denote the convex hull of the points in  $B(\mathbf{a})$ . Observe that  $B$  is a  $d$ -dimensional cube with side length  $(2c_4 + 4)n^{(d-1)/(2d-1)}$ . Consider the line  $l$  given by

$$X_i = x_i - p_i w, \quad (1 \leq i \leq d, -\infty < w < \infty).$$

This line contains the point  $(x_1, \dots, x_d)$  (corresponding to  $w = 0$ ). Let  $(v_1, \dots, v_d)$  be the point corresponding to  $w = z$  in this line, (i.e.,  $v_i = x_i - p_i z$  for  $1 \leq i \leq d$ ). By (7) this point lies inside the cube  $B$ . Moreover, its distance from each of the facets of  $B$  is at least  $2n^{(d-1)/(2d-1)}$ . Since every integral value of the parameter  $w$  gives a lattice point on  $l$ , this, together with the fact that the integers  $p_i$  satisfy (6) imply that the lattice points on  $l$  corresponding to, say,  $w = \lceil z \rceil$  and to  $w = \lceil z \rceil + 1$  both belong to  $B(\mathbf{a})$ , completing the proof.  $\square$

**Proof of Theorem 1.1** The lower bound is proved in Lemma 2.1. To prove the upper bound, fix  $d \geq 2$  and let  $c_4 = c_4(d)$  and  $c_5 = c_5(d)$  be as in Lemma 3.2. Suppose  $n \geq c_5$  and choose randomly

and independently (with repetitions),  $\lceil d \log n \rceil + 1$  lattice points  $(a_1, \dots, a_d)$  satisfying (5), where each point is chosen independently according to a uniform distribution on the points satisfying (5). Let  $A$  be the random set of all chosen points and let  $S = S(A)$  be the union of all the sets  $B(\mathbf{a})$ , where  $\mathbf{a} \in A$ . Clearly

$$|S| < (\lceil d \log n \rceil + 1)(2c_4 + 5)^d n^{d(d-1)/(2d-1)} \leq c_6 n^{d(d-1)/(2d-1)} \log n,$$

where  $c_6 = c_6(d)$  is a constant depending only on  $d$ .

By Lemma 3.2, for every fixed lattice point  $(x_1, \dots, x_d)$  in  $L(n, d)$ , the probability that there is **no** point  $\mathbf{a} = (a_1, \dots, a_d)$  among the chosen ones so that the assertion of the lemma holds for this point is at most  $(\frac{1}{2})^{\lceil d \log n \rceil + 1} < 1/n^d$ . Hence, the expected number of points  $(x_1, \dots, x_d)$  for which there is no such  $\mathbf{a} \in A$  is less than 1 and hence there is a choice for  $A$  such that there are no such points.

However, by Lemma 3.3, for such an  $A$ ,  $S(A)$  is an  $(n, d)$  covering set. We have thus proved that if  $n \geq c_5(d)$  then  $t(n, d) \leq c_6(d) n^{d(d-1)/(2d-1)} \log n$ , and this clearly implies the upper bound in Theorem 1.1 with, e.g.,  $c_2(d) = \max\{c_6(d), c_5(d)\}$ .  $\square$

It would be interesting to decide if the  $\log n$  factor in the upper bound in Theorem 1.1 is necessary.

## References

- [1] Richard K. Guy, *Unsolved Problems in Number Theory*, Springer Verlag, New York and Heidelberg, 1981, p. 133.
- [2] G. H. Hardy and E. M. Wright, *An Introduction to the Theory of Numbers*, (Forth Edition), Oxford University Press, London, 1959, Chapter XI.