# Correction: Basic Network Creation Games 

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#### Abstract

We prove a previously stated but incorrectly proved theorem: there is a diameter-3 graph in which replacing any edge $\{v, w\}$ of the graph with $\left\{v, w^{\prime}\right\}$, for any vertex $w^{\prime}$, does not decrease the total sum of distances from $v$ to all other nodes (a property called sum equilibrium).


Theorem 5 in [1] states that there exists a diameter-3 sum equilibrium graph, that is, an undirected graph such that, for every edge $\{v, w\}$ and every node $w^{\prime}$, replacing edge $\{v, w\}$ with $\left\{v, w^{\prime}\right\}$ does not decrease the total sum of distances from $v$ to all other nodes (and thus no vertex $v$ has incentive to swap an incident edge). In this short note, we observe an error in the original construction and proof, but present a different example that is indeed a diameter- 3 sum equilibrium graph, thereby restoring the theorem.

First we describe why Figure 3 of [1] is not in sum equilibrium. Specifically, vertex $d_{1}$ has an incentive to replace the edge $\left\{d_{1}, c_{1,1}\right\}$ with $\left\{d_{1}, c_{2,1}\right\}$, as the total distance is 27 in the first case and 26 in the last. The original proof ignores that $c_{2,1}$ is a neighbor of $c_{1,1}$ and, hence, Lemma 8 of [1] implies that the distance from $d_{1}$ to $c_{1,1}$ increases by 1 and not by 2 as claimed.

Figure 1 below presents a diameter- 3 sum equilibrium graph $G$ (which is also simpler than the original construction). In this instance, vertices $v_{2}, v_{4}, v_{5}$, and $v_{7}$ have local diameter 2 so, by Lemma 6 of [1], they have no incentive to swap any edge. (Lemma 6 states that a vertex of local diameter 2 never has incentive to swap an incident edge, as the number of distance- 1 neighbors remains fixed, and thus the number of nodes at distance $\geq 2$ remains fixed, so keeping their distances equal to 2 is optimal.) Among the remaining vertices, by symmetry, it suffices to prove that $v_{1}$ and $v_{3}$ do not have an incentive to swap edges.

Consider vertex $v_{i}$ for $i \in\{1,3\}$. Let $G_{-i}$ be the graph obtained by removing vertex $v_{i}$ and its incident edges; refer to Figure 2. The sum of distances for $v_{i}$ in $G$ is 13 . Because $v_{i}$ has degree 2 , the smallest possible sum of distances for $v_{i}$ is 12 , which can be obtained if $v_{i}$ were connected to two vertices that form a dominating set in $G_{-i}$. (A dominating set of cardinality larger than

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Figure 1: A diameter-3 sum equilibrium graph.
two can safely be ignored because $v_{i}$, having degree 2 , cannot connect to all its vertices in order to reduce the sum of distances to less than 13.) Furthermore, the only dominating set in $G_{-i}$ with cardinality 2 consists of vertices with degree 3 in $G_{-i}$, i.e., vertices $v_{4}$ and $v_{7}$ for $G_{-1}$ and vertices $v_{5}$ and $v_{7}$ for $G_{-3}$. (To see that, note that, because $G_{-i}$ contains 7 vertices, the dominating set should contain at least one vertex of degree 3, and the subgraph of $G_{-i}$ obtained after removing a vertex of degree 3 and its neighbors consists of a line of three vertices; clearly, the middle vertex of the line, which in all cases has degree 3 in $G_{-i}$, is the only possible choice for inclusion in the dominating set.)


Figure 2: Graphs $G_{-1}$ (top) and $G_{-3}$ (bottom). Grey vertices form the only dominating sets of cardinality two.

We conclude that, in order for vertex $v_{i}$ to reduce the sum of distances to $12, v_{i}$ should connect to both vertices of $G$ that form a dominating set in $G_{-i}$. The claim follows by noticing that, in $G, v_{i}$ is not connected to any of these two vertices, and hence cannot improve its sum of distances with a single edge swap. This concludes the proof of Theorem 5 in [1].

We have verified by exhaustive computer search that no graph with fewer than eight vertices is in sum equilibrium. Among graphs with eight vertices, Figure 1 has the fewest number 10 of edges, along with one other graph in which edge $\left\{v_{2}, v_{7}\right\}$ is replaced by $\left\{v_{3}, v_{6}\right\}$; there are also examples with 11 and 12 edges.

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## References

[1] N. Alon, E. D. Demaine, M. T. Hajiaghayi, and T. Leighton. Basic network creation games. SIAM Journal on Discrete Mathematics, 27(2):656-668, 2013.


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