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# On Cilleruelo's conjecture for the least common multiple of polynomial sequences

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**Abstract.** A conjecture due to Cilleruelo states that for an irreducible polynomial  $f$  with integer coefficients of degree  $d \geq 2$ , the least common multiple  $L_f(N)$  of the sequence  $f(1), f(2), \dots, f(N)$  has asymptotic growth  $\log L_f(N) \sim (d-1)N \log N$  as  $N \rightarrow \infty$ . We establish a version of this conjecture for almost all shifts of a fixed polynomial, the range of  $N$  depending on the range of shifts.

## 1. Introduction

### 1.1. Background

It is a well known and elementary fact that the least common multiple of all integers  $1, 2, \dots, N$  is exactly given by

$$\log \operatorname{lcm}\{1, 2, \dots, N\} = \psi(N) := \sum_{n \leq N} \Lambda(n),$$

with  $\Lambda(n)$  being the von Mangoldt function, and hence by the prime number theorem,

$$\log \operatorname{lcm}\{1, 2, \dots, N\} \sim N.$$

For a polynomial  $f \in \mathbb{Z}[X]$ , set

$$L_f(N) := \operatorname{lcm}\{f(n) : n = 1, \dots, N\}.$$

The goal is to understand the asymptotic growth of  $\log L_f(N)$  as  $N \rightarrow \infty$ .

In the linear case  $\deg f = 1$ , we still have  $\log L_f(N) \sim c_f N$  from the prime number theorem in arithmetic progressions, see e.g. [1]. A similar growth occurs for products of linear polynomials, see [3]. However, in the case of irreducible polynomials of higher degree, Cilleruelo [2] conjectured that the growth is faster than linear, precisely:

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**Conjecture 1.1.** *If  $f$  is an irreducible polynomial with  $\deg f \geq 2$ , then*

$$\log L_f(N) \sim (\deg f - 1)N \log N, \quad N \rightarrow \infty.$$

Cilleruelo proved Conjecture 1.1 for quadratic polynomials. Moreover, in that case there is a secondary main term (see also [6]):

$$\log L_f(N) = N \log N + b_f N + o(N).$$

No other case of Conjecture 1.1 is known to date. We do know that for any irreducible  $f$  of degree  $d \geq 3$ , we have an upper bound precisely compatible with the conjecture:  $\log L_f(N) \lesssim (d-1)N \log N$ , and a lower bound of the correct order of magnitude:  $\log L_f(N) \gg N \log N$ , see [4].

We will show that Conjecture 1.1 holds for *almost all*  $f$  in a suitable sense.

## 1.2. General setup

We fix a polynomial  $f_0(x) \in \mathbb{Z}[x]$  of degree  $d \geq 3$ , which we assume is monic, and for  $a \in \mathbb{Z}$  we set

$$f_a(x) = f_0(x) - a.$$

It is known that  $f_a(x)$  is generically irreducible. Set

$$L_a(N) = \text{lcm}\{f_0(n) - a : n = 1, \dots, N\}.$$

We want to show that:

**Theorem 1.2.** *For almost all  $|a| \leq T$ , and for all  $N$  satisfying*

$$T^{1/(d-1)} < N < \frac{T}{\log T},$$

*we have*

$$(1.1) \quad \log L_a(N) \sim (d-1)N \log N.$$

**Remark.** What one would like to show is that (1.1) holds for *all*  $N > N_0(a)$ , for all but  $o(T)$  values of  $|a| \leq T$ . At this time we do not know how to do this.

## 1.3. Plan

Let

$$P_a(N) = \prod_{n \leq N} |f_0(n) - a|.$$

We write down the prime power factorization

$$P_a(N) = \prod_p p^{\alpha_p(a; N)}.$$

If  $T \ll N^{d-1}$ , then  $\alpha_p(a; N) = 0$  for  $p \gg N^d$ , and (see Lemma 2.3 below)

$$\log P_a(N) = dN \log N + O(N).$$

We also write the prime power factorization of  $L_a(N)$  as

$$L_a(N) = \prod_p p^{\beta_p(N)}.$$

Let

$$D(a) = \text{disc}(f_0(x) - a)$$

be the discriminant of  $f_0(x) - a$ . It is a polynomial in  $a$  of degree  $d - 1$ , with integer coefficients.

We will show (see Proposition 2.2) that

$$(1.2) \quad \log L_a(N) = dN \log N - \text{Bad}_N(a) - \Delta_N(a) - NC_N(a) + O(N),$$

where

$$\begin{aligned} \text{Bad}_N(a) &= \sum_{\substack{p \leq N \\ p \nmid D(a)}} \alpha_p(N) \log p, \\ \Delta_N(a) &= \sum_{N < p \ll N^d} (\alpha_p(N) - \beta_p(N)) \log p, \\ C_N(a) &= \sum_{\substack{p \leq N \\ p \nmid D(a)}} \frac{\log p}{p-1} \rho(a; p), \end{aligned}$$

with

$$\rho(a; d) = \#\{n \bmod d : f_0(n) - a = 0 \bmod d\}.$$

We will show that for almost all  $|a| \leq T$ , with  $N \log N < T < N^{d-1}$ , we have

$$(1.3) \quad C_N(a) \sim \log N,$$

$$(1.4) \quad \text{Bad}_N(a) \ll N(\log \log N)^{1+o(1)},$$

$$(1.5) \quad \Delta_N(a) \ll N(\log \log N)^{1+o(1)}.$$

Inserting these into (1.2) will prove Theorem 1.2.

To prove (1.3), (1.4) and (1.5) we use averaging: denoting by  $\langle \bullet \rangle$  the average over all  $|a| \leq T$  such that  $f_0(x) - a$  is irreducible, we show that for  $N \log N < T < N^{d-1}$ ,

$$(1.6) \quad \langle |C_N(a) - \log N|^2 \rangle \ll (\log \log N)^2,$$

$$(1.7) \quad \langle \text{Bad}_N(a) \rangle \ll N \log \log N,$$

$$(1.8) \quad \langle \Delta_N(a) \rangle \ll N \log \log N.$$

Noting that  $\Delta_N(a), \text{Bad}_N(a) \geq 0$  are non-negative, we obtain (1.3), (1.4), (1.5) from the Chebyshev/Markov inequality.

**Remark.** In the deterministic case ( $a$  fixed,  $N \rightarrow \infty$ ), the quantities  $\text{Bad}_N$  and  $C_N$  can be handled easily, as in the quadratic case  $d = 2$ , see [2]. It is the quantity  $\Delta_N(a)$  which we do not know how to show is  $o(N \log N)$  individually (though the upper bound  $O(N \log N)$  is easy). This is why we need to average over  $a$ . However, letting  $a$  grow with  $N$  introduces new problems, in particular for the study of  $C_N$ , which may need the generalized Riemann hypothesis to overcome individually. The results (1.6) and (1.7) for random  $a$  are much easier and this is the method that we use.

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## 2. Background

### 2.1. Generic irreducibility

Fix  $f_0(x) \in \mathbb{Z}[x]$  monic, of degree  $d \geq 2$ . It is known that  $f_0(x) - a$  is generically irreducible; in fact (see Section 9.7 of [8]):

**Lemma 2.1.** *Fix  $f_0(x) \in \mathbb{Z}[x]$  of degree  $d \geq 2$ . Then the number of  $|a| \leq T$  for which  $f_0(x) - a$  is reducible is  $O(\sqrt{T})$ .*

This is sharp in this generality, since for even degree  $d = 2m$ , for the polynomial  $f_0(x) = x^{2m}$  we have  $x^{2m} - a$  is reducible whenever  $a = b^2$  is a perfect square.

Denote

$$D(a) = \text{disc}(f_0(x) - a)$$

the discriminant of  $f_a(x)$ , which is a polynomial in  $a$  of degree  $\leq d - 1$  with integer coefficients (depending on the coefficients of  $f_0$ ). We assume that  $a$  is such that  $f_0(x) - a$  is irreducible, and therefore  $D(a)$  is not zero, i.e.,  $f_a$  has no multiple roots.

Examples:

- i)  $f_0(x) = x^3$ , then  $\text{disc}(f_0(x) - a) = \text{disc}(x^3 - a) = -27a^2$ .
- ii)  $f_0(x) = x^3 - 3x$ , then  $\text{disc}(f_0(x) - a) = -27(a - 2)(a + 2)$ .

### 2.2. A decomposition

**Proposition 2.2.** *For  $|a| \leq N^{d-1}$  such that  $f_0(x) - a$  is irreducible, we have*

$$\log L_a(N) = d \log N - \text{Bad}_N(a) - NC_N(a) - \Delta_N(a) + O(N).$$

For  $a \in \mathbb{Z}$  such that  $f_a(x) = f_0(x) - a$  is irreducible, let

$$P_a(N) := \prod_{n \leq N} |f_a(n)|,$$

which is nonzero since  $f_a$  has no rational roots, and write the prime power decomposition as

$$P_a(N) = \prod_p p^{\alpha_p(N)}$$

so that

$$\alpha_p(N) = \sum_{n \leq N} \nu_p(f_a(n)),$$

where  $\nu_p(m) := \max(k \geq 0 : p^k \mid m)$ .

Following Cilleruelo [2], we want to relate  $\log L_a(N)$  to  $\log P_a(N)$ , which is clearly bigger. We write the prime power decomposition of  $L_a(N)$  as

$$L_a(N) = \prod_p p^{\beta_p(N)}, \quad \text{with } \beta_p(N) = \max\{\nu_p(f_a(n)) : n \leq N\}.$$

Using the prime factorization of  $L_a(N)$  and  $P_a(N)$ , we have

$$(2.1) \quad \begin{aligned} \log L_a(N) &= \log P_a(N) - \sum_{p \leq N} \alpha_p(N) \log p + \sum_{p \leq N} \beta_p(N) \log p \\ &\quad - \sum_{p > N} (\alpha_p(N) - \beta_p(N)) \log p, \end{aligned}$$

where we have separated out the contribution of primes  $p \leq N$ , and the larger ones. We further break off the contribution of primes  $p \leq N$  which divide the discriminant  $D(a) = \text{disc}(f_a)$ , by setting

$$\text{Bad}_N(a) := \sum_{\substack{p \leq N \\ p \nmid D(a)}} \alpha_p(N) \log p,$$

and abbreviate the contribution of big primes  $p > N$  as

$$\Delta_N(a) := \sum_{p > N} (\alpha_p(N) - \beta_p(N)) \log p.$$

Note that  $\text{Bad}_N$  and  $\Delta_N \geq 0$  are both non-negative. We obtain the expression

$$(2.2) \quad \begin{aligned} \log L_a(N) &= \log P_a(N) + \sum_{p \leq N} \beta_p(N) \log p \\ &\quad - \text{Bad}_N(a) - \sum_{\substack{p \leq N \\ p \nmid D(a)}} \alpha_p(N) \log p - \Delta_N(a). \end{aligned}$$

### 2.3. The quantity $P_a(N)$

**Lemma 2.3.** *For  $f_0(x) \in \mathbb{Z}[x]$  monic of degree  $d$ , and for  $|a| \ll N^{d-1}$  so that  $f_0(x) - a$  is irreducible, we have*

$$\log P_a(N) = dN \log N + O(N).$$

*Proof.* Write

$$\log P_a(N) = \sum_{n \leq N} \log |f_0(n) - a|.$$

Since we assume  $f_0(x) - a$  is irreducible, none of the factors  $f_0(n) - a$  can vanish, so that  $\log P_a(N)$  is well defined. If  $f_0(x) = x^d + c_{d-1}x^{d-1} + \dots$ , we have for  $n \leq N$ ,

$$f_0(n) - a = n^d \left( 1 + \frac{c_{d-1}}{n} + \frac{c_{d-2}}{n^2} + \dots - \frac{a}{n^d} \right).$$

Consider first the  $n$ 's satisfying  $N/\log N < n \leq N$ , for which we use (recall that  $|a| \leq N^{d-1}$ )

$$\log \left( 1 + \frac{c_{d-1}}{n} + \frac{c_{d-2}}{n^2} + \dots - \frac{a}{n^d} \right) = O\left(\frac{(\log N)^d}{N}\right),$$

so that

$$\sum_{\frac{N}{\log N} < n \leq N} \log |f_0(n) - a| = \sum_{\frac{N}{\log N} < n \leq N} \left( d \log n + O\left(\frac{(\log N)^d}{N}\right) \right) = dN \log N + O(N).$$

For  $1 \leq n \leq N/\log N$ , we just use  $1 \leq |f_0(n) - a| \ll N^d$  so that  $0 \leq \log |f_0(n) - a| \ll \log N$ , and

$$\sum_{n \leq N/\log N} \log |f_0(n) - a| \ll \sum_{n \leq N/\log N} \log N \ll N.$$

Hence

$$\log P_a(N) = dN \log N + O(N)$$

as claimed.  $\square$

**2.3.1. Dealing with  $\beta_p(N)$ .** For  $a$  such that  $f_0(x) - a$  is irreducible, we have

$$\beta_p(N) \ll \frac{\log N}{\log p}$$

because

$$\beta_p(N) = \max_{n \leq N} \max(k \geq 0 : p^k \mid f_0(n) - a),$$

and since  $f_0(n) - a \neq 0$  for all  $n$ , if  $p^k \mid f_0(n) - a \neq 0$ , then

$$k \leq \frac{\log |f_0(n) - a|}{\log p} \ll \frac{\log n + \log |a|}{\log p}.$$

Hence, since  $|a| \ll N^{d-1}$ ,

$$\beta_p(N) \ll \frac{\log N}{\log p},$$

and hence the contribution of primes  $p \leq N$  to (2.2) is

$$(2.3) \quad \sum_{p \leq N} \beta_p(N) \log p \ll \sum_{p \leq N} \log N \ll N.$$

**2.3.2. Dealing with  $\alpha_p(N)$ .** Using Hensel's lemma, it is easy to check that (see [5], and also Lemma 4 in [2]):

**Lemma 2.4.** *For  $p \nmid D(a) = \text{disc}(f_0(x) - a)$ , we have*

$$\alpha_p(N) = N \frac{\rho(a; p)}{p-1} + O\left(\frac{\log N}{\log p}\right),$$

where  $\rho(a; p) = \#\{n \bmod p : f_0(n) - a = 0 \bmod p\}$ .

Consequently, we find that in (2.2),

$$\sum_{\substack{p \leq N \\ p \nmid D(a)}} \alpha_p(N) \log p = NC_N(a) + O(N), \quad \text{where} \quad C_N(a) := \sum_{\substack{p \leq N \\ p \nmid D(a)}} \frac{\log p}{p-1} \rho(a; p).$$

Therefore we have proven Proposition 2.2.

### 3. Bounding $\text{Bad}_N$ almost surely

Recall that we defined

$$\text{Bad}_N(a) = \sum_{\substack{p \leq N \\ p \nmid D(a)}} \log p \sum_{n \leq N} \#\{k \geq 1 : p^k \mid f_0(n) - a\}$$

(we assume that  $f_0(x) - a$  is irreducible).

We denote the averaging operator over  $|a| \leq T$  such that  $f_0(x) - a$  is irreducible by

$$\langle \bullet \rangle = \frac{1}{\#\{|a| \leq T : f_0(x) - a \text{ is irreducible}\}} \sum_{\substack{|a| \leq T \\ f_0(x) - a \text{ irreducible}}} \bullet$$

The number of  $|a| \leq T$  for which  $f_0(x) - a$  is reducible is  $O(\sqrt{T})$  (Lemma 2.1), so that

$$(3.1) \quad \langle \bullet \rangle = \frac{1}{2T + O(\sqrt{T})} \sum_{\substack{|a| \leq T \\ f_0(x) - a \text{ irreducible}}} \bullet$$

**Proposition 3.1.** *If  $T \geq N$  but  $\log T \ll \log N$ , then*

$$\langle \text{Bad}_N \rangle \ll N \log \log N.$$

*Proof.* We separate out the contribution  $B_1(a)$  of  $k = 1$  and the contribution  $B_2(a)$  of the remaining  $k \geq 2$ :

$$\text{Bad}_N(A) = B_1(a) + B_2(a),$$

where

$$B_1(a) = \sum_{\substack{p \leq N \\ p \nmid D(a)}} \log p \#\{n \leq N : f_0(n) = a \bmod p\}$$

and

$$B_2(a) = \sum_{\substack{p \leq N \\ p|D(a)}} \log p \sum_{n \leq N} \#\{k \geq 2 : p^k \mid f_0(n) - a\}.$$

We will show that

$$(3.2) \quad B_1(a) \ll N \log \log N$$

and that

$$\langle B_2 \rangle \ll N,$$

proving Proposition 3.1.

We first show that

$$B_1(a) \ll N \log \log |D(a)|,$$

which suffices for (3.2) since  $\log |D(a)| \ll \log T \ll \log N$ .

Indeed, for  $p \leq N$  we have

$$\begin{aligned} \#\{n \leq N : f_0(n) = a \pmod{p}\} &= \left( \frac{N}{p} + O(1) \right) \#\{n \pmod{p} : f_0(n) = a \pmod{p}\} \\ &\ll \frac{N}{p} \rho(a; p), \end{aligned}$$

where

$$\rho(a; p) := \#\{n \pmod{p} : f_0(n) = a \pmod{p}\},$$

which we see by dividing the interval  $[1, N]$  into consecutive intervals of length  $p$ .

Since  $f_0(x)$  is a monic polynomial of degree  $d$ , it is nonzero modulo  $p$  and still of degree  $\leq d$ , hence  $\rho(a; p) \leq d$ . Thus

$$B_1(a) \ll \sum_{\substack{p \leq N \\ p|D(a)}} \log p \frac{N}{p} \rho(a; p) \ll N \sum_{p|D(a)} \frac{\log p}{p}.$$

We use:

**Lemma 3.2.** For  $k > 1$ ,

$$\sum_{p|k} \frac{\log p}{p} \ll \log \log k.$$

*Proof.* Indeed, splitting the sum into small primes  $p \leq \log k$ , and the rest (where the summands are at most  $\log \log k / \log k$ ), we get

$$\begin{aligned} \sum_{p|k} \frac{\log p}{p} &\leq \sum_{\substack{p|k \\ p \leq \log k}} \frac{\log p}{p} + \sum_{\substack{p|k \\ p > \log k}} \frac{\log p}{p} \ll \sum_{p \leq \log k} \frac{\log p}{p} + \frac{\log \log k}{\log k} \sum_{p|k} 1 \\ &\ll \log \log k + \frac{\log \log k}{\log k} \cdot \frac{\log k}{\log \log k} \ll \log \log k, \end{aligned}$$

since the number of distinct prime divisors of  $k$  is  $\ll \log k / \log \log k$ .  $\square$



Therefore

$$\sum_{p|D(a)} \frac{\log p}{p} \ll \log \log |D(a)| \ll \log \log |a|,$$

and we obtain

$$B_1(a) \ll N \log \log |D(a)|.$$

Next we bound the mean value of  $B_2(a)$ :

$$\langle B_2 \rangle = \frac{1}{2T + O(\sqrt{T})} \sum_{\substack{|a| \leq T \\ f_0(x) - a \\ \text{irreducible}}} \sum_{\substack{p \leq N \\ p|D(a)}} \log p \sum_{k \geq 2} \mathbf{1}(f_0(n) = a \bmod p^k).$$

Now if  $f_0(x) - a$  is irreducible, then  $f_0(n) - a \neq 0$ , and so if  $p^k | f_0(n) - a$  with  $n \leq N$ , then  $k \ll \log N / \log p$ , so we restrict the summation to  $2 \leq k \ll \log N / \log p$ . Moreover, given  $n$ , the condition  $f_0(n) = a \bmod p^k$  determines  $a$  modulo  $p^k$ , so there are  $\ll T/p^k + 1$  choices for  $a$ . Hence we may bound

$$\begin{aligned} \langle B_2 \rangle &\ll \frac{1}{T} \sum_{p \leq N} \log p \sum_{n \leq N} \sum_{2 \leq k \ll \log N / \log p} \left( \frac{T}{p^k} + 1 \right) \\ &= \frac{N}{T} \sum_{p \leq N} \log p \sum_{2 \leq k \ll \log N / \log p} \left( \frac{T}{p^k} + 1 \right). \end{aligned}$$

We have

$$\frac{N}{T} \sum_{p \leq N} \log p \sum_{2 \leq k \ll \log N / \log p} \frac{T}{p^k} \ll N \sum_{p \leq N} \log p \sum_{k \geq 2} \frac{1}{p^k} \ll N \sum_{p \leq N} \frac{\log p}{p^2} \ll N$$

and

$$\frac{N}{T} \sum_{p \leq N} \log p \sum_{2 \leq k \ll \log N / \log p} 1 \ll \frac{N}{T} \sum_{p \leq N} \log p \cdot \frac{\log N}{\log p} \ll \frac{N^2}{T}.$$

Altogether we find

$$\langle B_2 \rangle \ll N + \frac{N^2}{T},$$

which is  $O(N)$  if  $T \geq N$ . □

#### 4. Averaging $\Delta_N(a)$

Let

$$\Delta_N(a) = \sum_{p > N} \log p (\alpha_p(N) - \beta_p(N)).$$

Then clearly  $\Delta_N \geq 0$ , and we want to show:

**Proposition 4.1.** *Assume that  $T \geq N \log N$ , but  $\log T \ll \log N$ . Then*

$$\langle \Delta_N \rangle \ll_{f_0} N \log \log N.$$

#### 4.1. Preparations

Let

$$G(m, n) = \frac{f_0(m) - f_0(n)}{m - n},$$

which, given  $n$ , is a (nonzero) polynomial in  $m$ , of degree  $\leq d - 1$ . If  $f_0$  is monic then so is  $G(m, n)$ , so its degree is exactly  $d - 1$ .

**Lemma 4.2.** *There is some  $C_1 = C_1(f_0)$  so that if  $m, n \geq 1$  and  $\max(m, n) > C_1$ , then  $G(m, n) \neq 0$ .*

*Proof.* We have

$$G(m, n) = \sum_{j=1}^d c_j \frac{m^j - n^j}{m - n},$$

and if  $j \geq 2$  then, for  $n = \max(m, n)$ ,

$$\frac{m^j - n^j}{m - n} = n^{j-1} + n^{j-2}m + \cdots + m^{j-1} \leq jn^{j-1},$$

while

$$\frac{m^d - n^d}{m - n} = n^{d-1} + n^{d-2}m + \cdots + m^{d-1} > n^{d-1},$$

so that (assuming  $f_0$  monic, so  $c_d = 1$ )

$$G(m, n) \geq \frac{m^d - n^d}{m - n} - \sum_{j=2}^{d-1} |c_j| \frac{m^d - n^d}{m - n} - |c_1| > n^{d-1} - \sum_{j=1}^{d-1} |c_j| j n^{j-1},$$

which is clearly positive once  $n$  is sufficiently large in terms of the coefficients  $c_1, \dots, c_{d-1}$  of  $f_0$ .  $\square$

**Lemma 4.3.** *There is some  $C(d) > 0$  so that for all  $|a| \leq N^d$  such that  $f_a(x) = f_0(x) - a$  is irreducible, we have  $\alpha_p(N) \leq C(d)$  if  $p > N$ . Moreover,  $\alpha_p(N) = 0$  unless  $p \ll N^d + |a|$ .*

*Proof.* We have, by definition,

$$\alpha_p(N) = \sum_{n \leq N} \sum_{k \geq 1} \mathbf{1}(f_0(n) = a \bmod p^k) = \sum_{k \geq 1} \#\{n \leq N : f_0(n) = a \bmod p^k\}.$$

Since we assume that  $f_a(x) = f_0(x) - a$  is irreducible, hence has no rational zeros, we must have, if  $p \mid f_a(n)$ , that  $p \leq |f_a(n)| \ll N^d + |a| \ll N^d$  uniformly in  $|a| \leq T$  (recall  $T \leq N^d$ ). Hence  $\alpha_p(N) = 0$  for  $p \gg N^d$ .

Given  $n$  so that  $p \mid f_a(n)$ , with  $p > N$ , we claim that there are at most  $d$  such integers:

$$\#\{m \leq N : f_a(m) = f_a(n) \bmod p\} \leq d.$$

Indeed, for any  $c \in \mathbb{Z}/p\mathbb{Z}$ , the number of solutions  $m \bmod p$  of  $f_a(m) = c \bmod p$  is at most  $d$ , and since  $p > N$ , this certainly applies to those  $m \leq N$  which solve  $f_a(m) = c$  with  $c = f_a(n)$ .

Moreover, if  $p > N$ , the maximal  $k$  so that  $p^k \mid f_0(n) - a$  for some  $n \leq N$  is, because we assume  $f_a(n) \neq 0$ ,

$$\ll \frac{\log(N^d + |a|)}{\log p} = O_d(1)$$

since we assume that  $|a| \leq T$  with  $\log T \ll \log N$ .

Therefore

$$\alpha_p(N) = \sum_{k \geq 1} \#\{n \leq N : f(n) = 0 \bmod p^k\} \leq \sum_{1 \leq k \ll O_d(1)} d = O_d(1)$$

as claimed.  $\square$

## 4.2. A preliminary bound on $\Delta_N(a)$

**Lemma 4.4.** *If  $a$  is such that  $f_0(x) - a$  has no rational zeros, and  $\log |a| \ll \log N$ , then*

$$(4.1) \quad \Delta_N(a) \ll \sum_{\substack{1 \leq m < n \leq N \\ G(m,n) \neq 0}} \sum_{\substack{N < p \ll N^d \\ p \mid f_0(m) - a \\ p \mid G(m,n)}} \log p + O(\log N).$$

*Proof.* We have  $\alpha_p(N) \neq \beta_p(N)$  if and only if there are two distinct integers  $m, n \leq N$  so that  $p \mid f_a(m)$  and  $p \mid f_a(n)$ . Using Lemma 4.3, we see that  $\alpha_p(N) - \beta_p(N) = O_d(1)$  for  $p > N$ , and hence applying a union bound we obtain, if  $a$  is such that  $f_a(x)$  has no rational zeros,

$$\Delta_N(a) \ll_d \sum_{1 \leq m < n \leq N} \sum_{\substack{N < p \ll N^d \\ p \mid f_0(m) - a \\ p \mid f_0(n) - a}} \log p.$$

Note that if  $p \mid f_a(m)$  and  $p \mid f_a(n)$  then  $p \mid f_a(m) - f_a(n) = (m - n)G(m, n)$ , and so since  $p \nmid m - n$  (because  $1 \leq n - m \leq N - 1 < p$ ), we must have  $p \mid G(m, n)$ . Thus

$$(4.2) \quad \Delta_N(a) \ll \sum_{1 \leq m < n \leq N} \sum_{\substack{N < p \ll N^d \\ p \mid f_0(m) - a \\ p \mid G(m,n)}} \log p.$$

We break off the terms corresponding to  $G(m, n) = 0$ . According to Lemma 4.2, the condition  $G(m, n) = 0$  forces  $m, n \leq C_1$  to be bounded. Hence the contribution of such pairs to (4.2) is bounded by

$$\ll \sum_{m, n \leq C_1} \sum_{\substack{N < p \ll N^d \\ p \mid f_0(m) - a}} \log p \ll \log N \max_{m \leq C_1} \#\{p > N : p \mid a - f_0(m)\}.$$

Note that  $0 < |f_0(m) - a| \ll |a| + 1$  if  $m \leq C_1$  (we assume that  $a$  is such that  $f_0(x) - a$  has no rational zeros, hence  $f_0(m) - a \neq 0$ ), and hence the number of primes  $p > N$  dividing  $f_0(m) - a$  is at most  $\ll \log |a| / \log N$ . Hence the contribution of pairs  $m < n$  with  $G(m, n) = 0$  to (4.2) is at most  $\ll \log |a|$ . Thus

$$\Delta_N(a) \ll \sum_{\substack{1 \leq m < n \leq N \\ G(m, n) \neq 0}} \sum_{\substack{N < p \ll N^d \\ p | f_0(m) \\ p | G(m, n)}} \log p + O(\log |a|).$$

Finally, the assumption  $\log |a| \ll \log N$  gives (4.1).  $\square$

### 4.3. Proof of Proposition 4.1

Now to average over  $|a| \leq T$  (such that  $f_0(x) - a$  is irreducible). Using (4.1), noting that  $\log |a| \ll \log T \ll \log N$  gives

$$\langle \Delta_N \rangle \ll \sum_{\substack{1 \leq m < n \leq N \\ G(m, n) \neq 0}} \sum_{\substack{N < p \ll N^d \\ p | G(m, n)}} \log p \frac{1}{T} \#\{|a| \leq T : p \mid a - f_0(m)\} + O(\log N).$$

Given  $1 \leq m < N$  and  $N < p \ll N^d$ , the number of  $|a| \leq T$  with  $a = f_0(m) \pmod p$  is  $\ll T/p + 1$ . Hence

$$\begin{aligned} \langle \Delta_N \rangle &\ll \sum_{\substack{1 \leq m < n \leq N \\ G(m, n) \neq 0}} \sum_{\substack{N < p \ll N^d \\ p | G(m, n)}} \frac{\log p}{p} + \frac{1}{T} \sum_{\substack{1 \leq m < n \leq N \\ G(m, n) \neq 0}} \sum_{\substack{N < p \ll N^d \\ p | G(m, n)}} \log p + O(\log N) \\ &=: I + II + O(\log N). \end{aligned}$$

To treat the sum  $II$ , we note if  $m, n \leq N$ , then  $|G(m, n)| \leq C(f_0)N^{d-1}$  and so there are at most  $d - 1$  distinct primes  $p > N$  which divide  $G(m, n)$  (which we assume is non-zero), and for these,  $\log p \ll \log N$ . Therefore

$$II \ll \frac{\log N}{T} \sum_{1 \leq m < n \leq N} (d - 1) \ll \frac{N^2 \log N}{T},$$

which is  $O(N)$  if  $T > N \log N$ .

To treat the sum  $I$ , we separate the prime sum into primes with  $N < p \leq N \log N$  and the remaining large primes  $N \log N < p \ll N^{d-1}$  to get

$$I \ll \sum_{\substack{1 \leq m < n \leq N \\ G(m, n) \neq 0}} \sum_{\substack{N < p \leq N \log N \\ p | G(m, n)}} \frac{\log p}{p} + \sum_{\substack{1 \leq m < n \leq N \\ G(m, n) \neq 0}} \sum_{\substack{N \log N < p \ll N^d \\ p | G(m, n)}} \frac{\log p}{p}.$$

We treat the sum over small primes by switching the order of summation:

$$\begin{aligned} & \sum_{\substack{1 \leq m < n \leq N \\ G(m,n) \neq 0}} \sum_{\substack{N < p < N \log N \\ p | G(m,n)}} \frac{\log p}{p} \\ & \leq \sum_{N < p < N \log N} \frac{\log p}{p} \#\{1 \leq m < n \leq N : G(m,n) = 0 \pmod p\}. \end{aligned}$$

Now given  $m$ , the congruence  $G(m,n) = 0 \pmod p$  (if solvable) determines  $n \pmod p$  up to  $d-1$  possibilities, since  $G(m,n)$  is a monic polynomial of degree  $d-1$  in  $n$ , and since  $n \leq N < p$  means that  $n$  is determined as an integer up to  $d-1$  possibilities. Hence

$$\#\{1 \leq m < n \leq N : G(m,n) = 0 \pmod p\} \leq (d-1)N,$$

and the sum over small primes is bounded by

$$\ll \sum_{N < p < N \log N} \frac{\log p}{p} N = N \{(\log(N \log N) + O(1)) - (\log N + O(1))\} \sim N \log \log N$$

on using Mertens' theorem.

The sum over large primes is treated by using  $\log p/p \ll 1/N$  for  $p > N \log N$ , giving

$$\sum_{\substack{1 \leq m < n \leq N \\ G(m,n) \neq 0}} \sum_{\substack{N \log N < p \ll N^d \\ p | G(m,n)}} \frac{\log p}{p} \ll \frac{1}{N} \sum_{\substack{1 \leq m < n \leq N \\ G(m,n) \neq 0}} \#\{p > N \log N : p | G(m,n)\}.$$

Now given  $1 \leq m < n \leq N$  with  $G(m,n) \neq 0$ , there are at most  $d-1$  primes  $p > N \log N$  dividing  $G(m,n) \ll N^{d-1}$ , so that the contribution of large primes is bounded by

$$\ll \frac{1}{N} \sum_{1 \leq m < n \leq N} (d-1) \ll N.$$

This gives  $I \ll N \log \log N$ , and hence,

$$\langle \Delta_N \rangle \ll N \log \log N,$$

as claimed.  $\square$

## 5. Almost sure behaviour of $C_N$

### 5.1. An identity involving $C_N(a)$

Let  $f \in \mathbb{Z}[x]$  be an irreducible polynomial, and let  $\rho_f(p)$  be the number of distinct roots of the polynomial  $f$  modulo a prime  $p$ . It is well known [5] that for fixed  $f$ , the mean value of  $\rho_f(p)$  over all primes is 1:

$$\frac{1}{\pi(x)} \sum_{p \leq x} \rho_f(p) = 1 + o_f(1).$$

We write

$$\rho_f(p) = 1 + \sigma_f(p),$$

where  $\sigma_f(p)$  is a fluctuating quantity, having mean zero.

Now fix

$$f_0(x) = x^d + c_{d-1}x^{d-1} + \cdots + c_1x \in \mathbb{Z}[x],$$

a monic polynomial of degree  $d$ , and for  $a \in \mathbb{Z}$  set

$$f_a(x) = f_0(x) - a.$$

Write  $\rho(a; p) = \rho_{f_a}(p)$  and  $\sigma(a; p) = \sigma_{f_a}(p)$ . Note that  $0 \leq \rho(a; p) \leq d$ .

We write

$$C_N(a) := \sum_{\substack{p \leq N \\ p \nmid \text{disc}(f)}} \frac{\log p}{p-1} \rho(a; p) = \sum_{p \leq N} \frac{\log p}{p} - E_N(a) + D_N(a) + O(1),$$

where

$$D_N(a) := \sum_{\substack{p \leq N \\ p \mid \text{disc}(f_a)}} \frac{\log p}{p} \sigma(a; p) \quad \text{and} \quad E_N(a) := \sum_{\substack{p \leq N \\ p \mid D(a)}} \frac{\log p}{p}.$$

By Mertens' theorem,

$$\sum_{p \leq N} \frac{\log p}{p} = \log N + O(1).$$

The contribution  $E_N(a)$  of primes dividing the discriminant  $D(a) = \text{disc}(f_0(x) - a)$  can be bounded individually, for  $|a| \leq T \ll N^d$ , using Lemma 3.2 (assuming  $D(a) \neq 0$ ):

$$E_N(a) = \sum_{\substack{p \leq N \\ p \mid D(a)}} \frac{\log p}{p} \leq \sum_{p \mid D(a)} \frac{\log p}{p} \ll \log \log |D(a)|.$$

Since  $D(a)$  is a polynomial of degree  $d-1$  in  $a$ , and  $|a| \leq T \ll N^d$ , we find

$$\sum_{\substack{p \leq N \\ p \mid D(a)}} \frac{\log p}{p} \ll \log \log N,$$

which is negligible relative to the main term. Hence

$$C_N(a) = \log N + D_N(a) + O(\log \log N).$$

In the following part, we will establish the following upper bound on the second moment of  $D_N(a)$ .

**Proposition 5.1.** *For  $T \geq N \log N$ , the second moment of  $D_N(a)$  satisfies*

$$\langle |D_N|^2 \rangle \ll 1.$$

Using the triangle inequality and Cauchy–Schwartz, we obtain:

**Proposition 5.2.**

$$\langle |C_N - \log N|^2 \rangle \ll (\log \log N)^2.$$

As a consequence, we deduce our main objective for this section.

**Proposition 5.3.** *For almost all  $|a| \leq T$  (with  $N \log N \leq T \ll N^{d-1}$ ),*

$$C_N(a) = \log N + O(\log \log N).$$

**5.2. Proof of Proposition 5.1**

*Proof.* Expanding, we have

$$\langle (D_N)^2 \rangle = \sum_{p \leq N} \sum_{q \leq N} \frac{\log p \log q}{pq} \langle \sigma(a; p) \sigma(a; q) \rangle.$$

The diagonal contribution  $p = q$  gives

$$\sum_{p \leq N} \frac{(\log p)^2}{p^2} \langle \sigma(a; p)^2 \rangle.$$

Now note that

$$-1 \leq \sigma(a; p) \leq d - 1$$

is uniformly bounded. This is because the polynomial  $f_0(x) - a$  is monic of degree  $d$ , hence has at most  $d$  zeros modulo  $p$ , so that  $0 \leq \rho(a; p) \leq d$  and so  $-1 \leq \sigma(a; p) \leq d - 1$ . Thus we obtain a bound for the diagonal sum:

$$\sum_{p \leq N} \frac{(\log p)^2}{p^2} \langle \sigma(a; p)^2 \rangle \ll \sum_{p \leq N} \frac{(\log p)^2}{p^2} \ll 1.$$

For the off-diagonal terms, we use:

**Lemma 5.4.** *For distinct primes  $p \neq q$ ,*

$$|\langle \sigma(\bullet; p) \sigma(\bullet; q) \rangle| \ll \frac{\sqrt{pq} \log(pq)}{T} + \frac{1}{\sqrt{T}}.$$

Therefore, given Lemma 5.4, we obtain

$$\begin{aligned} \sum_{p \neq q \leq N} \frac{\log p \log q}{pq} |\langle \sigma(a; p) \sigma(a; q) \rangle| &\ll \sum_{p \neq q \leq N} \frac{\log p \log q}{pq} \left( \frac{\sqrt{pq} \log(pq)}{T} + \frac{1}{\sqrt{T}} \right) \\ &\ll \frac{\log N}{T} \left( \sum_{p \leq N} \frac{\log p}{\sqrt{p}} \right)^2 + \frac{1}{\sqrt{T}} \left( \sum_{p \leq N} \frac{\log p}{p} \right)^2 \\ &\ll \frac{N \log N}{T} + \frac{(\log N)^2}{\sqrt{T}}, \end{aligned}$$

which is  $O(1)$  if  $T \geq N \log N$ , proving Proposition 5.1.  $\square$

### 5.3. Proof of Lemma 5.4

For the argument, it will be important to have  $a$  run over an interval. So we first remove the restriction on  $a$  in the averaging, that  $f_0(x) - a$  is irreducible. Since  $-1 \leq \sigma(a; p) \leq d - 1$ , this introduces an error bounded by

$$\ll \frac{1}{T} \sum_{\substack{|a| \leq T \\ f_0(x) - a \text{ reducible}}} (d-1)^2 \ll \frac{1}{T} \#\{|a| \leq T : f_0(x) - a \text{ reducible}\} \ll \frac{1}{\sqrt{T}},$$

and so

$$\langle \sigma(a; p) \sigma(a; q) \rangle = \frac{1}{2T + O(\sqrt{T})} \sum_{|a| \leq T} \sigma(a; p) \sigma(a; q) + O\left(\frac{1}{\sqrt{T}}\right).$$

We express  $\rho(a; p)$  as an exponential sum:

$$\rho(a; p) = \#\{x \bmod p : f_0(x) - a = 0 \bmod p\} = \sum_{x \bmod p} \frac{1}{p} \sum_{t \bmod p} e\left(\frac{t(f_0(x) - a)}{p}\right).$$

The term  $t = 0$  contributes the main term of 1, and we obtain the following expression for  $\sigma(a; p) = \rho(a; p) - 1$ :

$$(5.1) \quad \sigma(a; p) = \frac{1}{p} \sum_{t \neq 0 \bmod p} e\left(-\frac{at}{p}\right) \sum_{x \bmod p} e\left(\frac{tf_0(x)}{p}\right),$$

where  $e(z) := e^{2\pi iz}$ . Set

$$\mathcal{S}_{f_0}(b, n) := \sum_{x \bmod n} e\left(\frac{bf_0(x)}{n}\right).$$

Using (5.1), we have on switching orders of summation,

$$\begin{aligned} & \frac{1}{2T + O(\sqrt{T})} \sum_{|a| \leq T} \sigma(a; p) \sigma(a; q) \\ &= \frac{1}{2T + O(\sqrt{T})} \frac{1}{pq} \sum_{\substack{0 \neq t \bmod p \\ 0 \neq s \bmod q}} \sum_{|a| \leq T} e\left(-a\left(\frac{t}{p} + \frac{s}{q}\right)\right) \mathcal{S}_{f_0}(t, p) \mathcal{S}_{f_0}(s, q). \end{aligned}$$

Weil's bound [9], [7] shows that there is a constant  $c(d) > 0$  so that, for all primes  $p$  and all  $b$  coprime to  $p$ ,

$$(5.2) \quad |\mathcal{S}_{f_0}(b, p)| \leq c(d)\sqrt{p}.$$

In fact, for any  $f_0 \in \mathbb{Z}[x]$  with  $f_0(x)$  primitive of degree  $d$ , if  $p > d$  then

$$|\mathcal{S}_{f_0}(b, p)| \leq (d-1)\sqrt{p}.$$



Hence we find

$$\begin{aligned} |\langle \sigma(\bullet; p) \sigma(\bullet; q) \rangle| &\ll_d \frac{1}{T\sqrt{pq}} \sum_{\substack{0 \neq t \pmod p \\ 0 \neq s \pmod q}} \left| \sum_{|a| \leq T} e\left(-a\left(\frac{t}{p} + \frac{s}{q}\right)\right) \right| + O\left(\frac{1}{\sqrt{T}}\right) \\ &= \frac{1}{T\sqrt{pq}} \sum_{\substack{m \pmod{pq} \\ \gcd(m, pq)=1}} \left| \sum_{|a| \leq T} e\left(-\frac{am}{pq}\right) \right| + O\left(\frac{1}{\sqrt{T}}\right), \end{aligned}$$

where we have used that if  $p \neq q$  are distinct primes, then as  $t$  and  $s$  vary over all invertible residues modulo  $p$  (resp., modulo  $q$ ),  $tq + sp \pmod{pq}$  covers all invertible residues modulo  $pq$  exactly once.

We sum the geometric progression

$$\left| \sum_{|a| \leq T} e\left(-\frac{am}{pq}\right) \right| \ll \min\left(T, \left\| \frac{m}{pq} \right\|^{-1}\right)$$

where  $\|\alpha\| = \text{dist}(\alpha, \mathbb{Z})$ . We may take  $1 \leq m < pq/2$ , and then the bound is

$$\ll pq/m.$$

This gives

$$|\langle \sigma(\bullet; p) \sigma(\bullet; q) \rangle| \ll \frac{1}{T\sqrt{pq}} \sum_{\substack{1 \leq m \leq pq/2 \\ \gcd(m, pq)=1}} \frac{pq}{m} + O\left(\frac{1}{\sqrt{T}}\right) \ll \frac{\sqrt{pq} \log(pq)}{T} + O\left(\frac{1}{\sqrt{T}}\right),$$

proving Lemma 5.4. □

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