

## Contents

<b>5</b>	<b>Markov chains</b>	<b>47</b>
5a	Random walk on a regular graph . . . . .	47
5b	Finite Markov chains . . . . .	50
5c	Return time . . . . .	52

## 5 Markov chains

### 5a Random walk on a regular graph

Assume that a weakly connected<sup>1</sup> finite directed graph has  $m$  vertices and is regular (that is, each vertex has  $k$  outgoing edges and  $k$  incoming edges, the same  $k$  for all vertices). In addition, we assume *aperiodicity*: there exists no  $p \in \{2, 3, \dots\}$  such that the length of every cycle is divisible by  $p$ . A random walk started at a given vertex. Denote by  $S_n$  the position of the walk after  $n$  steps.

**5a1 Proposition.** For each vertex  $x$  of the graph,

$$\mathbb{P}(S_n = x) \rightarrow \frac{1}{m} \quad \text{as } n \rightarrow \infty.$$

This fact is a special case of a convergence theorem for Markov chains (see 5b3).<sup>2</sup>

Now assume that the random walk, started at a given vertex, stops on the first return to this vertex.

**5a2 Proposition.** The expected number of moves is equal to  $m$  (the number of vertices).

The proof uses Markov chains.<sup>3</sup> (Aperiodicity is not needed.)

Think, what happens if the graph consists of two large pieces connected by a thin neck.

Prop. 5a2 will be proved in Sect. 5c; Prop. 5a1 — in the end of this Sect. 5a.

First, some graph theory.

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<sup>1</sup>That is, the corresponding undirected graph is connected.

<sup>2</sup>[D, Sect. 5.4, Example 4.5; Sect. 5.5(a)]; [KS, Sect. 5.3].

<sup>3</sup>[D, Sect. 5.4, (4.6) and Example 4.5].

We consider an aperiodic regular weakly connected finite directed graph. The graph has a set  $V$  of vertices and a set  $E \subset V \times V$  of edges.<sup>1</sup> Weak connectedness:

$$\emptyset \subsetneq A \subsetneq V \implies E \cap ((A \times (V \setminus A)) \cup ((V \setminus A) \times A)) \neq \emptyset$$

for all  $A \subset V$ . Regularity:

$$\#\{y : (x, y) \in E\} = k = \#\{y : (y, x) \in E\}$$

for all  $x \in V$ .

**5a3 Lemma.** For every  $A \subset V$  the number of incoming edges is equal to the number of outgoing edges; that is,

$$\#(E \cap (A \times (V \setminus A))) = \#(E \cap ((V \setminus A) \times A)).$$

*Proof.* Denoting  $B = V \setminus A$  we have

$$\begin{aligned} E \cap (A \times V) &= E \cap (A \times B) \uplus E \cap (A \times A), \\ E \cap (V \times A) &= E \cap (B \times A) \uplus E \cap (A \times A), \end{aligned}$$

thus  $\#(E \cap (A \times B)) = k \cdot (\#A) - \#(E \cap (A \times A)) = \#(E \cap (B \times A))$ .  $\square$

**5a4 Corollary.** Strong connectedness:

$$\emptyset \subsetneq A \subsetneq V \implies E \cap (A \times (V \setminus A)) \neq \emptyset$$

for all  $A \subset V$ . (Closed sets:  $\emptyset$  and  $V$  only.)

**5a5 Corollary.** For all  $x, y \in V$  there exists a path (of *some* length) from  $x$  to  $y$ .

**5a6 Lemma.** There exists  $n$  such that for all  $x, y \in V$ , every  $t \geq n$  is the length of some (at least one) path from  $x$  to  $y$ .

*Proof.* The set  $L_x$  of lengths of all loops from  $x$  to  $x$  is a semigroup, therefore  $L_x - L_x$  is a group,  $L_x - L_x = p_x \mathbb{Z}$  for some  $p_x$ . By 5a5,  $L_x - L_x$  does not depend on  $x$ . Thus,  $p_x = 1$  for all  $x$ . It means existence of  $N_x$  such that  $N_x \in L_x$  and  $N_x + 1 \in L_x$ . We take  $n_x = N_x^2$  and note that<sup>2</sup>  $N_x^2 + kN_x + r = N_x(N_x + k) + r = N_x(N_x + k) - N_x r + (N_x + 1)r = N_x(N_x + k - r) + (N_x + 1)r \in L_x$ . We take  $m$  such that a path of length  $\leq m$  exists from every  $x$  to every  $y$ ; then  $n = m + \max_x n_x$  fits.  $\square$

<sup>1</sup>May intersect the diagonal. Multiple edges are excluded, but all said can be easily generalized to graphs with multiple edges.

<sup>2</sup>Example:  $\{10k + 11l\} \not\ni 78, 79, 89$ .

Now we return to probability.

We want to show that the initial point  $x_0$  is ultimately forgotten by the random walk  $(S_n)$ .

Given another starting point  $x'_0 \in V$ , we introduce the probability space  $\Omega'$  of paths (of length  $n$ ) starting at  $x'_0$ , and random variables  $S'_0, \dots, S'_n : \Omega' \rightarrow V$ . We take the product

$$\tilde{\Omega} = \Omega \times \Omega'$$

and treat  $S_t, S'_t$  as maps  $\tilde{\Omega} \rightarrow V$ . We get two *independent* random walks, one starting at  $x_0$ , the other at  $x'_0$ . In addition, we let  $\tilde{S}_t = (S_t, S'_t) : \tilde{\Omega} \rightarrow \tilde{V} = V \times V$ .

The reflection helps again! The transformation  $(x, y) \mapsto (y, x)$  of  $\tilde{V}$  will be treated as reflection, and the diagonal of  $\tilde{V}$  as the barrier. We define  $M_n : \tilde{\Omega} \rightarrow \{0, 1\}$  by

$$M_n = \begin{cases} 0 & \text{if } S_0 \neq S'_0, S_1 \neq S'_1, \dots, S_n \neq S'_n, \\ 1 & \text{otherwise.} \end{cases}$$

**5a7 Exercise.**  $\mathbb{E}(f(\tilde{S}_n)\mathbb{1}_{M_n=1}) = 0$  for every antisymmetric function  $f : \tilde{V} \rightarrow \mathbb{R}$  (“antisymmetric” means  $f(y, x) = -f(x, y)$ ).

Prove it.

Hint: similar to the proof of Lemma 4a2.

That is, the conditional distribution of  $\tilde{S}_n$  given  $M_n = 1$  is symmetric (if defined).

And again (recall 4a3),  $\mathbb{E}f(\tilde{S}_n) = \mathbb{E}(f(\tilde{S}_n)\mathbb{1}_{M_n=0})$ .

**5a8 Lemma.**  $|\mathbb{P}(S_n = x) - \mathbb{P}(S'_n = x)| \leq \mathbb{P}(M_n = 0)$ .

*Proof.* Take  $f(a, b) = \mathbb{1}_{\{x\}}(a) - \mathbb{1}_{\{x\}}(b)$  in 5a7. □

The probability of the event  $M_n = 0$  depends on  $n$ ,  $x_0$  and  $x'_0$ . We maximize it in  $x_0, x'_0$ :

$$\varepsilon_n = \max_{x_0, x'_0 \in V} \mathbb{P}(M_n = 0).$$

**5a9 Lemma.**  $\varepsilon_n \rightarrow 0$  as  $n \rightarrow \infty$ .

The proof will be given later.

Let  $p_n(x, y)$  denote the  $n$ -step transition probability from  $x$  to  $y$ . (Thus,  $\mathbb{P}(S_t = y) = p_t(x_0, y)$  and  $\mathbb{P}(S'_t = y) = p_t(x'_0, y)$ .)

Clearly,  $\sum_{y \in V} p_1(x, y) = 1$  for all  $x \in V$ ; but regularity ensures also  $\sum_{x \in V} p_1(x, y) = 1$  for all  $y \in V$ . By induction,  $\sum_{y \in V} p_n(x, y) = 1$  for all  $x \in V$ , and  $\sum_{x \in V} p_n(x, y) = 1$  for all  $y \in V$ .

**Proof of Prop. 5a1.** By Lemma 5a8,  $|p_n(x_0, y) - p_n(x'_0, y)| \leq \varepsilon_n$ . We average it in  $x'_0$ ; taking into account that  $\frac{1}{m} \sum_{x'_0 \in V} p_n(x'_0, y) = \frac{1}{m}$  we get  $|p_n(x_0, y) - \frac{1}{m}| \leq \varepsilon_n$ ; finally,  $\varepsilon_n \rightarrow 0$  by Lemma 5a9.  $\square$

*Proof of Lemma 5a9.* Lemma 5a6 gives us  $n$  such that  $p_n(x, y) \neq 0$  for all  $x, y$ . Clearly,  $p_n(x, y) \geq k^{-n}$ . Thus,

$$\mathbb{P}(M_n = 1) \geq \mathbb{P}(S_n = y, S'_n = y) \geq k^{-2n},$$

no matter which  $y$  is used. We put  $\theta = 1 - k^{-2n}$  and see that  $\mathbb{P}(M_n = 0) \leq \theta$ . But moreover,  $\mathbb{P}(M_{t+n} = 0 \mid M_t = 0, S_t = a, S'_t = b) \leq \theta$  for all  $a, b$  (provided that the condition is of non-zero probability). It follows that

$$\begin{aligned} \mathbb{P}(M_{t+n} = 0 \mid M_t = 0) &\leq \theta \quad \text{for all } t; \\ \mathbb{P}(M_{t+n} = 0) &\leq \theta \cdot \mathbb{P}(M_t = 0) \quad \text{for all } t; \\ \mathbb{P}(M_{jn} = 0) &\leq \theta^j \quad \text{for all } j \end{aligned}$$

and, of course,  $\theta^j \rightarrow 0$  as  $j \rightarrow \infty$ .  $\square$

Interestingly,  $\varepsilon_n \rightarrow 0$  exponentially fast. However, the constant  $nk^{2n}$  can be quite large.

## 5b Finite Markov chains

A *Markov chain* (discrete in space and time, and homogeneous in time) is described by a *transition probability matrix*

$$(p(x, y))_{x, y \in V}$$

satisfying

$$p(x, y) \geq 0; \quad \forall x \quad \sum_y p(x, y) = 1.$$

The set  $V$  is assumed to be finite. We turn  $V$  into a graph putting

$$E = \{(x, y) \in V^2 : p(x, y) \neq 0\}$$

and define the probability of a path  $(s_0, \dots, s_n)$  as the product of  $n$  probabilities

$$p(s_0, \dots, s_n) = p(s_0, s_1) \dots p(s_{n-1}, s_n);$$

as before,  $s_0$  must be equal to a given initial point  $x_0 \in V$ . Here are some definitions that depend on the graph only.

A set  $A \subset V$  is *closed* if  $E \cap (A \times (V \setminus A)) = \emptyset$ .

A Markov chain is *irreducible* if  $\emptyset$  and  $V$  are the only closed sets. In other words: the graph is strongly connected. Equivalently: for all  $x, y \in V$  there exists a path from  $x$  to  $y$  (recall 5a5).

An irreducible Markov chain is *aperiodic*, if there exists no  $p \in \{2, 3, \dots\}$  such that every loop length is divisible by  $p$ . (This property does not depend on the initial point; recall the proof of 5a6.)

Here are some results stated here without proofs.

**5b1 Theorem.** If a Markov chain is irreducible and aperiodic then the limit

$$\lim_n \mathbb{P}(S_n = x)$$

exists for each  $x \in V$ .

**5b2 Definition.** A probability measure  $\mu$  on  $V$  is *stationary*, if

$$\mu(y) = \sum_{x \in V} \mu(x)p(x, y) \quad \text{for all } y \in V.$$

Irreducibility implies that  $\mu(x) > 0$  for all  $x$  (since the set  $\{x : \mu(x) > 0\}$  is closed).

**5b3 Theorem.** If a Markov chain is irreducible and aperiodic then it has one and only one stationary probability measure  $\mu$ , and

$$\forall y \quad \sum_{x \in V} \nu(x)p_n(x, y) \rightarrow \mu(y) \quad \text{as } n \rightarrow \infty$$

for every probability measure  $\nu$  on  $V$ .

If a Markov chain  $(V, p)$  is irreducible but periodic, with the (least) period  $d$ , then  $V = V_1 \uplus \dots \uplus V_d$  and  $p_1(x, y) \neq 0$  only when  $x \in V_i, y \in V_{i+1}$  for some  $i$  (here  $n + 1 = 1$ , of course). The Markov chain  $(V_1, p_d)$  is irreducible and aperiodic, its stationary probability measure is  $\mu(x) = \lim_n \mathbb{P}(S_{nd} = x)$  (assuming  $x_0 \in V_1$ ), and the measure

$$\nu(x) = \lim_n \frac{1}{d} (\mathbb{P}(S_{nd} = x) + \mathbb{P}(S_{nd+1} = x) + \dots + \mathbb{P}(S_{nd+d-1} = x))$$

is stationary for the original Markov chain  $(V, p)$ .

Here is another property related to the graph only.

**5b4 Definition.** <sup>1</sup> A state  $x \in V$  is *transient*, if there exists  $y \in V$  such that a path from  $x$  to  $y$  exists, but a path from  $y$  to  $x$  does not exist. Otherwise,  $x$  is called *recurrent*.

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<sup>1</sup>Only for *finite* Markov chains.

If  $x$  is transient then  $\mathbb{P}(S_n = x) \rightarrow 0$  as  $n \rightarrow \infty$ .

Recurrent states  $x, y$  are called equivalent, if there exists a path from  $x$  to  $y$ , and a path from  $y$  to  $x$ . (Well, the latter follows from the former.) Equivalence classes are irreducible closed sets...

## 5c Return time

Similarly to Sect. 5a we consider a regular (weakly) connected finite directed (but maybe periodic) graph  $(V, E)$ , and the random walk  $(S_n)$  on it, starting at a given  $x_0 \in V$ .

We introduce the “return time” random variable<sup>1</sup>  $T = \inf\{n > 0 : S_n = x_0\}$ .

**5c1 Lemma.**  $T < \infty$  almost surely, and moreover,  $\mathbb{E}T < \infty$ .

*Proof.*

$$\begin{aligned} \exists n \forall t \quad \mathbb{P}(T \leq t + n | S_0, \dots, S_t) &> 0 \text{ a.s.}; \\ \exists n \exists \varepsilon \forall t \quad \mathbb{P}(T \leq t + n | S_0, \dots, S_t) &\geq \varepsilon \text{ a.s.}; \\ \mathbb{P}(T > t + n | S_0, \dots, S_t) &\leq (1 - \varepsilon)\mathbb{1}_{T > t} \text{ a.s.}; \\ \mathbb{P}(T > t + n) &\leq (1 - \varepsilon)\mathbb{P}(T > t); \\ \forall j \quad \mathbb{P}(T > jn) &\leq (1 - \varepsilon)^j. \end{aligned}$$

□

Treating the (one step) transition function  $p(\cdot, \cdot)$  as a matrix and measures on  $V$  as row vectors we write  $\mu p = \nu$  rather than  $\nu(\{y\}) = \sum_x \mu(\{x\})p(x, y) = \int p(\cdot, y) d\mu$ , and in particular,  $\delta_x p$  rather than  $\sum_y p(x, y)\delta_y$ . Thus, distributions of  $S_n$  are:  $\text{Distr}(S_0) = \delta_{x_0}$ ,  $\text{Distr}(S_1) = \delta_{x_0} p$ , and so on. We also use expectations of random vectors (in the  $m$ -dimensional linear space of signed measures on  $V$ ):  $\text{Distr}(S_n) = \mathbb{E} \delta_{S_n}$  (and in general,  $\text{Distr}(X) = \mathbb{E} \delta_X$ ).

“The cycle trick”:  $\sum_{n=0}^{T-1} \delta_{S_n} = \sum_{n=1}^T \delta_{S_n}$  a.s. (just because  $S_0 = x_0 = S_T$  a.s.).

**5c2 Lemma.**  $\mathbb{E} \sum_{n=1}^T \delta_{S_n} = (\mathbb{E} \sum_{n=0}^{T-1} \delta_{S_n})p$ .

*Proof.*

$$\begin{aligned} \mathbb{E}(\delta_{S_{n+1}} - \delta_{S_n} p | S_0, \dots, S_n) &= 0; \\ \mathbb{E}((\delta_{S_{n+1}} - \delta_{S_n} p)\mathbb{1}_{T > n}) &= 0; \end{aligned}$$

<sup>1</sup>A priori, taking on values in  $\{1, 2, \dots\} \cup \{\infty\}$ .

taking into account that  $\sum_{n=0}^{\infty} \mathbb{P}(T > n) = \mathbb{E}T < \infty$  (and vectors  $\delta_{S_{n+1}} - \delta_{S_n}p$  are a bounded set) we have

$$\begin{aligned} \mathbb{E} \sum_{n=0}^{\infty} (\delta_{S_{n+1}} - \delta_{S_n}p) \mathbb{1}_{T>n} &= 0; \\ \mathbb{E} \sum_{n=0}^{T-1} (\delta_{S_{n+1}} - \delta_{S_n}p) &= 0; \\ \mathbb{E} \sum_{n=0}^{T-1} \delta_{S_{n+1}} &= \left( \mathbb{E} \sum_{n=0}^{T-1} \delta_{S_n} \right) p. \end{aligned}$$

□

**Proof of Prop. 5a2.** The measure  $\mathbb{E} \sum_{n=0}^{T-1} \delta_{S_n}$  is invariant, therefore, proportional to the uniform (or the counting) measure. The measure at  $x_0$  is equal to 1 ( $n = 0$  only...); thus the measure of the whole  $V$  must be  $m$ . On the other hand, it is  $\mathbb{E} \sum_{n=0}^{T-1} 1 = \mathbb{E}T$ ; thus,  $\mathbb{E}T = m$ . □