



CS 3200 – Introduction to Scientific Computing

Instructor: Chris Johnson

Topic: Conjugate Gradient (CG) Method

For Symmetric Positive-Definite Matrix A

- Solve the system Ax = b
- Graph of the quadratic form $f(x) = \frac{1}{2}x^T A x b^T x + c$:



• The minimum point of the surface is the solution to Ax = b

from J. R. Shewchuk "Painless CG"

• Quadratic function of vector *x*:

$$f(x) = \frac{1}{2}x^T \mathbf{A}x - b^T x + c$$

• Derivative:

$$f'(x) = \frac{1}{2}\mathbf{A}^T x + \frac{1}{2}\mathbf{A}x - b$$

• If A is symmetric

$$f'(x) = \mathbf{A}x - b$$

- If A is also positive-definite (i.e., $x^T A x > 0$ for any nonzero vector x)
 - Minimize f(x) by setting f'(x) to 0

$$\mathbf{A}x = b$$

Gradient of Quadratic Form



For A Not Positive-Definite



Non-Stationary Iterative Method

 Start from initial guess x₀; adjust until close enough to the exact solution:

$$x_{i+1} = x_i + a_i p_i$$

- *p_i* Adjustment Direction
- a_i Step Size
- How to choose direction and step size?

 Choose the direction in which *f* decreases most quickly – the direction opposite to *f*'(*x_i*):

$$r_i = -f'(x_i) = b - Ax_i$$

• Which is also the direction of the residual:

$$x_{i+1} = x_i + a_i r_i$$

- Direction of steepest descent is plane in 3 dimensions
- Intersection of plane with function produces parabola
- Minimum of parabola determines desired step size



from J. R. Shewchuk "Painless CG"

- How to choose step size? *a_i* should minimize *f*, along the direction of *r_i*, which means

$$\frac{d}{da}f(x_{i+1}) = 0$$

$$\frac{d}{da}f(x_{i+1}) = f'(x_{i+1})^T \frac{d}{da}x_{i+1} = f'(x_{i+1})^T r_i = 0$$

$$\Rightarrow r_{i+1}^{T} r_{i} = 0$$

$$\Rightarrow (b - Ax_{i+1})^{T} r_{i} = 0$$

$$\Rightarrow (b - A(x_{i} + a_{i}r_{i}))^{T} r_{i} = 0$$

$$\Rightarrow (b - Ax_{i})^{T} r_{i} = a_{i}(Ar_{i})^{T} r_{i}$$

$$\Rightarrow a_{i} = \frac{r_{i}^{T} r_{i}}{r_{i}^{T} Ar_{i}}$$

Steepest Descent Algorithm

 Given x₀, iterate until the residue is smaller than the error tolerance:

$$r_i = b - Ax_i$$
$$a_i = \frac{r_i^T r_i}{r_i^T A r_i}$$

$$x_{i+1} = x_i + a_i r_i$$

Iterations of Steepest Descent Method



from J. R. Shewchuk "Painless CG"

- The steepest descent method is very reliable - it can always make progress provided the gradient is nonzero
- Depending on the function, however, steepest descents can take a long time to converge to the true solution
- Instead of looking at a gradient direction, we can speed up minimization by searching in the conjugate direction
 - The error after *i* steps is orthogonal to all previous search directions

Conjugate Gradient (CG) Method



Conjugate Gradient Method

- The search directions are orthogonal (socalled A-orthogonal or conjugate)
- The error is minimal over the space spanned by the search directions
- Minimum error implies reaching an exact solution in at most n steps
 - Because of rounding errors, there is a loss of orthogonality

Orthogonal Directions

• Pick orthogonal search directions: d_0, d_1, \dots, d_{n-1}

$$x_{i+1} = x_i + a_i d_i$$
$$d_i^T e_{i+1} = 0$$
$$d_i^T (e_i + a_i d_i) = 0$$
$$a_i = -\frac{d_i^T e_i}{d_i^T d_i}$$

• Problem – we don't know e_i

A-Orthogonal Directions

 Instead of orthogonal search directions, we make the search directions A-orthogonal (conjugate)



These pairs of vectors are A-orthogonal because these pairs of vectors are orthogonal.

Conjugate Gradient Method

• New search directions:

$$x_{i+1} = x_i + \alpha_i d_i$$

• Optimum orthogonality:

$$0 = d_i^T e_{i+1}$$

- We say that two non-zero vectors u and v are conjugate (with respect to A) if $u^T A v = 0$
- Look at the conjugation:

$$0 = d_i^T A e_{i+1}$$

Search Step Size

•
$$\frac{d}{da}f(x_{i+1}) = f'(x_{i+1})^T \frac{d}{da}x_{i+1} = f'(x_{i+1})^T d_i = 0$$

$$\Rightarrow r_{i+1}^T d_i = 0$$

$$\Rightarrow (b - Ax_{i+1})^T d_i = 0$$

$$\Rightarrow (b - A(x_i + a_i d_i))^T d_i = 0$$

$$\Rightarrow (b - Ax_i)^T d_i = a_i (Ad_i)^T d_i$$

$$\Rightarrow a_i = \frac{r_i^T d_i}{d_i^T A d_i}$$

Conjugate Gradient Algorithm

Start with

$$d_0 = r_0 = b - Ax_0$$

• Iterate:

$$\begin{aligned} \alpha_{i} &= \frac{d_{i}^{T} r_{i}}{d_{i}^{T} A d_{i}} & \leftarrow \text{scalar} \\ x_{i+1} &= x_{i} + \alpha_{i} d_{i} & \leftarrow \text{vector} \\ r_{i+1} &= r_{i} - \alpha_{i} A d_{i} & \leftarrow \text{vector} \\ \beta_{i+1} &= \frac{r_{i+1}^{T} r_{i+1}}{r_{i}^{T} r_{i}} & \leftarrow \text{scalar} \\ d_{i+1} &= r_{i+1} + \beta_{i+1} d_{i} & \leftarrow \text{vector} \end{aligned}$$

Conjugate Gradient Algorithm

• Initialization:

$$x_{0} = \begin{bmatrix} 0 & 0 \end{bmatrix}'$$
$$d_{0} = r_{0} = b - Ax_{0}$$
Iteration:
$$\alpha_{0} = \frac{d_{0}^{T}r_{0}}{d_{0}^{T}Ad_{0}}$$
$$x_{1} = x_{0} + \alpha_{0}d_{0}$$
$$r_{1} = r_{0} - \alpha_{0}Ad_{0}$$
$$\beta_{1} = \frac{r_{1}^{T}r_{1}}{r_{0}^{T}r_{0}}$$
$$d_{1} = r_{1} + \beta_{1}d_{0}$$
$$\alpha_{1} = \frac{d_{1}^{T}r_{1}}{d_{1}^{T}Ad_{1}}$$
$$x_{2} = x_{1} + \alpha_{1}d_{1}$$



Preconditioning

- In exact arithmetic, CG converges in n steps (completely unrealistic)
- The convergence rate of CG can be accelerated by preconditioning
- Apply the CG algorithm to M⁻¹A
 Choose M so that M⁻¹A is better conditioned and systems of the form My = z are easily solved • Then solve $(M^{-1}A)x = M^{-1}b$
- Possible choices for preconditioning include
 - Diagonalization
 - Incomplete LU factorization
 - Approximate inverse
 - SSOR

Conjugate Gradient – Other resources

- Shewchuk's introduction to steepest descent and CG (cited in these slides):
 - <u>http://www.cs.cmu.edu/~quake-papers/painless-conjugate-gradient.pdf</u>
- Another set of slides on CG, with information on preconditioners:
 - <u>http://www.stanford.edu/class/ee364b/lectures/</u> <u>conj_grad_slides.pdf</u>
- A good online text for iterative methods (see section 5.3 for steepest descent & 6.7 for CG):
 - <u>http://www-users.cs.umn.edu/~saad/IterMethBook_2ndEd.pdf</u>

Conjugate Gradient – Other resources

- Applets illustrating the steepest descent and CG methods, from a site with an <u>extensive collection of applets</u> for scientific computing concepts and methods:
 - <u>http://www.cse.illinois.edu/iem/optimization/SteepestDescent/</u>
 - <u>http://www.cse.illinois.edu/iem/optimization/ConjugateGradient/</u>

Steepest Descent

- Suppose that we want to find a local minimum for the scalar function f of the vector variable x, starting from an initial point x_0 . Picking an appropriate x_0 is crucial but also very problem-dependent. We start from x_0 and go downhill. At every step of the way, we must make the following decisions:
 - When to stop?
 - In what direction to proceed?
 - How long a step to take?

Quadratic function of vector x

$$f(x) = \frac{1}{2}x^T A x - b^T x + c$$

- If $x^T \mathbf{A} x > 0$ for any nonzero vector x
 - Matrix *A* is positive-definite
- If *A* is symmetric and positive-definite
 - f(x) is minimized by the solution Ax = b

Symmetric Positive-Definite Matrix

- If A is symmetric and positive-definite:
 - p is an arbitrary point
 - x is the solution point

$$x = A^{-1}b$$

$$f(p) = f(x) + \frac{1}{2}(p - x)^{T}A(p - x)$$

Since

$$\frac{1}{2}(p-x)^T A(p-x) > 0$$

• We have

f(p) > f(x)

Steepest Descent: Example

•
$$\begin{bmatrix} 3 & 2 \\ 2 & 6 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 2 \\ -8 \end{bmatrix}$$

- a) Starting at (-2, -2), take the direction of steepest descent of *f*
- b) Find the point on the intersection of these two surfaces that minimizes *f*
- c) Intersection of surfaces
- d) The gradient at the bottommost point is orthogonal to the gradient of the previous step



Case Study

- a) Large κ , small μ
- b) An example of poor convergence $-\kappa$ and μ and both large
- c) Small κ and μ
- d) Small κ , large μ



Other Krylov Subspace Methods

- Nonsymmetric linear systems:
 - GMRES:
 - for i = 1, 2, 3, . . .

find $x_i \in K_i$ (A, b) such that $r_i = (Ax_i - b) \perp K_i$ (A, b) But, no short recurrence => save old vectors => lots more space (usually "restarted" every k iterations to use less space)

- BiCGStab, QMR, etc.: Two spaces K_i (A, b) and K_i (A^T, b) w/ mutually orthogonal bases Short recurrences => O(n) space, but less robust
- Convergence and preconditioning more delicate than CG
- Active area of current research
- Eigenvalues: Lanczos (symmetric), Arnoldi (nonsymmetric)

Conjugate Gradient: Convergence

- In exact arithmetic, CG converges in n steps (completely unrealistic)
- Accuracy after k steps of CG is related to:
 - Consider polynomial of degree k that is equal to 1 at step 0
 - How small can such a polynomial be at all the eigenvalues of A?
- Condition number: $\kappa(A) = \|A\|_2 \|A^{-1}\|_2 = \lambda_{max}(A) / \lambda_{min}(A)$
- Residual is reduced by a constant factor over $\mathcal{O}(K^{1/2}(A))$ iterations of CG

Preconditioning

- The convergence rate of CG can be accelerated by preconditioning
- Apply the CG algorithm to M⁻¹A where M is chosen so that M⁻¹A is better conditioned and systems in the form of Mz=y are easily solved
- Possible choices for preconditioners include:
 - Diagonal
 - Incomplete LU factorization
 - Approximate inverse
 - SSOR

Preconditioners

- Suppose you had a matrix B such that:
 - 1. Condition number κ (B⁻¹A) is small
 - 2. By = z is easy to solve
- Then you could solve (B⁻¹A)x = B⁻¹b instead of Ax = b
- B = A is great for (1), not for (2)
- B = I is great for (2), not for (1)
- Domain-specific approximations sometimes work
- B = diagonal of A sometimes works
- Better: blend in some direct-methods ideas...

$$\mathbf{x}_0 = 0, r_0 = b, d_0 = B^{-1}r_0, y_0 = B^{-1}r_0$$

for
$$k = 1,2,3,...$$

 $\alpha k = (y_{k-1}^T r_{k-1}) / (d_{k-1}^T A d_{k-1})$ step length
 $x_k = x_{k-1} + \alpha_k d_{k-1}$ approx solution
 $r_k = r_{k-1} - \alpha_k A d_{k-1}$ residual
 $y_k = B^{-1} r_k$ preconditioning solve
 $\beta_k = (y_k^T r_k) / (y_{k-1}^T r_{k-1})$ improvement
 $d_k = y_k + \beta_k d_{k-1}$ search direction

- One matrix-vector multiplication per iteration
- One solve with preconditioner per iteration

Enter the Conjugate Gradient Method

 The conjugate gradient method was originally proposed in Magnus R. Hestenes and Eduard Stiefel (1952), Methods of conjugate gradients for solving linear systems, J. Research Nat. Bur. Standards 49, 409–436.